



⑪ Publication number : 0 622 741 A2

12

EUROPEAN PATENT APPLICATION

② Application number : 94302323.4

(51) Int. Cl.⁵: G06F 15/332, H04N 7/13,
G06F 15/64

(22) Date of filing : 30.03.94

The application is published incomplete as filed (Article 93 (2) EPC). The point in the description or the claim(s) at which the omission obviously occurs has been left blank.

③ Priority : 30.03.93 US 40301
30.07.93 US 100747
01.10.93 US 130571

④3 Date of publication of application :
02.11.94 Bulletin 94/44

84 Designated Contracting States :
AT BE CH DE DK ES FR GB GR IE IT U LU MC
NL PT SE

71 Applicant : KLICS, Ltd.
P.P. Box 570,
No.1, Le Courteur Court,
Mulcaster Street,
St Helier
Jersey JE4 8X2, Channel Islands (GB)

(72) Inventor : Knowles, Gregory P.
Calle Menorca 18-2-B
E-07011 Palma (ES)

⑦ Representative : Jones, Ian
W.P. THOMSON & CO.
Calcon House
289-293 High Holborn
London WC1V 7HU (GB)

(54) Device and method for data compression/decompression.

(57) An apparatus produces an encoded and compressed digital data stream from an original input digital data stream using a forward discrete wavelet transform and a tree encoding method. The input digital data stream may be a stream of video image data values in digital form. The apparatus is also capable of producing a decoded and decompressed digital data stream, closely resembling the originally input digital data stream from an encoded and compressed digital data stream using a corresponding tree decoding method and a corresponding inverse discrete wavelet transform. A dual convolver is disclosed which performs both boundary and nonboundary filtering for forward transform discrete wavelet processing and which also performs filtering of corresponding inverse transform discrete wavelet processes. A portion of the dual convolver is also usable to filter an incoming stream of digital video image data values before forward discrete wavelet processing. Methods and structures for generating the addresses to read/write data values from/to memory as well as for reducing the total amount of memory necessary to store data values are also disclosed.

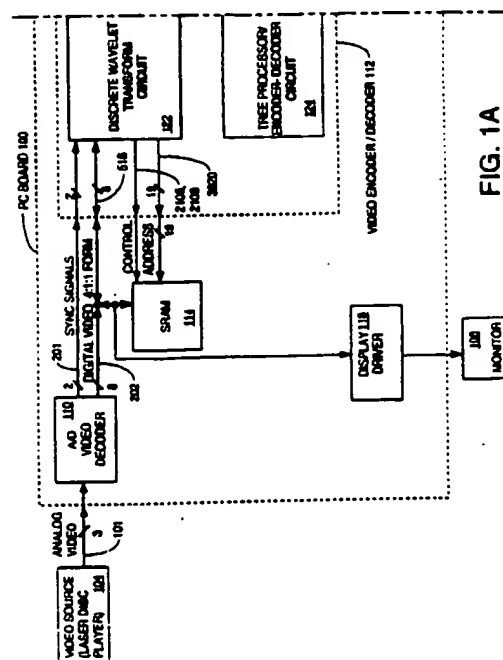


FIG 1A

EP 0 622 741 A2

Best Available Copy
~~This Page Blank (uspto)~~

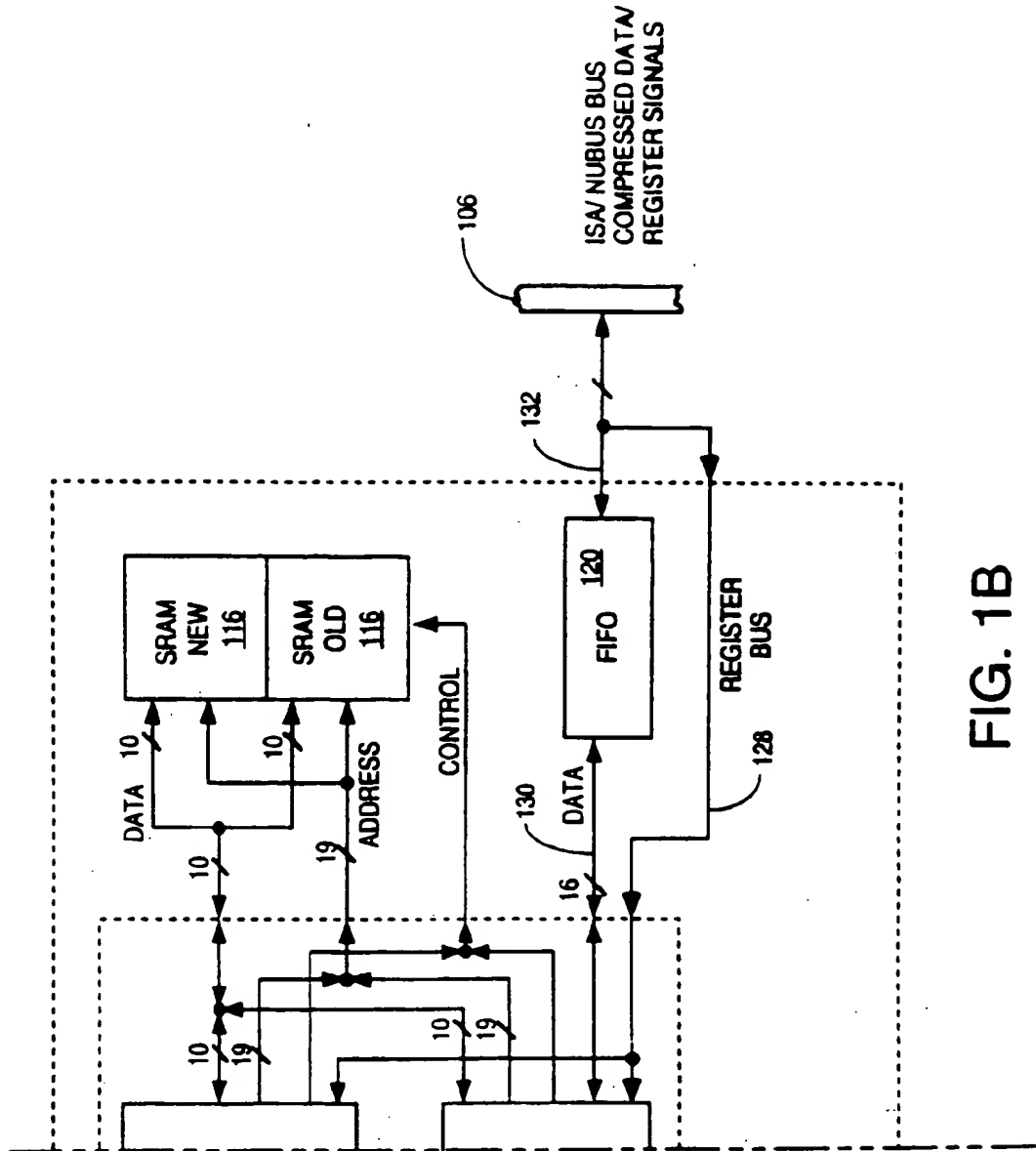


FIG. 1B

CROSS REFERENCE TO PAPER APPENDICES

Appendix A, which is a part of the present disclosure, is a paper appendix of 6 pages. Appendix A is a description of a CONTROL_ENABLE block contained in the tree processor/encoder-decoder portion of a video encoder/decoder integrated circuit chip, written in the VHDL hardware description language.

Appendix B, which is a part of the present disclosure, is a paper appendix of 10 pages. Appendix B is a description of a MODE_CONTROL block contained in the tree processor/encoder-decoder portion of a video encoder/decoder integrated circuit chip, written in the VHDL hardware description language.

Appendix C, which is a part of the present disclosure, is a paper appendix of 11 pages. Appendix C is a description of a CONTROL_COUNTER block contained in the tree processor/encoder-decoder portion of a video encoder/decoder integrated circuit chip, written in the VHDL hardware description language.

A portion of the disclosure of this patent document contains material which is subject to copyright protection. The copyright owner has no objection to the facsimile reproduction by anyone of the patent document or the patent disclosure, as it appears in the Patent and Trademark Office patent files or records, but otherwise reserves all copyright rights whatsoever. The VHDL hardware description language of Appendices A, B and C is an international standard, IEEE Standard 1076-1987, and is described in the "IEEE Standard VHDL Language Reference Manual".

Appendix D, which is a part of the present disclosure, is a paper appendix of 181 pages. Appendix D is a description of one embodiment of a video encoder/decoder integrated circuit chip in the VHDL hardware description language. The VHDL hardware description language of Appendix D is an international standard, IEEE Standard 1076-1987, and is described in the "IEEE Standard VHDL Language Reference Manual". The "IEEE Standard VHDL Language Reference Manual" can be obtained from the Institute of Electrical and Electronics Engineers, Inc., 445 Hoes Lane, Piscataway, New Jersey 08855, telephone 1-800-678-4333.

DESCRIPTION

This invention relates to a method and apparatus for compressing, decompressing, transmitting, and/or storing digitally encoded data. In particular, this invention relates to the compression and decompression of digital video image data.

An apparatus produces an encoded/compressed digital data stream from an original input digital data stream using a discrete wavelet transform and a tree encoding method. The apparatus is also capable of producing a decoded/decompressed digital data stream closely resembling the originally input digital data stream from an encoded/compressed digital data stream using a corresponding tree decoding method and a corresponding inverse discrete wavelet transform.

The apparatus comprises a discrete wavelet transform circuit which is capable of being configured to perform either a discrete wavelet transform or a corresponding inverse discrete wavelet transform. The discrete wavelet transform circuit comprises an address generator which generates the appropriate addresses to access data values stored in memory. Methods and structures for reducing the total amount of memory necessary to store data values and for taking advantage of various types of memory devices including dynamic random access memory (DRAM) devices are disclosed. A convolver circuit of the discrete wavelet transform circuit performs both boundary and non-boundary filtering for the forward discrete wavelet transform and performs start, odd, even and end reconstruction filtering for the inverse discrete wavelet transform. The convolver may serve the dual functions of 1) reducing the number of image data values before subsequent forward discrete wavelet transforming, and 2) operating on the reduced number of image data values to perform the forward discrete wavelet transform.

The apparatus also comprises a tree processor/ encoder-decoder circuit which is configurable in an encoder mode or in a decoder mode. In the encoder mode, the tree processor/encoder-decoder circuit generates addresses to traverse trees of data values of a sub-band decomposition, generates tokens, and quantizes and Huffman encodes selected transformed data values stored in memory. In the decoder mode, the tree processor/decoder-encoder circuit receives Huffman encoded data values and tokens, Huffman decodes and inverse quantizes the encoded data values, recreates trees of transformed data values from the tokens and data values, and stores the recreated trees of data values in memory.

The apparatus is useful in, but not limited to, the fields of video data storage, video data transmission, television, video telephony, computer networking, and other fields of digital electronics in which efficient storage and/or transmission and/or retrieval of digitally encoded data is needed. The apparatus facilitates the efficient and inexpensive compression and storage of video and/or audio on compact laser discs (commonly known as CDs) as well as the efficient and inexpensive storage of video and/or audio on digital video tapes (commonly known as VCR or "video cassette recorder" tapes). Similarly, the invention facilitates the efficient

and inexpensive retrieval and decompression of video and/or audio from digital data storage media including CDs and VCR tapes.

The invention is further described below, by way of example, with reference to the accompanying drawings, in which:

Figure 1 is a block diagram of an expansion printed circuit board which is insertable into a card slot of a personal computer.

Figure 2 is a block diagram of an embodiment of the analog/digital video decoder chip depicted in Figure 1.

Figures 3A-C illustrate a 4:1:1 luminance-chrominance-chrominance format (Y:U:V) used by the expansion board of Figure 1.

Figure 4 is an illustration of a timeline of the output values output from the analog/digital video decoder chip of Figures 1 and 2.

Figure 5 is a block diagram of the discrete wavelet transform circuit of the video encoder/decoder chip of Figure 1.

Figure 6 is a block diagram of the row convolver block of Figure 5.

Figure 7 is a block diagram of the column convolver block of Figure 5.

Figure 8 is a block diagram of the wavelet transform multiplier circuit blocks of Figures 6 and 7.

Figure 9 is a block diagram of the row wavelet transform circuit block of Figure 6.

Figure 10 is a diagram illustrating control signals which control the row convolver of Figure 5 and signals output by the row convolver of Figure 5 during a forward octave 0 transform.

Figure 11 is a diagram showing data flow in the row convolver of Figure 5 during a forward octave 0 transform.

Figure 12 is a diagram illustrating data values output by the row convolver of Figure 5 during the forward octave 0 transform.

Figure 13 is a block diagram of the column wavelet transform circuit block of Figure 7.

Figure 14 is a diagram illustrating control signals which control the column convolver of Figure 5 and signals output by the column convolver of Figure 5 during a forward octave 0 transform.

Figure 15 is a diagram showing data flow in the column convolver of Figure 5 during a forward octave 0 transform.

Figure 16 is a diagram illustrating data values present in memory unit 116 of Figure 1 after operation of the column convolver of Figure 5 during the forward octave 0 transform.

Figure 17 is a diagram showing control signals controlling the row convolver of Figure 5 and signals output by the row convolver of Figure 5 during a forward octave 1 transform.

Figure 18 is a diagram showing data flow in the row convolver of Figure 5 during a forward octave 1 transform.

Figure 19 is a diagram showing control signals controlling the column convolver of Figure 5 and signals output by the column convolver of Figure 5 during a forward octave 1 transform.

Figure 20 is a diagram showing data flow in the column convolver of Figure 5 during a forward octave 1 transform.

Figure 21 is a block diagram of one embodiment of the control block 506 of the discrete wavelet transform circuit of Figure 5.

Figure 22 is a diagram showing control signals controlling the column convolver of Figure 5 and signals output by the column convolver of Figure 5 during an inverse octave 1 transform.

Figure 23 is a diagram showing data flow in the column convolver of Figure 5 during a forward octave 1 transform.

Figure 24 is a diagram showing control signals controlling the row convolver of Figure 5 and signals output by the row convolver of Figure 5 during an inverse octave 1 transform.

Figure 25 is a diagram showing data flow in the row convolver of Figure 5 during an inverse octave 1 transform.

Figure 26 is a diagram showing control signals controlling the column convolver of Figure 5 and signals output by the column convolver of Figure 5 during an inverse octave 0 transform.

Figure 27 is a diagram showing data flow in the column convolver of Figure 5 during an inverse octave 0 transform.

Figure 28 is a diagram showing control signals controlling the row convolver of Figure 5 and signals output by the row convolver of Figure 5 during an inverse octave 0 transform.

Figure 29 is a diagram showing data flow in the row convolver of Figure 5 during an inverse octave 0 transform.

Figure 30 is a block diagram of the DWT address generator block of the discrete wavelet transform circuit

of Figure 5.

Figure 31 is a block diagram of the tree processor/encoder-decoder circuit 124 of Figure 1, simplified to illustrate an encoder mode.

Figure 32 is a block diagram of the tree processor/encoder-decoder circuit 124 of Figure 1, simplified to illustrate a decoder mode.

Figure 33 is a block diagram of the decide circuit block 3112 of the tree processor/encoder-decoder of Figures 31-32.

Figure 34 is a block diagram of the tree processor address generator TP_ADDR_GEN block 3114 of the tree processor/encoder-decoder of Figures 31-32.

Figure 35 illustrates the state table for the CONTROL_ENABLE block 3420 of the tree processor address generator of Figure 34.

Figure 36 is a graphical illustration of the tree decomposition process, illustrating the states and corresponding octaves of Figure 35.

Figure 37 is a block diagram of the quantizer circuit block 3116 of the tree processor/encoder-decoder of Figures 31-32.

Figure 38 is a block diagram of the buffer block 3122 of the tree processor/encoder-decoder of Figures 31-32.

Figure 39 is a diagram of the buffer block 3122 of Figure 38 which has been simplified to illustrate buffer block 3122 operation in the encoder mode.

Figure 40 illustrates the output of barrel shifter 3912 of buffer block 3122 when buffer block 3122 is in the encoder mode as in Figure 39.

Figure 41 is a diagram of the buffer block 3122 of Figure 38 which has been simplified to illustrate buffer block 3122 operation in the decoder mode.

Figure 42 illustrates a pipelined encoding-decoding scheme used by the tree processor/encoder-decoder 124 of Figures 31 and 32.

Figure 43 is a block diagram of another embodiment in accordance with the present invention in which the Y:U:V input is in a 4:2:2 format.

Figure 44 illustrates a sequence in which luminance data values are read from and written to the new portion of memory unit 116 of the PC board 100 in a first embodiment in accordance with the invention in which memory unit 116 is realized as a static random access memory (SRAM).

Figure 45 illustrates a sequence in which luminance data values are read from and written to the new portion of memory unit 116 of the PC board 100 in a second embodiment in accordance with the present invention in which memory unit 116 is realized as a dynamic random access memory (DRAM).

Figure 46 illustrates a third embodiment in accordance with the present invention in which memory unit 116 of the PC board 100 is realized as a dynamic random access memory and in which a series of static random access memories are used as cache buffers between tree processor/encoder-decoder 124 and memory unit 116.

Figure 47 illustrates a time line of the sequence of operations of the circuit illustrated in Figure 46.

Figure 1 illustrates a printed circuit expansion board 100 which is insertable into a card slot of a personal computer. Printed circuit board 100 may be used to demonstrate features in accordance with various aspects of the present invention. Printed circuit board 100 receives an analog video signal 101 from an external video source 104 (such as a CD player), converts information in the analog video signal into data in digital form, transforms and compresses the data, and outputs compressed data onto a computer data bus 106 (such as an ISA/NUBUS parallel bus of an IBM PC or IBM PC compatible personal computer). While performing this compression function, the board 100 can also output a video signal which is retrievable from the compressed data. This video signal can be displayed on an external monitor 108. This allows the user to check visually the quality of images which will be retrievable later from the compressed data while the compressed data is being generated. Board 100 can also read previously compressed video data from data bus 106 of the personal computer, decompress and inverse-transform that data into an analog video signal, and output this analog video signal to the external monitor 108 for display.

Board 100 comprises an analog-to-digital video decoder 110, a video encoder/decoder integrated circuit chip 112, two static random access memory (SRAM) memory units 114 and 116, a display driver 118, and a first-in-first-out memory 120. Analog-to-digital (A/D) video decoder 110 converts incoming analog video signal 101 into a digital format. Video encoder/decoder chip 112 receives the video signal in the digital format and performs a discrete wavelet transform (DWT) function, and then a Huffman encoding function to produce a corresponding compressed digital data stream. Memory unit 116 stores "new" and "old" DWT-transformed video frames.

Video encoder/decoder chip 112 comprises a discrete wavelet transform circuit 122 and a tree proces-

5 sor/ encoder-decoder circuit 124. The discrete wavelet transform circuit 122 performs either a forward discrete wavelet transformation or an inverse discrete wavelet transformation, depending on whether the chip 112 is configured to compress video data or to decompress compressed video data. Similarly, the tree processor/encoder-decoder circuit 124 either encodes wavelet-transformed images into a compressed data stream or decodes a compressed data stream into decompressed images in wavelet transform form, depending on whether the chip 112 is configured to compress or to decompress video data. Video encoder/decoder chip 112 is also coupled to computer bus 106 via a download register bus 128 so that the discrete wavelet transform circuit 122 and the tree processor/encoder-decoder circuit 124 can receive control values (such as a value indicative of image size) from ISA bus 106. The control values are used to control the transformation, tree processing, and encoding/decoding operations. FIFO buffer 120 buffers data flow between the video encoder/decoder chip 112 and the data bus 106. Memory unit 114 stores a video frame in uncompressed digital video format. Display driver chip 118 converts digital video data from either decoder 110 or from memory unit 114 into an analog video signal which can be displayed on external monitor 108.

10 Figure 2 is a block diagram of analog/digital video decoder 110. Analog/digital video decoder 110 converts the analog video input signal 101 into one 8-bit digital image data output signal 202 and two digital video SYNC output signals 201. The 8-bit digital image output signal 202 contains the pixel luminance values, Y, time multiplexed with the pixel chrominance values, U and V. The video SYNC output signals 201 comprise a horizontal synchronization signal and a vertical synchronization signal.

15 Figures 3A-C illustrate a 4:1:1 luminance-chrominance-chrominance format (Y:U:V) used by board 100. Because the human eye is less sensitive to chrominance variations than to luminance variations, chrominance values are subsampled such that each pixel shares an 8-bit chrominance value U and an 8-bit chrominance value V with three of its neighboring pixels. The four pixels in the upper-left hand corner of the image, for example, are represented by $\{Y_{00}, U_{00}, V_{00}\}$, $\{Y_{01}, U_{00}, V_{00}\}$, $\{Y_{10}, U_{00}, V_{00}\}$, and $\{Y_{11}, U_{00}, V_{00}\}$. The next four pixels to the right are represented by $\{Y_{02}, U_{01}, V_{01}\}$, $\{Y_{03}, U_{01}, V_{01}\}$, $\{Y_{12}, U_{01}, V_{01}\}$, and $\{Y_{13}, U_{01}, V_{01}\}$. A/D video decoder 110 serially outputs all the 8-bit Y-luminance values of a frame, followed by all the 8-bit U-chrominance values of the frame, followed by all the 8-bit V-chrominance values of the frame. The Y, U and V values for a frame are output every 1/30 of a second. A/D video decoder 110 outputs values in raster-scan format so that a row of pixel values $Y_{00}, Y_{01}, Y_{02} \dots$ is output followed by a second row of pixel values $Y_{10}, Y_{11}, Y_{12} \dots$ and so forth until all the values of the frame of Figure 3A are output. The values of Figure 3B are then output row by row and then the values of Figure 3C are output row by row. In this 4:1:1 format, each of the U and V components of the image contains one quarter of the number of data values contained in the Y component.

20 Figure 4 is a diagram of a timeline of the output of A/D video decoder 110. The bit rate of the decoder output is equal to 30 frames/sec x 12 bits/pixel. For a 640 x 400 pixel image, for example, the data rate is approximately 110×10^6 bits/second. A/D video decoder 110 also detects the horizontal and vertical synchronization signals in the incoming analog video input signal 102 and produces corresponding digital video SYNC output signals 201 to the video encoder/decoder chip 112.

25 The video encoder/decoder integrated circuit chip 112 has two modes of operation. It can either transform and compress ("encode") a video data stream into a compressed data stream or it can inverse transform and decompress ("decode") a compressed data stream into a video data stream. In the compression mode, the digital image data 202 and the synchronization signals 201 are passed from the A/D video decoder 110 to the discrete wavelet transform circuit 122 inside the video encoder/decoder chip 112. The discrete wavelet transform circuit 122 performs a forward discrete wavelet transform operation on the image data and stores the resulting wavelet-transformed image data in the "new" portion of memory unit 116. At various times during this forward transform operation, the "new" portion of memory unit 116 stores intermediate wavelet transform results, such that certain of the memory locations of memory unit 116 are read and overwritten a number of times. The number of times the memory locations are overwritten corresponds to the number of octaves in the wavelet transform. After the image data has been converted into a sub-band decomposition of wavelet-transformed image data, the tree processor/encoder-decoder circuit 124 of encoder/decoder chip 112 reads wavelet-transformed image data of the sub-band decomposition from the "new" portion of memory unit 116, processes it, and outputs onto lines 130 a compressed ("encoded") digital data stream to FIFO buffer 120. During this tree processing and encoding operation, the tree processor/encoder-decoder circuit 124 also generates a quantized version of the encoded first frame and stores that quantized version in the "old" portion of memory unit 116. The quantized version of the encoded first frame is used as a reference when a second frame of wavelet-transformed image data from the "new" portion of memory unit 116 is subsequently encoded and output to bus 106. While the second frame is encoded and output to bus 106, a quantized version of the encoded second frame is written to the "old" portion of memory unit 116. Similarly, the quantized version of the encoded second frame in the "old" portion of memory unit 116 is later used as a reference for encoding a third frame of image data.

30 In the decompression mode, compressed ("encoded") data is written into FIFO 120 from data bus 106 and

is read from FIFO 120 into tree processor/encoder-decoder circuit 124 of the video encoder/decoder chip 112. The tree processor/encoder-decoder circuit 124 decodes the compressed data into decompressed wavelet-transformed image data and then stores the decompressed wavelet-transformed image data into the "old" portion of memory unit 116. During this operation, the "new" portion of memory unit 116 is not used. Rather, the tree processor/encoder-decoder circuit 124 reads the previous frame stored in the "old" portion of memory unit 116 and modifies it with information from the data stream received from FIFO 120 in order to generate the next frame. The next frame is written over the previous frame in the same "old" portion of the memory unit 116. Once the decoded wavelet-transformed data of a frame of image data is present in the "old" portion of memory unit 116, the discrete wavelet transform circuit 122 accesses memory unit 116 and performs an inverse discrete wavelet transform operation on the frame of image data. For each successive octave of the inverse transform, certain of the memory locations in the "old" portion of memory unit 116 are read and overwritten. The number of times the locations are overwritten corresponds to the number of octaves in the wavelet transform. On the final octave of the inverse transform which converts the image data from octave-0 transform domain into standard image domain, the discrete wavelet transform circuit 122 writes the resulting decompressed and inverse-transformed image data into memory unit 114. The decompressed and inverse-transformed image data may also be output to the video display driver 118 and displayed on monitor 108.

Figure 5 is a block diagram of the discrete wavelet transform circuit 122 of video encoder/decoder chip 112. The discrete wavelet transform circuit 122 shown enclosed by a dashed line comprises a row convolver block CONV_ROW 502, a column convolver block CONV_COL 504, a control block 506, a DWT address generator block 508, a REGISTERS block 536, and three multiplexers, mux1 510, mux2 512, and mux3 514. In order to transform a frame of digital video image data received from A/D video decoder 110 into the wavelet transform domain, a forward two dimensional discrete wavelet transform is performed. Similarly, in order to return the wavelet transform digital data values of the frame into a digital video output suitable for displaying on a monitor such as 108, an inverse two dimensional discrete wavelet transform is performed. In the presently described embodiment of the present invention, four coefficient quasi-Daubechies digital filters are used as set forth in the copending Patent Cooperation Treaty (PCT) application filed March 30, 1994 entitled "Data Compression and Decompression".

The discrete wavelet transform circuit 122 shown in Figure 5 performs a forward discrete wavelet transform as follows. First, a stream of 8-bit digital video image data values is supplied, one value at a time, to the discrete wavelet transform circuit 122 via eight leads 516. The digital video image data values are coupled through multiplexer mux1 510 to the input leads 518 of the row convolver CONV_ROW block 502. The output leads 520 of CONV_ROW block 502 are coupled through multiplexer mux2 512 to input leads 522 of the CONV_COL block 504. The output leads 524 of CONV_COL block 504 are coupled to data leads 526 of memory unit 116 through multiplexer mux3 so that the data values output from CONV_COL block 504 can be written to the "new" portion of frame memory unit 116. The writing of the "new" portion of memory unit 116 completes the first pass, or octave, of the forward wavelet transform. To perform the next pass, or octave, of the forward wavelet transform, low pass component data values of the octave 0 transformed data values are read from memory unit 116 and are supplied to input leads 518 of CONV_ROW block 502 via input leads 526, lines 528 and multiplexer mux1 510. The flow of data proceeds through row convolver CONV_ROW block 502 and through column convolver CONV_COL block 504 with the data output from CONV_COL block 504 again being written into memory unit 116 through multiplexer mux3 514 and leads 526. Control block 506 provides control signals to mux1 510, mux2 512, mux3 514, CONV_ROW block 502, CONV_COL block 504, DWT address generator block 508, and memory unit 116 during this process. This process is repeated for each successive octave of the forward transform. The data values read from memory unit 116 for the next octave of the transform are the low pass values written to the memory unit 116 on the previous octave of the transform.

The operations performed to carry out the inverse discrete wavelet transform proceed in an order substantially opposite the operations performed to carry out the forward discrete wavelet transform. The frame of image data begins in the transformed state in memory unit 116. For example, if the highest octave in the forward transform (OCT) is octave 1, then transformed data values are read from memory unit 116 and are supplied to the input leads 522 of the CONV_COL block 504 via leads 526, lines 528 and multiplexer mux2 512. The data values output from CONV_COL block 504 are then supplied to the input leads 518 of CONV_ROW block 502 via lines 525 and multiplexer mux1 510. The data values output from CONV_ROW block 502 and present on output leads 520 are written into memory unit 116 via lines 532, multiplexer mux3 514 and leads 526. The next octave, octave 0, of the inverse transform proceeds in similar fashion except that the data values output by CONV_ROW block 502 are the fully inverse-transformed video data which are sent to memory unit 114 via lines 516 rather than to memory unit 116. Control block 506 provides control signals to multiplexer mux1 510, multiplexer mux2 512, multiplexer mux3 514, CONV_ROW block 502, CONV_COL block 504, DWT address generator block 508, memory unit 116, and memory unit 114 during this process.

In both forward wavelet transform and inverse wavelet transform operations, the control block 506 is timed by the external video sync signals 201 received from A/D video decoder 110. Control block 506 uses these sync signals as well as register input values ximage, yimage, and direction to generate the appropriate control signals mentioned above. Control block 506 is coupled to: multiplexer mux1 510 via control leads 550, multiplexer mux2 512 via control leads 552, multiplexer mux3 via control leads 554, CONV_ROW block 502 via control leads 546, CONV_COL block 504 via control leads 548, DWT address generator block 508 via control leads 534, 544, and 556, memory unit 116 via control leads 2108, and memory unit 114 via control leads 2106.

As shown in Figure 5, multiplexer mux1 510 couples one of the following three sets of input signals to input leads 518 of CONV_ROW block 502, depending on the value of control signals on leads 550 supplied from CONTROL block 506: digital video input data values received on lines 516 from A/D video decoder 110, data values from memory unit 116 or data values from multiplexer mux3 514 received on lines 528, or data values from CONV_COL block 504 received on lines 525. Multiplexer mux2 512 couples either the data values being output from row convolver CONV_ROW block 502 or the data values being output from multiplexer mux3 514 received on lines 528 to input leads 522 of CONV_COL block 504, depending on the value of control signals on lead 552 generated by CONTROL block 506. Multiplexer mux3 514 passes either the data values being output from CONV_ROW block 502 received on lines 532 or the data values being output from CONV_COL block 504 onto lines 523 and leads 526, depending on control signals generated by CONTROL block 506. Blocks CONV_ROW 502, CONV_COL 504, CONTROL 506, DWT address generator 508, and REGISTERS 536 of Figure 5 are described below in detail in connection with a forward transformation of a matrix of digital image data values. Lines 516, 532, 528 and 525 as well as input and output leads 518, 520, 522, 524 and 526 are each sixteen bit parallel lines and leads.

Figure 6 is a block diagram of the row convolver CONV_ROW block 502. Figure 7 is a block diagram of the column convolver CONV_COL block 504. Figure 21 is a block diagram of the CONTROL block 506 of Figure 5. Figure 30 is a block diagram of the DWT address generator block 508 of Figure 5.

As illustrated in Figure 6, CONV_ROW block 502 comprises a wavelet transform multiplier circuit 602, a row wavelet transform circuit 604, a delay element 606, a multiplexer MUX 608, and a variable shift register 610. To perform a forward discrete wavelet transform, digital video values are supplied one-by-one to the discrete wavelet transform circuit 122 of the video encoder/decoder chip 112 illustrated in Figure 1. In one embodiment in accordance with the present invention, the digital video values are in the form of a stream of values comprising 8-bit Y (luminance) values, followed by 8-bit U (chrominance) values, followed by 8-bit V (chrominance) values. The digital video data values are input in "raster scan" form. For clarity and ease of explanation, a forward discrete wavelet transform of an eight-by-eight matrix of luminance values Y as described is represented by Table 1. Extending the matrix of Y values to a larger size is straightforward. If the matrix of Y values is an eight-by-eight matrix, then the subsequent U and V matrices will each be four-by-four matrices.

D_{00}	D_{01}	D_{02}	D_{07}
D_{10}	D_{11}	D_{12}	D_{17}
.
.
.
D_{70}	D_{71}	D_{77}

Table 1.

The order of the Y values supplied to the discrete wavelet transform circuit 122 is $D_{00}, D_{01}, \dots, D_{07}$ in the first row, then $D_{10}, D_{11}, \dots, D_{17}$ in the second row, and so forth row by row through the values in Table 1. Multiplexer 510 in Figure 5 is controlled by control block 506 to couple this stream of data values to the row convolver CONV_ROW block 502. The row convolver CONV_ROW block 502 performs a row convolution of the row data values $D_{00}, D_{01}, D_{02}, \dots, D_{07}$ with a high pass four coefficient quasi-Daubechies digital filter $G = (d, c, -b, a)$ and a low pass four coefficient quasi-Daubechies digital filter $H = (a, b, c, -d)$ where $a = 11/32, b = 19/32, c = 5/32, d = 3/32$. The coefficients a, b, c, d are related to a four coefficient Daubechies wavelet as described in the copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression".

The operation of CONV_ROW block 502 on the data values of Table 1 is explained with reference to Figures 6, 8, 9, 10 and 11. Figure 8 is a detailed block diagram of the wavelet transform multiplier circuit 602 of the CONV_ROW block. Figure 9 is a detailed block diagram of the row wavelet transform circuit 604 of the CONV_ROW block. Figure 10 shows a sequence of control signals supplied by the control block 506 of Figure 5 to the row wavelet transform circuit 604 of Figure 9. This sequence of control signals effects a forward one dimensional wavelet transform on the rows of the matrix Table 1. The wavelet transform multiplier circuit 602 of Figure 8 comprises combinatorial logic which multiplies each successive input data value x by various scaled combinations of coefficients 32a, 32b, 32c, and 32d. This combinatorial-logic block comprises shift registers 802, 804, 806, and 808 which shift the multibit binary input data value x to the left by 1, 2, 3, and 4 bits, respectively. Various combinations of these shifted values, as well as the input value x itself, are supplied to multibit adders 810, 812, 814, 816, and 818. The data outputs 32dx, 32(c-d)x, 32cx, 32ax, 32(a+b)x, 32bx, and 32(c+d)x are therefore available to the row wavelet transform circuit 604 on separate sets of leads as shown in detail in Figures 6 and 9.

The row wavelet transform circuit 604 of Figure 9 comprises sets of multiplexers, adders, and delay elements. Multiplexer mux1 902, multiplexer mux2 904, and multiplexer mux3 906 pass selected ones of the data outputs of the wavelet transform multiplier circuit 602 of Figure 8 as determined by control signals on leads 546 from CONTROL block 506 of Figure 5. These control signals on leads 546 are designated muxsel(1), muxsel(2), and muxsel(3) on Figure 9. The remainder of the control signals on leads 546 supplied from CONTROL block 506 to the row wavelet transform circuit 604 comprise andsel(1), andsel(2), andsel(3), andsel(4), addsel(1), addsel(2), addsel(3), addsel(4), muxandsel(1), muxandsel(2), muxandsel(3), centermuxsel(1) and centermuxsel(2).

Figure 10 shows values of the control signals at different times during a row convolution of the forward transform. For example, at time $t=0$, the control input signal to multiplexer mux2 904, muxsel(2), is equal to 2. Multiplexer mux2 904 therefore couples its second input leads carrying the value 32(a+b)x to its output leads. Each of multiplexers 908, 910, 912, and 914 either passes the data value on its input leads, or passes a zero, depending on the value of its control signal. Control signals andsel(1) through andsel(4) are supplied to select input leads of multiplexers 908, 910, 912, and 914, respectively. Multiplexers 916, 918, and 920 have similar functionality. The outputs of multiplexers 916, 918, and 920 depend on the values of control signals muxandsel(1) through muxandsel(3), respectively. Multiplexers 922 and 924 pass either the value on their "left" input leads or the value on their "right" input leads, as determined by control select inputs centermuxsel(1) and centermuxsel(2), respectively. Adder/subtractors 926, 928, 930, and 932 either pass the sum or the difference of the values on their left and right input leads, depending on the values of the control signals addsel(1) through addsel(4), respectively. Elements 934, 936, 938, and 940 are one-cycle delay elements which output the data values that were at their respective input leads during the previous time period.

Figure 11 is a diagram of a data flow through the row convolver CONV_ROW 502 during a forward transform operation on the data values of Table 1 when the control signals 546 controlling the row convolver CONV_ROW 502 are as shown in Figure 10. At the left hand edge of the matrix of the data values of Table 1, start forward low pass and start forward high pass filters G_3 and H_3 are applied in accordance with equations 22 and 24 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" as follows:

$$32H_{00} = 32\{(a+b)D_{00} + cD_{01} - dD_{02}\}$$

$$32G_{00} = 32\{(c+d)D_{00} - bD_{01} + aD_{02}\}$$

The row wavelet transform circuit of Figure 9 begins applying these start forward low and high pass filters when the control signals for this circuit assume the values at time $t=0$ as illustrated in Figure 10.

At time $t=0$, muxsel(2) has a value of 2. Multiplexer mux2 904 therefore outputs the value $32(a+b)D_{00}$ onto its output leads. Muxsel(3) has a value of 3 so multiplexer mux3 906 outputs the value $32(c+d)D_{00}$ into its output leads. Because the control signals andsel(2) and andsel(3) cause multiplexers 910 and 912 to output zeros at $t=0$ as shown in Figure 10, the output leads of adder/subtractor blocks 928 and 930 carry the values $32(a+b)D_{00}$ and $32(c+d)D_{00}$, respectively, as shown in Figure 11. These values are supplied to the input leads of delay elements 936 and 938. Delay elements 936 and 938 in the case of the row transform are one time unit delay elements. The control signals centermuxsel(1) and centermuxsel(2) have no effect at $t=0$, because control signals andsel(2) and andsel(3) cause multipliers 910 and 912 to output zeros.

At time $t=1$, input data value x is the data value D_{01} . Control signal muxsel(2) is set to 1 so that multiplexer mux2 904 outputs the value $32bD_{01}$. The select signal centermuxsel(1) for adder/subtractor block 922 is set to pass the value on its right input leads. The value $32(c+d)D_{00}$, the output of adder/subtractor block 930 at $t=0$, is therefore passed through multiplexer mux4 922 due to the one time unit delay of delay element 938. The control signal andsel(2) is set to pass, so the two values supplied to the adder/subtractor block 928 are $32(c+d)D_{00}$ and $32bD_{01}$. Because the control signal addsel(2) is set to subtract, the value output by adder/sub-

tractor block 928 is $32((c+d)D_{00}-bD_{01})$ as shown in Figure 11. Similarly, with the values of control signals $\text{centermuxsel}(2)$, $\text{andsel}(3)$, $\text{muxsel}(3)$, $\text{muxandsel}(2)$, and $\text{addsel}(3)$ given in Figure 10, the value output by adder/subtractor block 930 is $32((a+b)D_{00} + cD_{01})$ as shown in Figure 11.

At time $t=2$, input data value x is data value D_{02} . The control signals $\text{andsel}(1)$, $\text{muxsel}(1)$, and $\text{muxandsel}(1)$ are set so that the inputs to adder/subtractor block 926 are $32aD_{02}$ and $32((c+d)D_{00}-bD_{01})$. The value $32((c+d)D_{00}-bD_{01})$ was the previous output from adder/subtractor block 928. Because control signal $\text{addsel}(1)$ is set to add as shown in Figure 10, the output of block 926 is $32((c+d)D_{00}-bD_{01}+aD_{02})$ as shown in Figure 11. Similarly, with the value of control signals $\text{addsel}(4)$, $\text{andsel}(4)$ and $\text{muxandsel}(3)$, the value output by adder/subtractor block 932 is $32((a+b)D_{00} + cD_{01} - dD_{02})$ as shown in Figure 11.

As illustrated in Figure 10, output leads OUT2 (which are the output leads of delay element 940) carry a value of $32H_{00}$ at time $t=3$. The value $32((a+b)D_{00}-bD_{01}+aD_{02})$ is equal to $32H_{00}$ because $32H_{00}=32((a+b)D_{00}+cD_{01}-dD_{02})$ as set forth above. Similarly, output leads OUT1 (which are the output leads of delay element 934) carry a value of $32G_{00}$ at $t=3$ because output leads of block 926 have a value of $32((c+d)D_{00}-bD_{01}+aD_{02})$ one time period earlier. Because $32H_{00}$ precedes $32G_{00}$ in the data stream comprising the high and low pass components in a one-dimensional row convolution, delay element 606 is provided in the CONV_ROW row convolver of Figure 6 to delay $32G_{00}$ so that $32G_{00}$ follows $32H_{00}$ on the leads which are input to the multiplexer 608. Multiplexer 608 selects between the left and right inputs shown in Figure 6 as dictated by the value mux_608 , which is provided on one of the control leads 546 from control block 506. The signal mux_608 is timed such that the value $32H_{00}$ precedes the value $32G_{00}$ on the output leads of multiplexer 608.

The output leads of multiplexer 608 are coupled to a variable shift register 610 as shown in Figure 6. The function of the variable shift register 610 is to normalize the data values output from the CONV_ROW block by shifting the value output by multiplexer 608 to the right by m_row bits. In this instance, for example, it is desirable to divide the value output of multiplexer 608 by 32 to produce the normalized values H_{00} and G_{00} . To accomplish this, the value m_row provided by control block 506 via one of the control leads 546 is set to 5. The general rule followed by the control block 506 of the discrete wavelet transform circuit is to: (1) set m_row equal to 5 to divide by 32 during the forward transform, (2) set m_row equal to 4 to divide by 16 during the middle of a row during an inverse transform, and (3) set m_row equal to 3 to divide by 8 when generating a start or end value of a row during the inverse transform. In the example being described, the start values of a transformed row during a forward transform are being generated, so m_row is appropriately set equal to 5.

As illustrated in Figure 10, the $\text{centermuxsel}(1)$ and $\text{centermuxsel}(2)$ control signals alternate such that the values on the right and the left input leads of multiplexers 922 and 924 are passed to their respective output leads for each successive data value convolved. This reverses data flow through the adder/subtractor blocks 928 and 930 in alternating time periods. In time period $t=0$, for example, Figure 11 indicates that the value $32aD_{01}$ in the column designated "Output of Block 926" in time period $t=1$ is added to $32bD_{02}$ to form the value $32(aD_{01} + bD_{02})$ in the column designated "Output of Block 928" at time $t=2$. Then, in time period $t=3$, the value $32(dD_{01} + cD_{02})$ in the column designated "Output of Block 930" is added to $32bD_{03}$ to form the value $32(dD_{01} + cD_{02} - bD_{03})$ in the column designated "Output of Block 928".

Accordingly, in time period $t=2$, the two values supplied to block 928 are $32bD_{02}$ and the previous output from block 926, $32bD_{01}$. Because $\text{addsel}(2)$ is set to add as shown in Figure 10, the value output by block 928 is $32(aD_{01} + bD_{02})$.

Similarly, the output of block 930 is $32(dD_{01} + cD_{02})$. In this way it can be seen the sequence of control signals in Figure 10 causes the circuit of Figure 9 to execute the data flow in Figure 11 to generate, after passage through multiplexer mux_608 and shift register 610 with m_row set equal to 5, the low and high pass non-boundary components H_{01} , G_{01} , H_{02} , and G_{02} . To implement the end forward low and high pass filters beginning at $t=7$ when the last data value of the first row of Table 1, D_{07} , is input to the row convolver, the control signal $\text{muxsel}(2)$ is set to 3, so that $32(b-a)D_{07}$ is passed to block 928. Control signal $\text{muxsel}(3)$ is set to 4, so that $32(c-d)D_{07}$ is passed to block 930. Control signal $\text{addsel}(2)$ is set to subtract and control signal $\text{addsel}(3)$ is set to add. Accordingly, the output of adder/subtractor 928 is $32(dD_{06}+cD_{06}-(b-a)D_{07})$. Similarly, the output of adder/subtractor 930 is $32(aD_{06} + bD_{06} + (c-d)D_{07})$.

As shown in Figure 11, these values are output from blocks 926 and 932 at the next time period when $t=8$ by setting $\text{muxandsel}(1)$ and $\text{muxandsel}(3)$ to be both zero so that adder/subtractor blocks 926 and 932 simply pass the values unchanged. Delay elements 934 and 940 cause the values $32G_{03}$ and $32H_{03}$ to be output from output leads OUT1 and OUT2 at time $t=9$. Multiplexer 608, as shown in Figure 6, selects between the output of delay unit 606 and the OUT2 output as dictated by CONTROL block 506 of Figure 5. Shift register 610 then normalizes the output as described previously, with m_row set equal to 5 for the end of the row. The resulting values G_{03} and H_{03} are the values output by the end low pass and end high pass forward transform digital filters in accordance with equations 26 and 28 of copending Patent Cooperation Treaty (PCT) application filed March

30, 1994, entitled "Data Compression and Decompression". Thus, a three coefficient start forward transform low pass filter and a three coefficient start forward transform high pass filter have generated the values H_{00} and G_{00} . A four coefficient quasi-Daubecheis low pass forward transform filter and a four coefficient quasi-Daubecheis high pass forward transform filter have generated the values $H_{01} \dots G_{02}$. A three coefficient end forward transform low pass filter and a three coefficient end forward transform high pass filter have generated the values H_{03} and G_{03} .

The same sequence is repeated for each of the rows of the matrix in Table 1. In this way, for each two data values input there is one high pass (G) data value generated and there is one low pass (H) data value generated. The resulting output data values of CONV_ROW block 502 are shown in Figure 12.

As illustrated in Figure 5, the values output from row convolver CONV_ROW block 502 are passed to the column convolver CONV_COL block 504 in order to perform column convolution using the same filters in accordance with the method set forth in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression".

Figure 7 is a block diagram of the column convolver CONV_COL block 504 of Figure 5. The CONV_COL block 504 comprises a wavelet transform multiplier circuit 702, a column wavelet transform circuit 704, a multiplexer 708, and a variable shift register 710. In general, the overall operation of the circuit shown in Figure 7 is similar to the overall operation of the circuit shown in Figure 6. The wavelet transform multiplier circuit 702 of the column convolver is identical to the wavelet transform multiplier circuit 602 of Figure 6. The dashed line in Figure 8. Therefore, is designated with both reference numerals 602 and 702.

Figure 13 is a detailed block diagram of the column wavelet transform circuit 704 of Figure 7 of the column convolver. The CONV_COL block 504, as shown in Figure 13, is similar to the CONV_ROW block 502, except that the unitary delay elements 934, 936, 938, and 940 of the CONV_ROW block 502 are replaced by "line delay" blocks 1334, 1336, 1338, and 1340, respectively. The line delay blocks represent a time delay of one row which, in the case of the matrix of the presently described example, is eight time units. In some embodiments in accordance with the present invention, the line delays are realized using random access memory (RAM).

To perform a column convolution on the values of the matrix of Figure 12, the first three values H_{00} , H_{10} , H_{20} of the first column are processed to generate, after a bit shift in shift register 710 of Figure 7, low and high pass values HH_{00} and HG_{00} of Figure 16. The first three values G_{00} , G_{10} , G_{20} of the second column of the matrix of Figure 12 are then processed to likewise produce GH_{00} and GG_{00} , and so on, to produce the top two rows of values of the matrix of Figure 16. Three values in each column are processed because the start low and high pass filters are three coefficient filters rather than four coefficient filters.

Figure 14 is a diagram illustrating control signals which control the column convolver during the forward transform of the data values of Figure 12. Figure 15 is a diagram illustrating data flow through the column convolver. Corresponding pairs of data values are output from line delays 1334 and 1340 of the column wavelet transform circuit 704. For this reason, the low pass filter output values are supplied from the output leads of the adder/subtractor block 1332 at the input leads of line delay 1340 rather than from the output leads of the line delay 1340 so that a single transformed data value is output from the column wavelet transform circuit in each time period. In Figure 14, output data values $32HH_{00} \dots 32GH_{03}$ are output during time periods $t=16$ to $t=23$ whereas output data values $32HG_{00} \dots 32GG_{03}$ are output during time periods $t=24$ to $t=31$, one line delay later. After being passed through multiplexer 708 and variable shift register 710 of Figure 7, the column convolved data values $HH_{00} \dots GH_{03}$ and $HG_{00} \dots GG_{03}$ are written to memory unit 116 under the control of the address generator. After all the data values of Figure 16 are written to memory unit 116, an octave 0 sub-band decomposition exists in memory unit 116.

To perform the next octave of decomposition, only the low pass component HH values in memory unit 116 are processed. The HH values are read from memory unit 116 and passed through the CONV_ROW block 502 and CONV_COL block 504 as before, except that the control signals for control block 506 are modified to reflect the smaller matrix of data values being processed. The line delay in the CONV_COL block 504 is also shortened to four time units because there are now only four low pass component HH values per row. The control signals to accomplish the octave 1 forward row transform on the data values in Figure 16 are shown in Figure 17. The corresponding data flow for the octave 1 forward row transform is shown in Figure 18. Likewise, the control signals to accomplish the octave 1 forward column transform are shown in Figure 19, and the corresponding data flow for the octave 1 forward column transform is shown in Figure 20.

The resulting HHHH, HHHG, HHGH, and HHGG data values output from the column convolver CONV_COL block 504 are sent to memory unit 116 to overwrite only the locations in memory unit 116 storing corresponding HH data values as explained in connection with Figures 17 and 18 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". The result is an octave 1 sub-band decomposition stored in memory unit 116. This process can be performed on

large matrices of data values to generate sub-band decompositions having as many octaves as required. For ease of explanation and illustration, control inputs and dataflow diagrams are not shown for the presently described example for octaves higher than octave 1. However, control inputs and dataflows for octaves 2 and above can be constructed given the method described in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" along with the octave 0 and octave 1 implementation of that method described above.

Figure 21 illustrates a block diagram of one possible embodiment of control block 506 of Figure 5. Control block 506 comprises a counter 2102 and a combinatorial logic block 2104. The control signals for the forward and inverse discrete wavelet transform operations, as shown in Figures 10, 14, 17, 19, 22, 24, 26, and 28, are output onto the output leads of the combinatorial logic block 2104. The input signals to the control block 506 comprise the sync leads 201 which are coupled to A/D video decoder 110, the direction lead 538 which is coupled to REGISTERS block 536, and the image size leads 540 and 542 which are also coupled to REGISTERS block 536. The values of the signals on the register leads 538, 540, and 542 are downloaded to REGISTERS block 536 of the video encoder/decoder chip 112 from data bus 106 via register download bus 128. The output leads of control block 506 comprise CONV_ROW control leads 546, CONV_COL control leads 548, DWT control leads 550, 552, and 554, memory control leads 2106 and 2108, DWT address generator mux control lead 556, DWT address generator read control leads 534, and DWT address generator write control leads 544.

Counter block 2102 generates the signals row_count, row_carry, col_count, col_carry, octave, and channel, and provides these signals to combinatorial logic block 2104. Among other operations, counter 2102 generates the signals row_count and row_carry by counting the sequence of data values from 0 up to ximage, where ximage represents the horizontal dimension of the image received on leads 540. Similarly, counter 2102 generates the signals col_count and col_carry by counting the sequence of data values from 0 up to yimage, where yimage represents the vertical dimension of the image received on leads 542. The inputs to combinatorial logic block 2104 comprise the outputs of counter block 2102 as well as the inputs direction, ximage, yimage and sync to control block 506. The output control sequences of combinatorial logic block 2104 are combinatorially generated from the signals supplied to logic block 2104.

After the Y data values of an image have been transformed, the chrominance components U and V of the image are transformed. In the presently described specific embodiment of the present invention, a 4:1:1 format of Y:U:V values is used. Each of the U and V matrices of data values comprises half the number of rows and columns as does the Y matrix of data values. The wavelet transform of each of these components of chrominance is similar to the transformation of the Y data values except the line delays in the CONV_COL are shorter to accommodate the shorter row length and the size of the matrices corresponding to the matrix of Table 1 is smaller.

Not only does the discrete wavelet transform circuit of Figure 5 transform image data values into a multi-octave sub-band decomposition using a forward discrete wavelet transformation, but the discrete wavelet transform circuit of Figure 5 can be used to perform a discrete inverse wavelet transform on transformed-image data to convert a sub-band decomposition back into the image domain. In one octave of an inverse discrete wavelet transform, the inverse column convolver 504 of Figure 5 operates on transformed-image data values read from memory unit 116 via leads 526, lines 528 and multiplexer mux2 512 and the inverse row convolver 502 operates on the data values output by the column convolver supplied via leads 524, lines 525 and multiplexer mux1 510.

Figures 22 and 23 show control signals and data flow for the column convolver 504 of Figure 5 when column convolver 504 performs an inverse octave 1 discrete wavelet transform on transformed-image data located in memory unit 116. As illustrated in Figure 23, the data value output from adder/subtractor block 1326 of Figure 13 at time $t=4$ is $32((b-a)HHHH_{00} + (c-d)HHHG_{00})$. The column convolver therefore processes the first two values $HHHH_{00}$ and $HHHG_{00}$ in accordance with the two coefficient start reconstruction filter (inverse transform filter) set forth in equation 52 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". Subsequently, blocks 1332 and 1326 output values indicating that the column convolver performs the four coefficient odd and even reconstruction filters (interleaved inverse transform filters) of equations 20 and 19 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". Fig. 23 illustrates that the column convolver performs the two coefficient end reconstruction filter (inverse transform filter) on the last two data values $HHHH_{10}$ and $HHHG_{10}$ (see time $t=20$) of the first column of transformed data values in accordance with equation 59 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". The data values output from the column convolver of Figure 13 are supplied to the row convolver 502 of Figure 5 via lines 525 and multiplexer mux1 510.

Figures 24 and 25 show control signals and data flow for the row convolver 502 of Figure 5 when the row convolver performs an inverse octave 1 discrete wavelet transform on the data values output from the column

convolver. The column convolver 504 has received transformed values $HHHH_{00} \dots HHGH_{01}$ and so forth as illustrated in Fig. 23 and generated the values $HHH_{00} \dots HHG_{01}$ and so forth, as illustrated in Fig. 22, onto output leads 524. Row convolver 502 receives the values $HHH_{00} \dots HHG_{01}$ and so forth as illustrated in Fig. 24 and generates the values $HH_{00}, HH_{01}, HH_{02}$ and so forth as illustrated in Fig. 24 onto output leads 520 of row convolver 502. The data flow of Fig. 25 indicates that the row convolver performs the start reconstruction filter on the first two data values of a row, performs the odd and even reconstruction filters on subsequent non-boundary data values, and performs the end reconstruction filter on the last two data values of a row. The HH data values output from row convolver 502 are written to memory unit 116 into the memory locations corresponding with the HH data values shown in Fig. 16.

To inverse transform the octave 0 data values in memory unit 116 into the image domain, the column convolver 504 and the row convolver 502 perform an inverse octave 0 discrete wavelet transform. Figures 26 and 27 show the control signals and the data flow for the column convolver 504 of Figure 5 when the column convolver performs an inverse octave 0 discrete wavelet transform on transformed image data values in memory unit 116. The data values output from the column convolver are then supplied to the row convolver 502 of Figure 5 via lines 528 and multiplexer mux1 510.

Figures 28 and 29 show control signals and data flow for the row convolver 502 of Figure 5 when the row convolver performs an inverse octave 0 discrete wavelet transform on the data output from the column convolver to inverse transform the transformed-image data back to the image domain. Column convolver 504 receives transformed values $HH_{00} \dots GH_{03}$ and so forth as illustrated in Fig. 27 and generates the values $H_{00} \dots G_{03}$ and so forth as illustrated in Fig. 26 onto output leads 524. Row convolver 502 receives the values $H_{00} \dots G_{03}$ and so forth, as illustrated in Fig. 28, and generates the inverse transformed data values $D_{00}, D_{01}, D_{02} \dots D_{07}$ and so forth, as illustrated in Fig. 28, onto output leads 520 of row convolver 502. The inverse transformed data values output from row convolver 502 are written to memory unit 114.

The control signals and the data flows of Figures 22, 23, 24, 25, 26, 27, 28 and 29 comprise the inverse transformation from octave 1 to octave 0 and from octave 0 back into image domain. Inverse transformed data values which are substantially the same as the original data values of the matrix Table 1. The control signals which control the row convolver and column convolver to perform the inverse transform are generated by control block 506. The addresses and control signals used to read data values from and write data values to memory units 116 and 114 are generated by the DWT address generator block 508 under the control of control block 506.

After the inverse wavelet transform of the Y matrix of transformed data values is completed, the U and V matrices of transformed data values are inverse transformed one after the other in a similar way to the way the Y matrix was inverse transformed.

Figure 30 is a block diagram of the DWT address generator block 508 of Figure 5. The DWT address generator block 508 supplies read and/or write addresses to the memory units 116 and 114 for each octave of the forward and inverse transform. The DWT address generator block 508 comprises a read address generator portion and a write address generator portion. The read address generator portion comprises multiplexer 3006, adder 3010, multiplexer 3002, and resettable delay element 3014. The write address generator portion likewise comprises multiplexer 3008, adder 3012, multiplexer 3004, and resettable delay element 3016. The DWT address generator is coupled to the control block 506 via control leads 534, 556, and 544, to memory unit 116 via address leads 3022, and to memory unit 114 via address leads 3020. The input leads of DWT address generator 508 comprise the DWT address generator read control leads 534, the DWT address generator write control leads 544, and the muxcontrol lead 534. The DWT address generator read control leads 534, in turn, comprise 6 leads which carry the values col_end_R , $channel_start_R$, $reset_R$, $oct_add_factor_R$, $incr_R$, $base_u_R$, and $base_v_R$. The DWT address generator write control leads 544, in turn, comprise leads which carry the values col_end_W , $channel_start_W$, $reset_W$, $oct_add_factor_W$, $incr_W$, $base_u_W$, and $base_v_W$. All signals contained on these leads are provided by control block 506. The output leads of DWT address generator block 508 comprise address leads 3022 which provide address information to memory unit 116, and address leads 3020 which provide address information to memory unit 114. The addresses provided on leads 3022 can be either read or write addresses, depending on the cycle of the DWT transform circuit 122 as dictated by control signal muxcontrol provided by control block 506 on lead 556. The addresses provided on leads 3020 are write-only addresses, because memory unit 114 is only written to by the DWT transform circuit 122.

Memory locations of a two-dimensional matrix of data values such as the matrices of Table 1, Figure 12 and Figure 16 may have memory location addresses designated 0, 1, 2 and so forth, the addresses increasing by one left to right across each row and increasing by one to skip from the right most memory location at the end of a row to the left most memory location of the next lower row. To address successive data values in a matrix of octave 0 data values, the address is incremented by one to read each new data value D from the

matrix.

For octave 1, addresses are incremented by two because the HH values are two columns apart as illustrated in Figure 16. The row number, however, is incremented by two rather than one because the HH values are located on every other row. The DWT address generator 508 in octave 1 therefore increments by two until the end of a row is reached. The DWT address generator then increments once by $ximage + 2$ as can be seen from Figure 16. For example, the last HH value in row 0 of Figure 16 is HH_{03} at memory address 6 assuming HH_{00} has an address of 0 and that addresses increment by one from left to right, row by row, through the data values of the matrix. The next HH value is in row two, HH_{10} , at memory address 16. The increment factor in a row is therefore $incr = 2^{octave}$. The increment factor at the end of a row is $oct_add_factor = (2^{octave} - 1) \cdot ximage + 2^{octave}$ for octave ≥ 0 , where $ximage$ is the x dimension of the image.

In some embodiments, the transformed Y data values are stored in memory unit 116 from addresses 0 through $(ximage \cdot yimage - 1)$, where $yimage$ is the y dimension of the matrix of the Y data values. The transformed U data values are then stored in memory unit 116 from address $base_u$ up to $base_v - 1$, where:

$$base_u = ximage \cdot yimage$$

$$base_v = ximage \cdot yimage + \frac{ximage \cdot yimage}{4}$$

Similarly, the transformed V data values are stored in memory unit 116 at addresses beginning at address $base_v$.

The operation of the read address generator portion in Figure 30 is representative of both the read and write portions. In operation, multiplexer $base_mux$ 3002 of Figure 30 sets the read base addresses to be 0 for the Y channel, $base_u_R$ for the U channel, and $base_v_R$ for the V channel. Multiplexer 3002 is controlled by the control signals $channel_start_R$ which signifies when each Y, U, V channel starts. Multiplexer mux 3006 sets the increment factor to be $incr_R$, or, at the end of each row, to $oct_add_factor_R$. The opposite increment factor is supplied to adder 3010 which adds the increment factor to the current address present on the output leads of delay elements 3014 so as to generate the next read address, $next_addr_R$. The next read address $next_addr_R$ is then stored in the delay element 3014.

In some embodiments in accordance with the present invention, tables of $incr_R$ and $oct_add_factor_R$ for each octave are downloaded to REGISTERS block 536 on the video encoder/decoder chip 112 at initialization via download registers bus 128. These tables are passed to the control block 506 at initialization. To clarify the illustration, the leads which connect REGISTERS block 536 to control block 506 are not included in Figure 5. In other embodiments, values of $incr_R$ and $oct_add_factor_R$ are precalculated in hardware from the value of $ximage$ using a small number of gates located on-chip. Because the U and V matrices have half the number of columns as the Y matrix, the U and V jump tables are computed with $ximage$ replaced by $\frac{ximage}{2}$, a one bit shift. Because the tree encoder/decoder restricts $ximage$ to be a multiple of $2^{(OCT + 1)} > 2^{octave}$,

the addition of 2^{octave} in the oct_add_factor is, in fact, concatenation. Accordingly, only the factor $(2^{octave} - 1) \cdot ximage$ must be calculated and downloaded. The jump tables for the U and V addresses can be obtained from the Y addresses by shifting this factor one bit to the right and then concatenating with 2^{octave} . Accordingly, appropriate data values of a matrix can be read from a memory storing the matrix and processed data values can be written back into the matrix in the memory to the appropriate memory locations.

Figures 31 and 32 are block diagrams of one embodiment of the tree processor/encoder-decoder circuit 124 of Figure 1. Figure 31 illustrates the circuit in encoder mode and Figure 32 illustrates the circuit in decoder mode. Tree processor/encoder-decoder circuit 124 comprises the following blocks: DECIDE block 3112, TP_ADDR_GEN block 3114, quantizer block 3116, MODE_CONTROL block 3118, Huffman encoder-decoder block 3120, buffer block 3122, CONTROL_COUNTER block 3124, delay element 3126, delay element 3128, and VALUE_REGISTERS block 3130.

The tree processor/encoder-decoder circuit 124 is coupled to FIFO buffer 120 via input/output data leads 130. The tree processor/encoder-decoder circuit 124 is coupled to memory unit 116 via an old frame data bus 3102, a new frame data bus 3104, an address bus 3106, and memory control buses 3108 and 3110. The VALUE_REGISTERS block 3130 of the tree processor/encoder-decoder circuit 124 is coupled to data bus 106 via a register download bus 128. Figures 31 and 32 illustrate the same physical hardware; the encoder and decoder configurations of the hardware are shown separately for clarity. Although two data buses 3104 and 3102 are illustrated separately in Figure 31 to facilitate understanding, the new and old frame data buses may actually share the same pins on video encoder/decoder chip 112 so that the new and old frame data are time multiplexed on the same leads 526 of memory unit 116 as illustrated in Figure 5. Control buses 3108 and 3110 of Figure 31 correspond with the control lines 2108 in Figure 5. The DWT address generator block 508 of the discrete wavelet transform circuit 122 and the tree processor address generator block 3114 of the tree processor/encoder-decoder circuit 124 access memory unit 116 therefore may use the same physical address, data and

control lines.

Figure 33 illustrates an embodiment of DECIDE block 3112. A function of DECIDE block 3112 is to receive a two-by-two block of data values from memory unit 116 for each of the old and new frames and from these two-by-two blocks of data values and from the signals on leads 3316, 3318, 3320 and 3322, to generate seven flags present on leads 3302, 3304, 3306, 3308, 3310, 3312 and 3314. The MODE_CONTROL block 3118 uses these flags as well as values from VALUE_REGISTERS block 3130 supplied via leads 3316, 3318 and 3320 to determine the mode in which the new two-by-two block will be encoded. The addresses in memory unit 116 at which the data values of the new and old two-by-two blocks are located and determined by the address generator TP_ADDR_GEN block 3114.

The input signal on register lead 3316 is the limit value output from VALUE_REGISTERS block 3130. The input signal on register leads 3318 is the qstep value output from VALUE_REGISTERS block 3130. The input signal on register lead 3320 is the compare value output from VALUE_REGISTERS block 3130. The input signal on register lead 3322 is the octave value generated by TP_ADDR_GEN block 3114 as a function of the current location in the tree of the sub-band decomposition. As described in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" at equations 62-71, the values of the flags new_z, nz_flag, origin, noflag, no_z, oz_flag, and motion, produced on leads 3302, 3304, 3306, 3308, 3310, 3312, and 3314, respectively, are determined in accordance with the following equations:

$$nz = \sum_{0 \leq x, y \leq 1} |new(x)[y]| \quad (\text{equ. 1})$$

$$oz = \sum_{0 \leq x, y \leq 1} |old(x)[y]| \quad (\text{equ. 2})$$

$$no = \sum_{0 \leq x, y \leq 1} |new(x)[y] - old(x)[y]| \quad (\text{equ. 3})$$

$$nz_flag = nz < limit \quad (\text{equ. 4})$$

$$noflag = no < compare \quad (\text{equ. 5})$$

$$origin = nz \leq no \quad (\text{equ. 6})$$

$$motion = ((nz + oz) < octave) \leq no \quad (\text{equ. 7})$$

$$new_z = |new(x)[y]| < qstep,$$

$$\text{for } 0 \leq x, y \leq 1 \quad (\text{equ. 8})$$

$$no_z = |new(x)[y] - old(x)[y]| < qstep,$$

$$\text{for } 0 \leq x, y \leq 1 \quad (\text{equ. 9})$$

$$oz_flag = old(x)[y] = 0,$$

$$\text{for all } 0 \leq x, y, \leq 1 \quad (\text{equ. 10})$$

The DECIDE block 3112 comprises subtractor block 3324, absolute value (ABS) blocks 3326, 3328, and 3330, summation blocks 3332, 3334, and 3336, comparator blocks 3338, 3340, 3342, 3344, 3346, 3350, and 3352, adder block 3354, and shift register block 3356. The value output by ABS block 3326 is the absolute value of the data value new[x][y] on leads 3104. Similarly, the value output by ABS block 3328 is the absolute value of the data value old[x][y] on leads 3102. The value output by ABS block 3330 is the absolute value of the difference between the data values new[x][y] and old[x][y]. Comparator 3338, coupled to the output leads of ABS block 3326, unasserts new_z flag on lead output 3302 if qstep is less than the value output by block 3326. Block 3332 sums the last four values output from block 3326 and the value output by block 3332 is supplied to comparator block 3340. Comparator block 3340 compares this value to the value of limit 3316. The flag nz_flag 3304 is asserted on lead 3304 if limit is greater than or equal to the value output by block 3332. This value corresponds to nz_flag in equation 4. Summation block 3334 similarly sums the four most recent values output by block 3328. The values outputs by blocks 3332 and 3334 are added together by block 3354, the values output by block 3354 being supplied to shift register block 3356. The shift register block 3356 shifts the value received to the left by octave bits. Summation block 3336 adds the four most recent values output by block 3330. Comparator block 3342 compares the value output by block 3332 to the value output by block

3336 and asserts the motion flag in accordance with equation 7. The origin flag on output lead 3306 is asserted when the value output by block 3332 is less than the value output by 3336. This value corresponds to origin in equation 8 above. The value output by block 3336 is compared to the value compare by block 3344 such that flag noflag is asserted when compare is greater than the value output from block 3336. Block 3346 compares the value output by block 3330 to the value qstep such that flag no_z is unasserted when qstep is less. This corresponds to flag no_z in equation 9. The old input value on leads 3102 is compared to the value 0 by block 3350 such that flag oz_flag on lead 3312 is asserted when each of the values of the old block is equal to 0. This corresponds to oz_flag in equation 10 above. The seven flags produced by the DECIDE block of Figure 33 are passed to the MODE_CONTROL block 3118 to determine the next mode.

The tree processor/encoder-decoder circuit 124 of Figure 31 comprises delay elements 3126 and 3128. Delay element 3126 is coupled to the NEW portion of memory unit 116 via new frame data bus 3104 to receive the value new[x][y]. Delay element 3128 is coupled to the OLD portion of memory unit 116 via old frame data bus 3102 to receive the value old[x][y]. These delay elements, which in some embodiments of the invention are implemented in static random access memory (SRAM), serve to delay their respective input values read from memory unit 116 for four cycles before the values are supplied to quantizer block 3116. This delay is needed because the DECIDE block 3112 introduces a four-cycle delay in the dataflow as a result needing to read the four most recent data values before the new mode in which those data values will be encoded is determined. The delay elements therefore synchronize signals supplied to quantizer block 3116 by the MODE_CONTROL block 3118 with the values read from memory unit 116 which are supplied to quantizer block 3116.

The tree processor/encoder-decoder circuit 124 of Figures 31 and 32 comprises a VALUE_REGISTERS block 3130. The VALUE_REGISTERS block 3130 serves the function of receiving values from an external source and asserting these values onto leads 3316, 3318, 3320, 3132, 3134 and 3136, which are coupled to other blocks in the tree processor/encoder-decoder 124. In the presently described embodiment the external source is data bus 106 and VALUE_REGISTERS block 3130 is coupled to data bus 106 via a download register bus 128. Register leads 3316 carry a signal corresponding to the value of limit and are coupled to DECIDE block 3112 and to MODE_CONTROL block 3118. Register leads 3318 carry signals indicating the value of qstep and are coupled to DECIDE block 3112 and to MODE_CONTROL block 3118. Register leads 3320 carry signals indicating the value of compare and are coupled to DECIDE block 3112 and to MODE_CONTROL block 3118. Register leads 3132 carry signals indicating the value of ximage and are coupled to TP_ADDR_GEN block 3114 and to MODE_CONTROL block 3118. Register leads 3134 carry signals indicating the value of yimage and are coupled to TP_ADDR_GEN block 3114 and to MODE_CONTROL block 3118. Register lead 3136 carries a signal corresponding to the value of direction and is coupled to TP_ADDR_GEN block 3114, MODE_CONTROL block 3118, buffer block 3122, Huffman encoder-decoder block 3120, and quantizer block 3116. To clarify the illustration, only selected ones of the connections between the VALUE_REGISTERS block 3130 and other blocks of the tree processor/encoder-decoder circuit 124 are illustrated in Figures 31 and 32. VALUE_REGISTERS block 3130 is, in some embodiments, a memory mapped register addressable from bus 106.

Figure 34 is a block diagram of an embodiment of address generator TP_ADDR_GEN block 3114 of Figure 32. The TP_ADDR_GEN block 3114 of Figure 34 generates addresses to access selected two-by-two blocks of data values in a tree of a sub-band decomposition using a counter circuit (see Figures 27-29 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" and the corresponding text). Figure 34 illustrates a three-octave counter circuit. The signals supplied to TP_ADDR_GEN block 3114 are provided by MODE_CONTROL block 3118, CONTROL_COUNTER block 3124, and VALUE_REGISTERS block 3130. MODE_CONTROL block 3118 is coupled to TP_ADDR_GEN block 3114 by leads 3402 which carry the three bit value new_mode. CONTROL_COUNTER 3124 is coupled to TP_ADDR_GEN block 3114 by leads 3404 and 3406 which carry signals read_enable and write_enable, respectively. VALUE_REGISTER block 3130 is coupled to TP_ADDR_GEN block 3114 by register leads 3132 which carry a signal indicating the value of ximage. The output leads of TP_ADDR_GEN block 3114 comprise tree processor address bus 3106 and octave leads 3322. The address generator TP_ADDR_GEN block 3114 comprises a series of separate counters: counter TreeRoot_x 3410, counter TreeRoot_y 3408, counter C3 3412, counter C2 3414, counter C1 3416, and counter sub_count 3418. TP_ADDR_GEN block 3114 also comprises CONTROL_ENABLE block 3420, multiplexer 3428, multiplexer 3430, NOR gate 3436, AND gates 3422, 3424 and 3426, AND gates 3428, 3430 and 3432, multiplier block 3432 and adder block 3434.

Counter TreeRoot_x 3410 counts from 0 up to $\frac{ximage}{2^{OCT+1}} - 1$ and counter TreeRoot_y 3408 counts from 0 up to $\frac{yimage}{2^{OCT+1}} - 1$, where OCT is the maximum number of octaves in the decomposition. Counters C3, C2, C1, and sub_count are each 2-bit counters which count from 0 up to 3, and then return to 0. Each of these counters

takes on its next value in response to a respective count enable control signal supplied by CONTROL_ENABLE block 3420. Figure 34 shows count enable control signals x_en, y_en, c3_en, c2_en, c1_en, and sub_en, being supplied to the counters TreeRoot_x, TreeRoot_y, C3, C2, C1 and sub_count, respectively. When one of the counters reaches its maximum value, the counter asserts a carry out signal back to the CONTROL_ENABLE block 3420. These carry out signals are denoted in Figure 34 as x_carry, y_carry, c3_carry, c2_carry, c1_carry, and sub_carry.

CONTROL_ENABLE block 3420 responds to input signal new_mode on leads 3402 and to the carry out signals to generate the counter enable signals. The octave signal output by CONTROL_ENABLE is the value of the octave of the transform of the data values currently being addressed. The c1_carry, c2_carry, and c3_carry signals are logically ANDed with the write_enable signal supplied from CONTROL_COUNTER block 450 before entering the CONTROL_ENABLE block 3420. This AND operation is performed by AND gates 3422, 3424, and 3426 as shown in Figure 34. The counter enable signals from CONTROL_ENABLE block 3420 are logically ANDed with the signal resulting from the logical ORing of read_enable and write_enable by OR gate 3436. These ANDing operations are performed by AND gates 3428, 3430, and 3432 as shown in Figure 34. AND gates 3422, 3424, 3426, 3428, 3430, and 3432 function to gate the enable and carry signals with the read_enable and write_enable signals such that the address space is cycled through twice per state, once for reading and once for writing.

The CONTROL_ENABLE block 3420 outputs the enable signals enabling selected counters to increment when the count value reaches 3 in the case of the 2-bit counters 3412, 3414, and 3418, or when the count value reaches $\frac{ximage}{2^{oct+1}} - 1$ in the case of TreeRoot_x 3410, or when the count value reaches $\frac{yimage}{2^{oct+1}} - 1$ in the case of TreeRoot_y 3408. The resulting x and y addresses of a two-by-two block of data values of a given octave in a matrix of data values are obtained from the signals output by the various counters as follows:

For octave = 0:

$$\begin{array}{llllll} x = \text{TreeRoot}_x & C3(2) & C2(2) & C1(2) & \text{sub_count}(2) & \text{(equ. 11)} \\ y = \text{TreeRoot}_y & C3(1) & C2(1) & C1(1) & \text{sub_count}(1) & \text{(equ. 12)} \end{array}$$

For octave = 1:

$$\begin{array}{llllll} x = \text{TreeRoot}_x & C3(2) & C2(2) & \text{sub_count}(2) & 0 & \text{(equ. 13)} \\ y = \text{TreeRoot}_y & C3(1) & C2(1) & \text{sub_count}(1) & 0 & \text{(equ. 14)} \end{array}$$

For octave = 2:

$$\begin{array}{llllll} x = \text{TreeRoot}_x & C3(2) & \text{sub_count}(2) & 0 & 0 & \text{(equ. 15)} \\ y = \text{TreeRoot}_y & C3(1) & \text{sub_count}(1) & 0 & 0 & \text{(equ. 16)} \end{array}$$

Figure 34 and equations 11-16 illustrate how the x and y address component values are generated by multiplexers 3428 and 3430, respectively, depending on the value of octave. The (2) in equations 11-16 denotes the least significant bit of a 2-bit counter whereas the (1) denotes the most significant bit of a 2-bit counter. TreeRoot_x and TreeRoot_y are the multibit values output by counters 3410 and 3408, respectively. The output of multiplexer 3430 is supplied to multiplier 3432 so that the value output by multiplexer 3430 is multiplied by the value ximage. The value output by multiplier 3432 is added to the value output by multiplexer 3428 by adder block 3434 resulting in the actual address being output onto address bus 3106 and to memory unit 116.

Appendix A discloses one possible embodiment of CONTROL_ENABLE block 3420 of a three octave address generator described in the hardware description language VHDL. An overview of the specific implementation given in this VHDL code is provided below. The CONTROL_ENABLE block 3420 illustrated in Figure 34 and disclosed in Appendix A is a state machine which allows trees of a sub-band decomposition to be ascended or descended as required by the encoding or decoding method. The CONTROL_ENABLE block 3420 generates enable signals such that the counters generate four addresses of a two-by-two block of data values at a location in a tree designated by MODE_CONTROL block 3118. Instructions from the MODE_CONTROL block 3118 are read via leads 3402 which carry the value new_mode. Each state is visited for four consecutive cycles so that the four addresses of the block are output by enabling the appropriate counter C3 3412, C2 3414 or C1 3416. Once the appropriate counter reaches a count of 3, a carry out signal is sent back to CONTROL_ENABLE block 3420 so that the next state is entered on the next cycle.

Figure 35 is a state table for the TP_ADDR_GEN block 3114 of Figure 34 when the TP_ADDR_GEN block 3114 traverses all the blocks of the tree illustrated in Figure 36. Figure 35 has rows, each of which represents the generation of four address values of a block of data values. The (0-3) designation in Figure 35 represents the four values output by a counter. The names of the states (i.e., up0, up1, down1) do not indicate movement up or down the blocks of a tree but rather correspond with state names present in the VHDL code of Appendix A. (In Appendix A, the states down1, down2 and down3 are all referred to as down1 to optimize the implementation.) The state up0 in the top row of Figure 35, for example, corresponds to addressing the values of two-by-two block located at the root of the tree of Figure 36. In the tree of Figure 36 there are three octaves. After

these four addresses of the two-by-two block at the root of the tree are generated, the tree may be ascended to octave 1 by entering the state up1.

Figure 36 illustrates a complete traversal of all the data values of one tree of a 3-octave sub-band decomposition as well as the corresponding states of the CONTROL_ENABLE block of Figure 35. One such tree exists for each of the "GH", "HG" and "GG" sub-bands of a sub-band decomposition.

First, before a tree of the sub-band decomposition is traversed, all low pass HHHHHH component values of the decomposition are addressed by setting counter sub_count to output 00. Counter C3 3412 is incremented through its four values. Counter TreeRoot_x is then incremented and counter C3 3412 is incremented through its four values again. This process is repeated until TreeRoot_x reaches its maximum value. The process is then repeated with TreeRoot_y being incremented. In this manner, all HHHHHH low pass components are accessed. Equations 15 and 16 are used to compute the addresses of the HHHHHH low pass component data values.

Next, the blocks of the "GH" subband of a tree given by TreeRoot_x and TreeRoot_y are addressed. This "GH" subband corresponds to the value sub_count = 10 (sub_count (1) = 1 and sub_count (2) = 0). The up0 state shown in Figure 35 is used to generate the four addresses of the root block of the "GH" tree in accordance with equation 15. The up1 state shown in Figure 35 is then used such that addresses corresponding to equations 13 and 14 are computed to access the desired two-by-two block of data values in octave 1. The four two-by-two blocks in octave 0 are then accessed in accordance with equations 11 and 12. With TreeRoot_x and TreeRoot_y and sub_count untouched, the states zz0, zz1, zz2 and zz3 are successively entered, four addresses being generated in each state. After each one of these four states is exited, the C2 counter 3414 is incremented by CONTROL_ENABLE block 3420 via the c2_en signal once in order to move to the next octave 0 block in that branch of the tree. After incrementing in state zz3 is completed, the left hand branch of the tree is exhausted. To move to the next two-by-two block, the C3 counter 3412 is incremented and the C2 counter 3414 is cycled through its four values to generate the four addresses of the next octave 1 block in state down1 in accordance with equations 13 and 14. In this way, the TP_ADDR_GEN block 3114 generates the appropriate addresses to traverse the tree in accordance with instructions received from MODE_CONTROL block 3118. When the traversal of the "GH" sub-band tree is completed, the traversal of the sub-band decomposition moves to the corresponding tree of the next sub-band without changing the value of TreeRoot_x and TreeRoot_y. Accordingly, a "GH" "HG" and "GG" family of trees are traversed. After all the blocks of the three sub-band trees have been traversed, the TreeRoot_x and TreeRoot_y values are changed to move to another family of sub-band trees.

To move to the next family of sub-band trees, the counter TreeRoot_x 3410 is incremented, and the C3 3412, C2 3414, C1 3416 counters are returned to 0. The process of traversing the new "GH" tree under the control of the MODE_CONTROL block 3118 proceeds as before. Similarly, the corresponding "HG" and "GG" trees are traversed. After TreeRoot_x 3410 reaches its final value, a whole row of tree families has been searched. The counter TreeRoot_y 3408 is therefore incremented to move to the next row of tree families. This process may be continued until all of the trees in the decomposition have been processed.

The low pass component HHHHHH (when sub_count = 00) does not have a tree decomposition. In accordance with the present embodiment of the present invention, all of the low pass component data values are read first as described above and are encoded before the tree encoder reads and encodes the three subbands. The address of the data values in the HHHHHH subband are obtained from the octave 3 x and y addresses with sub_count = 00. Counters C3 3412, TreeRoot_x 3410, and TreeRoot_y 3408 run through their respective values. After the low pass component data values and all of the trees of all the sub-bands for the Y data values have been encoded, the tree traversal method repeats on the U and V data values.

Although all the blocks of the tree of Figure 36 are traversed in the above example of a tree traversal, the MODE_CONTROL block 3118 may under certain conditions decide to cease processing data values of a particular branch and to move to the next branch of the tree as set forth in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". This occurs, for example, when the value new_mode output by the MODE_CONTROL block 3118 indicates the mode STOP. In this case, the state machine of CONTROL_ENABLE block 3420 will move to, depending on the current location in the tree, either the next branch, or, if the branch just completed is the last branch of the last tree, the next tree.

Figure 34 illustrates control signal inputs read_enable and write_enable being supplied to TP_ADDR_GEN block 3114. These enable signals are provided because the reading of the new/old blocks and the writing of the updated values to the old frame memory occur at different times. To avoid needing two address generators, the enable signals of the counters C3 3412, C2 3414, and C1 3416 are logically ANDed with the logical OR of the read_enable and write_enable signals. Similarly, the carryout signals of these counters are logically ANDed with the write_enable signal. During time periods when the new/old blocks are read from memory, the

read_enable signal is set high and the write_enable signal is set low. This has the effect of generating the addresses of a two-by-two block, but disabling the change of state at the end of the block count. The counters therefore return to their original values they had the start of the block count so that the same sequence of four address values will be generated when the write_enable signal is set high. This time, however, the carry out is enabled into the CONTROL_ENABLE block 3420. The next state is therefore entered at the conclusion of the block count. In this manner, the address space is cycled through twice per state, once for reading and once for writing.

Figure 37 is a block diagram of one embodiment of quantizer block 3116 of Figure 31. As shown in Figure 31, quantizer block 3116 is coupled to MODE_CONTROL block 3118, a Huffman encoder-decoder block 3120, delay block 3126, delay block 3128, and VALUE_REGISTERS block 3130. Input lead 3702 carries the signal difference from MODE_CONTROL block 3118 which determines whether a difference between the new frame and old frame is to be quantized or whether the new frame alone is to be quantized. Values new[x][y] and old[x][y] are supplied on lines 3704 and 3706, respectively, and represent values from memory unit 116 delayed by four clock cycles. Input leads 3708 and 3710 carry the values sign_inv and qindex_inv from the Huffman encoder-decoder block 3120, respectively. Register leads 3318 and 3136 carry signals corresponding to the values qstep and direction from VALUE_REGISTERS block 3130, respectively.

During encoding, quantizer block 3116 performs quantization on the values new[x][y], as dictated by the signal difference and using the values old[x][y], and generates the output values qindex onto output leads 3712, sign onto output lead 3714, and a quantized and then inverse quantized value old[x][y] onto data bus 3102. The quantized and inverse quantized value old[x][y] is written back into memory unit 116.

During decoding, quantizer block 3116 performs inverse quantization on the values old[x][y], as dictated by the signal difference, sign_inv, and qindex_inv, and generates an inverse quantized value, old[x][y], which is supplied to the old portion of memory unit 116 via bus 3102. Lead 3136 carries the value direction supplied by the VALUE_REGISTERS 3130.

The value direction controls whether the quantizer operates in the encoder mode or the decoder mode. Figure 37 illustrates that multiplexers 3716 and 3718 use the direction signal to pass signals corresponding to the appropriate mode (sign and qindex for encoder mode; sign_inv and qindex_inv for decoder mode). Multiplexer 3720 passes either the difference of the new and old data values or passes the new value depending on the value of the difference signal. Absolute value block ABS 3722 converts the value output by multiplexer 3720 to absolute value form and supplies the absolute value form value to block 3724. The output leads of multiplexer 3720 are also coupled to sign block 3726. Sign block 3726 generates a sign signal onto lead 3714 and to multiplexer 3716.

Block 3724 of the quantizer block 3116 is an human visual system (HVS) weighted quantizer having a threshold of qstep. The value on input leads 3728 denoted mag in Figure 37 is quantized via a modulo-qstep division (see Figures 30 and 31 of copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" and the corresponding text). The resulting quantized index value qindex is output onto leads 3712 to the Huffman encoder block 3120. Multiplexer 3716 receives the sign signal on leads 3714 from block 3726 and also the sign_inv signal on lead 3708. Multiplexer 3716 passes the sign value in the encoder mode and passes the sign_inv value in the decoder mode. Likewise, multiplexer 3718 has as two inputs, the qindex signal on leads 3712 and the qindex_inv signal on leads 3710. Multiplexer 3718 passes the qindex value in the encoder mode and the qindex_inv value in the decoder mode. Inverse quantizer block 3730 inverse quantizes the value output by multiplexer 3718 by the value qstep to generate the value qvalue. Block NEG 3732 reverses the sign of the value on the output lead of block 3730, denoted qvalue in Figure 37. Multiplexer 3734 chooses between the positive and negative versions of qvalue as determined by the signal output from multiplexer 3716.

In the encoder mode, if the difference signal is asserted, then output leads 3712 qindex carry the quantized magnitude of the difference between the new and old data values and the output leads 3736 of multiplexer 3734 carry the inverse quantization of this quantized magnitude of the difference between the new and old values. In the encoder mode, if the difference input is deasserted, then the output leads 3712 qindex carry only the quantized magnitude of the new data value and the value on leads 3736 is the inverse quantization of the quantized magnitude of the new data value.

Adder block 3738 adds the inverse quantized value on leads 3736 to the old[x][y] data value and supplies the result to multiplexer 3740. Accordingly, when the difference signal is asserted, the difference between the old inverse quantized value on leads 3706 and the inverse quantized value produced by inverse quantizer 3730 is determined by adding in block 3738 the opposite of the inverse quantized output of block 3730 to the old inverse quantized value. Multiplexer 3740 passes the output of adder block 3738 back into the OLD portion of memory unit 116 via bus 3102. If, on the other hand, the difference signal is not asserted, then multiplexer 3740 passes the value on leads 3736 to the OLD portion of memory unit 116 via bus 3102. Accordingly, a frame

of inverse quantized values of the most recently encoded frame is maintained in the old portion of memory unit 118 during encoding.

In accordance with one embodiment of the present invention, the value of $qstep$ is chosen so that $qstep = 2^n$, where $0 \leq n \leq 7$, so that quantizer block 3724 and inverse quantizer 3730 perform only shifts by n bits.

Block 3724 then becomes in VHDL, where \gg denotes a shift to the left, and where mag denotes the value output by block 3722:

```
CASE n is
  WHEN 0 => qindex: = mag;
  WHEN 1 => qindex: = mag >> 1;
```

```
  WHEN 7 => qindex: = mag >> 7;
END CASE;
```

Similarly, block 3730 is described in VHDL as follows:

```
CASE n is
  WHEN 0 => qvalue : = qindex;
  WHEN 1 => qvalue : = (qindex << 1) & "0";
  WHEN 2 => qvalue : = (qindex << 2) & "01";
```

```
  WHEN 7 => qvalue : = (qindex << 7) & "0111111";
```

where \ll denotes a shift to the right and where $\&$ denotes concatenation. The factor concatenated after the shift is 2^{n-1} .

The tree processor/encoder-decoder circuit 124 of Figure 31 also includes a MODE_CONTROL block 3118. In the encoder mode, MODE_CONTROL block 3118 determines mode changes as set forth in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" when trees of data values are traversed to compress the data values into a compressed data stream. In the decoder mode, MODE_CONTROL block 3118 determines mode changes as set forth in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression" when trees of data values are recreated from an incoming compressed data stream of tokens and data values.

MODE_CONTROL block 3118 receives signals from DECIDE block 3112, CONTROL_COUNTER block 3124, TP_ADDR_GEN block 3114, and VALUE_REGISTERS block 3130. MODE_CONTROL block 3118 receives the seven flag values from DECIDE block 3112. The input from CONTROL_COUNTER block 3124 is a four-bit state vector 3138 indicating the state of the CONTROL_COUNTER block 3124. Four bits are needed because the CONTROL_COUNTER block 3124 can be in one of nine states. The input from TP_ADDR_GEN block 3114 is the octave signal carried by leads 3322. The VALUE_REGISTERS block 3130 supplies the values on leads 3316, 3318, 3320, 3132, 3134, and 3136 to MODE_CONTROL block 3118. Additionally, in the decoder mode, buffer 3122 supplies token values which are not Huffman decoded onto leads 3202 and to the MODE_CONTROL block 3118 as shown in Figure 32.

MODE_CONTROL block 3118 outputs a value new_mode which is supplied to TP_ADDR_GEN block 3114 via leads 3402 as well as a token length value T_L which is supplied to buffer block 3122 via leads 3140. In the encoder mode, MODE_CONTROL block 3118 also generates and supplies tokens to buffer block 3122 via leads 3202. Leads 3202 are therefore bidirectional to carry token values from MODE_CONTROL block 3118 to buffer block 3122 in the forward mode, and to carry token values from buffer 3122 to MODE_CONTROL block 3118 in the decoder mode. The token length value T_L , on the other hand, is supplied by MODE_CONTROL block 3118 to buffer block 3122 in both the encoder and decoder modes. MODE_CONTROL block 3118 also generates the difference signal and supplies the difference signal to quantizer block 3116 via lead 3142. MODE_CONTROL block 3118 asserts the difference signal when differences between new and old values are to be quantized and deasserts the difference signal when only new values are to be quantized. Appendix B is a VHDL description of an embodiment of the MODE_CONTROL block 3118 in the VHDL language.

In the encoding process, the MODE_CONTROL block 3118 initially assumes a mode, called pro_mode , from the block immediately below the block presently being encoded in the present tree. For example, the blocks in Figure 36 corresponding to states $zz0$, ..., $zz3$ in the left-most branch inherit their respective pro_modes from the left-most octave 1 block. Similarly, the left-most octave 1 block in Figure 36 inherits its pro_mode from the root of the tree in octave 2. After the data values of the new and old blocks are read and after the DECIDE block 3112 has generated the flags for the new block as described above, the state machine of

MODE_CONTROL block 3118 determines the new_mode for the new block based on the new data values, the flags, and the pro-mode. The value of new_mode, once determined, is then stored as the current mode of the present block in a mode latch. There is one mode latch for each octave of a tree and one for the low pass data values. The mode latches form a stack pointed to by octave so that the mode latches contain the mode in which each of the blocks of the tree was encoded.

The tree processor circuit of Figures 31 and 32 also comprises a Huffman encoder-decoder block 3120. In the encoder mode, inputs to the Huffman encoder-decoder block 3120 are supplied by quantizer block 3116. These inputs comprise the qindex value and the sign signal and are carried by leads 3712 and 3714, respectively. The outputs of Huffman encoder-decoder 3120 comprise the Huffman encoded value on leads 3142 and the Huffman length H_L on leads 3144, both of which are supplied to buffer block 3122.

In the decoder mode, the input to the Huffman encoder-decoder block 3120 is the Huffman encoded value carried by leads 3204 from buffer block 3122. The outputs of the Huffman encoder-decoder 3120 comprise the Huffman length H_L on leads 3144 and the sign_inv and qindex_inv values supplied to quantizer block 3116 via leads 3708 and 3710, respectively.

The Huffman encoder-decoder block 3120 implements the Huffman table shown in Table 2 using combinatorial logic.

qindex	Huffman code
-38 . . . -512	1 1 0 0 0 0 0 0 1 1 1 1 1 1 1 1
-22 . . -37	1 1 0 0 0 0 0 0 1 1 1 1 (qindex -22)
-7 . . -21	1 1 0 0 0 0 0 0 (qindex -7)
-6	1 1 0 0 0 0 0 1
.	.
.	.
.	.
-2	1 1 0 1
-1	1 1 1
0	0
1	1 0 1
2	1 0 0 1
.	.
.	.
.	.
6	1 0 0 0 0 0 0 1
7 . . 21	1 0 0 0 0 0 0 0 (qindex -7)
22 . . 37	1 0 0 0 0 0 0 0 1 1 1 1 (qindex -22)
38 . . 511	1 0 0 0 0 0 0 0 1 1 1 1 1 1 1 1

Table 2

In the encoder mode, qindex values are converted into corresponding Huffman codes for incorporation into the compressed data stream. Tokens generated by the MODE_CONTROL block 3118, on the other hand, are not encoded but rather are written directly into the compressed data stream.

Figure 38 illustrates one possible embodiment of buffer block 3122 of Figures 31 and 32. The function of buffer block 3122 in the encoder mode is to assemble encoded data values and tokens into a single serial compressed data stream. In the decoder mode, the function of buffer block 3122 is to deassemble a compressed

serial data stream into encoded data values and tokens. Complexity is introduced into buffer block 3122 due to the different lengths of different Huffman encoded data values. As illustrated in Figure 31, buffer 3122 is coupled to FIFO buffer 120 via input-output leads 130, to MODE_CONTROL block 3118 via token value leads 3202 and token length leads 3140, to Huffman encoder-decoder 3120 via leads 3144 and Huffman length leads 3144, to CONTROL_COUNTER 3124 via cycle select leads 3802, and to VALUE_REGISTERS 3130 via leads 3136.

The direction signal carried on leads 3136 from VALUE_REGISTERS block 3130 determines whether the buffer block 3122 operates in the encoder mode or in the decoder mode. In encoder mode, multiplexers 3804, 3806, 3808 and 3814 select the values corresponding to their "E" inputs in Figure 38. In the encoder mode, the buffer block 3122 processes the Huffman encoded value signal present on leads 3142, the token value signal present on leads 3202, the cycle select signal on leads 3802, the Huffman length signal H_L on leads 3144, and the token length signal T_L on leads 3140. The cycle select signal, supplied by CONTROL_COUNTER block 3124 via leads 3802, is supplied to multiplexers 3810 and 3812 to control whether a Huffman encoded value (received from Huffman encoder-decoder block 3120) or whether a non-encoded token value (received from MODE_CONTROL block 3118) is the value presently being assembled into the output data stream.

Figure 39 illustrates a simplified diagram of the buffer block 3122 of Figure 38 when configured in encoder mode. The value s_r is a running modulo sixteen sum of the input token length values and Huffman value length values. The circuit which determines s_r comprises adder block 3902, modulo sixteen divider block 3904, and delay block 3906. When the incoming length value added to the prior value s_r produces a length result of sixteen or greater, block 3904 subtracts sixteen from this length result to determine the new value of s_r . Comparator block 3908 also sends a signal high_low to input lead 3916 of multiplexer 3901 indicating that s_r has exceeded sixteen. Figure 39 shows a barrel shifter 3912 receiving data input values from the output data leads of multiplexer 3901 and from the output data leads of multiplexer 3810. Barrel shifter 3912 sends a 32-bit output signal to a 32-bit buffer 3914. The lower 16-bit output of 32-bit buffer 3914 constitutes the encoded bit stream output of the video encoder/decoder chip which is output onto input/output leads 130.

When the prior value of s_r plus the incoming value length is sixteen or greater, then the lower sixteen bits of buffer 3914 are sent out to FIFO buffer 120 and multiplexer 3901 is set to pass the upper sixteen bits of buffer 3914 back into the lower sixteen bit positions in barrel shifter 3912. The value s_r is then decremented by sixteen. These passed back bits will next become some of the bits in the lower sixteen bits of buffer 3914, on which a subsequent incoming encoded value or token received from multiplier 3810 will be stacked by the barrel shifter starting at location s_r to make sixteen or more packed bits.

Alternatively, if the value of s_r plus the length of the new incoming value is less than sixteen, then multiplexer 3901 is controlled to pass the lower sixteen bits of buffer 3914 back to barrel shifter 3912 and no bits are applied to FIFO buffer 120. The bits of a subsequent incoming encoded value or taken from multiplexer 3810 will be stacked on top of the bits of prior encoded data values or tokens in barrel shifter 3912. Because the value s_r did not exceed sixteen, s_r is not decremented by sixteen.

Figure 40 illustrates a typical output of the barrel shifter 3912 of the buffer 3122 in encoder mode. The maximum length of a Huffman encoded word is sixteen bits. All tokens are two bits in length, where length is the number of bits in the new encoded value or token. The value s in Figure 40 indicates the bit position in the barrel shifter 3912 immediately following the last encoded data value or token present in the barrel shifter. Accordingly, a new encoded value or token is written into barrel shifter 3912 at positions $s \dots s + \text{length}$. The resulting 32-bit output of the barrel shifter is rewritten to the 32-bit buffer 3914. The comparator block compares the new value of $s + \text{length}$ to sixteen. If this value $s + \text{length}$ is sixteen or greater as illustrated in Figure 40, then the control signal high_low on multiplexer input lead 3916 is asserted. The lower sixteen bits of the buffer are therefore already completely packed with either bits of data values and/or with bits of tokens. These lower sixteen bits are therefore output to comprise part of the output data stream. The upper sixteen bits, which are incompletely packed with data values and/or tokens, are sent back to the lower sixteen bit positions in the barrel shifter so that the remaining unpacked bits in the lower sixteen bits can be packed with new data bits or new token bits.

If, on the other hand, this value $s + \text{length}$ is fifteen or less, then there remain unpacked bits in the lower sixteen bit positions in barrel shifter 3912. These lower bits in barrel shifter 3912 can therefore not yet be output via buffer 3914 onto lines 130. Only when $s + \text{length}$ is sixteen or greater will the contents of barrel shifter 3912 be written to buffer 3914 so that the lower sixteen bits will be output via leads 130.

In the decoder mode, buffer 3122 receives an encoded data stream on leads 130, the token length signal T_L on leads 3140 from MODE_CONTROL block 3118, the Huffman encoded length signal H_L on leads 3144, and the control signal cycle select on lead 3802. Multiplexers 3804, 3806, and 3808 are controlled to select values on their respective "D" inputs. Cycle select signal 3802 selects between the Huffman encoded length H_L and the token length T_L depending on whether a data value or a token is being extracted from the in-

coming data stream.

Figure 41 illustrates a simplified diagram of the buffer block 3122 configured in the decoder mode. The value s_i is a running modulo thirty-two sum of the input token length values and Huffman value length values. The circuit which determines the value of s_i comprises adder block 4002, modulo thirty-two divider block 4004, and delay block 4006. When the incoming length value added to s_i results in a value greater than thirty-two, modulo thirty-two divider block 4004 subtracts thirty-two from this value. A comparator block 4008 sends a signal to buffer 3914 indicating when s_i has reached a value greater than or equal to thirty-two. Additionally, comparator block 4008 sends a signal to both buffer 3914 and to multiplexers 3901 and 4010 indicating when s_i has reached a value greater than or equal to sixteen.

Buffer 3122 in the decoder mode also comprises buffer 3914, multiplexers 3901 and 4010, and barrel shifter 4012. In the case of a Huffman encoded data value being the next value in the incoming data stream, sixteen bits of the encoded data stream that are present in barrel shifter 4012 are passed via output leads 3204 to the Huffman decoder block 3120. The number of bits in the sixteen bits that represent an encoded data value depends on the data value itself in accordance with the Huffman code used. In the case of a token being the next value in the incoming data stream, only the two most significant bits from barrel shifter 4012 are used as the token value which is output onto leads 3202 to MODE_CONTROL block 3118. The remaining fourteen bits are not output during this cycle. After a number of bits of either an encoded data value or a two-bit token is output, the value of s_i is updated to point directly to the first bit of the bits in barrel shifter 4012 which follows the bit last output. The circuit comprising adder block 4002, module block 4004, and delay element 4006 adds the length of the previously output value or token to s_i modulo thirty-two to determine the starting location of the next value or token in barrel shifter 4012. Comparator block 4008 evaluates the value of s_i plus the incoming length value, and transmits an active value on lead 4014 when this value is greater than or equal to sixteen and also transmits an active value on lead 4016 if this value is greater than or equal to thirty-two. When s_i is greater or equal to sixteen, the buffer 3914 will read in a new sixteen bits of encoded bit stream bits into its lower half. When $s_i \geq 32$, the buffer 3914 will read a new sixteen bits into its upper half. The two multiplexers 4010 and 3910 following the buffer 3914 rearrange the order of the low and high halves of the buffer 3914 to maintain at the input leads of barrel shifter 4012 the original order of the encoded data stream.

The tree processor/encoder-decoder circuit 124 of Figures 31 and 32 comprises a CONTROL_COUNTER block 3124. CONTROL_COUNTER block 3124 controls overall timing and sequencing of the other blocks of the tree processor/encoder/decoder circuit 124 by outputting the control signals that determine the timing of the operations that these blocks perform. In accordance with one embodiment of the present invention, the tree processor/encoder/decoder 112 is fully pipelined in a nine stage pipeline sequence, each stage occupying one clock cycle. Appendix C illustrates an embodiment of CONTROL_COUNTER block 3124 described in VHDL code.

The signals output by CONTROL_COUNTER block 3124 comprise a read_enable signal on lead 3404, which is active during read cycles, and a write_enable signal on lead 3406, which is active during write cycles. The signals output also comprise memory control signals on leads 3108 and 3110, which control the old and new portions of memory unit 116, respectively, for reading from memory or writing to memory. The signals output also comprise a 4-bit state vector on lead 3138, which supplies MODE_CONTROL block 3118 with the current cycle. The four-bit state vector counts through values 1 through 4 during the "skip" cycle, the value 5 during the "token" cycle, and the values 6-9 during the "data" cycle. The signals output by CONTROL_COUNTER block 3124 also comprise a cycle state value on leads 3802, which signals buffer 3122 when a token cycle or data cycle is taking place.

Figure 42 illustrates a pipelined encoding/decoding process controlled by CONTROL_COUNTER block 3124. Cycles are divided into three types: data cycles - when Huffman encoded/decoded data is being output/input into the encoded bit stream and when old frame values are being written back to memory; token cycles - when a token is being output/input; and skip cycles - the remaining case when no encoded/decoded data is output to or received from the encoded bit stream. A counter in CONTROL_COUNTER block 3124 counts up to 8 then resets to 0. At each sequence of the count, this counter decodes various control signals depending on the current MODE. The pipeline cycles are:

- 0) read old[0][0] and in encode new[0][0]; skip cycle.
- 1) read old[1][0] and in encode new[1][0]; skip cycle.
- 2) read old[0][1] and in encode new[0][1]; skip cycle.
- 3) read old[1][1] and in encode new[1][1]; skip cycle.
- 4) DECIDE blocks outputs flags MODE_CONTROL write/read token into/from coded data stream; generates new_mode, outputs tokens in encode; generates new_mode, inputs tokens in decode; token cycle.

- 5) Huffman encode/decode $qindex[0][0]$, and write $old[0][0]$; data cycle.
- 6) Huffman encode/decode $qindex[1][0]$, and write $old[1][0]$; data cycle.
- 7) Huffman encode/decode $qindex[0][1]$, and write $old[0][1]$; data cycle.
- 8) Huffman encode/decode $qindex[1][1]$, and write $old[1][1]$; data cycle.

Figure 42 illustrates that once the new_mode is calculated, another block of data values in the tree can be processed. The tree processor/encoder/decoder is thus fully pipelined, and can process four new transformed data values every five clock cycles. To change the pipeline sequence, it is only required that the control signals in the block CONTROL_COUNTER block 3124 be reprogrammed.

ADDITIONAL EMBODIMENTS

In accordance with the above-described embodiments, digital video in 4:1:1 format is output from A/D video decoder 110 on lines 202 to the discrete wavelet transform circuit 122 of video encoder/decoder circuit 112 row by row in raster-scan form. Figure 43 illustrates another embodiment in accordance with the present invention. Analog video is supplied from video source 104 to an A/D video decoder circuit 4300. The A/D video decoder circuit 4300, which may, for example, be manufactured by Philips, outputs digital video in 4:2:2 format on lines 4301 to a horizontal decimeter circuit 4302. For each two data values input to the horizontal decimeter circuit 4302, the horizontal decimeter circuit 4302 performs low pass filtering and outputs one data value. The decimated and low pass filtered output of horizontal decimeter circuit 4302 is supplied to a memory unit 114 such that data values are written into and stored in memory unit 114 as illustrated in Figure 43. The digital video in 4:2:2 format on lines 4301 occurs at a frame rate of 30 frames per second, each frame consisting of two fields. By discarding the odd field, the full 33.3 ms frame period is available for transforming and compressing/decompressing the remaining even field. The even fields are low-pass filtered by the horizontal decimeter circuit 4302 such that the output of horizontal decimeter circuit 4302 occurs at a rate of 30 frames per second, each frame consisting of only one field. Memory unit 114 contains 640 x 240 total image data values. There are 320 x 240 Y data values, as well as 160 x 240 U data values, as well as 160 x 240 V data values.

In order to perform a forward transform, the Y values from memory unit 114 are read by video encoder/decoder chip 112 as described above and are processed by the row convolver and column convolver of the discrete wavelet transform circuit 122 such that a three octave sub-band decomposition of Y values is written into memory unit 116. The three octave sub-band decomposition for the Y values is illustrated in Figure 43 as being written into a Y portion 4303 of the new portion of memory unit 116.

After the three octave sub-band decomposition for the Y values has been written into memory unit 116, the video encoder/decoder chip 112 reads the U image data values from memory unit 114 but bypasses the row convolver. Accordingly, individual columns of U values in memory unit 114 are digitally filtered into low and high pass components by the column convolver. The high pass component G is discarded and the low pass component H is written into U portion 4304 of the new portion of memory unit 116 illustrated in Figure 43. After the U portion 4304 of memory unit 116 has been written with the low pass H component of the U values, video encoder/decoder chip 112 reads these U values from U portion 4304 and processes these U data values using both the row convolver and column convolver of the discrete wavelet transform circuit 122 to perform an additional two octaves of transform to generate a U value sub-band decomposition. The U value sub-band decomposition is stored in U portion 4304 of memory unit 116. Similarly, the V image data values in memory unit 114 are read by video encoder/decoder chip 112 into the column convolver of the discrete wavelet transform circuit 122, the high pass component G being discarded and the low pass component H being written into V portion 4305 of the new portion of memory unit 116. The V data values of V portion 4305 are then read by the video encoder/decoder chip 112 and processed by both the row convolver and the column convolver of discrete wavelet transform circuit 122 to generate a V sub-band decomposition corresponding to the U sub-band decomposition stored in U portion 4304. This process completes a forward three octave discrete wavelet transform comparable to the 4:1:1 three octave discrete wavelet transform described above in connection with Figures 3A-3C. Y portion 4303 of memory unit 116 comprises 320 x 240 data value memory locations; U portion 4304 comprises 160 x 120 data value memory locations; and V portion 4305 comprises 160 x 120 data value memory locations.

The DWT address generator 508 illustrated in Figure 5 generates a sequence of 19-bit addresses on output lines OUT2. In accordance with the presently described embodiment, however, memory unit 114 is a dynamic random access memory (DRAM). This memory unit 114 is loaded from horizontal decimeter circuit 4302 and is either read from and written to by the video encoder/decoder chip 112. For example, in order for the video encoder/decoder chip 112 to access the Y data values in memory unit 114 the Inc_R value supplied to DWT address generator 508 by control block 506 is set to 2. This causes the DWT address generator 508 of the video encoder/decoder chip 112 to increment through even addresses as illustrated in Figure 43 such that only

the Y values in memory unit 114 are read. After all the Y values are read from memory unit 114 and are transformed into a Y sub-band decomposition, then base_u_R is changed to 1 and the Channel_start_r is set so that BASE_MUX 3002 of Figure 30 selects the base_u_R to address the first U data value in memory unit 114. Subsequent U data values are accessed because the inc_R value is set to 4 such that only U data values in memory unit 114 are accessed. Similarly, the V data values are accessed by setting the base_v_R value to 3 and setting the Channel_start_r value such that BASE_MUX 3002 selects the base_v_R Input leads. Successive V data values are read from memory unit 114 because the inc_R remains at 3.

Because in accordance with this embodiment the video encoder/decoder chip 112 reads memory unit 114, the DWT address generator 508 supplies both read addresses and write addresses to memory unit 114. The read address bus 3018 and the write address bus 3020 of Figure 30 are therefore multiplexed together (not shown) to supply the addresses on the OUT2 output lines of the DWT address generator.

To perform the inverse transform on a three octave sub-band decomposition stored in memory unit 116 of Figure 43, the row and column convolvers of the video encoder/decoder chip 112 require both low and high pass components to perform the inverse transform. When performing the octave 0 inverse transform on the U and V data values of the sub-band decomposition, zeros are inserted when the video encoder/decoder chip 112 is to read high pass transformed data values. In the octave 0 inverse transform, the row convolver is bypassed such that the output of the column convolver is written directly to the appropriate locations in the memory unit 114 for the U and V inverse transform data values. When the Y transform data values in memory unit 116 are to be inverse transformed, on the other hand, both the column convolver and the row convolver of the video encoder/decoder chip 112 are used on each of the three octaves of the inverse transform. The resulting inverse transformed Y data values are written into memory unit 114 in the appropriate locations as indicated in Figure 43.

Figure 44 illustrates a sequence of reading and writing Y data values from the Y portion of the new portion of memory unit 116 in accordance with the embodiment of the present invention illustrated in Figure 1 where memory unit 116 is a static random access memory (SRAM). The dots in Figure 44 represent individual memory locations in a two-dimensional matrix of memory locations adequately wide and deep to store an entire sub-band decomposition of the Y values in a single two-dimensional matrix. The discrete wavelet transform chip 122 reads the memory location indicated R0 during a first time period, outputs a transformed data value during a second time period to the memory location indicated W1, reads another data value from the memory location denoted R2, writes a transformed data value to the memory location denoted W3 and so forth. If memory unit 116 is realized as a dynamic random access memory (DRAM), addressing memory unit 116 in this manner results in a different row of the memory unit being accessed each successive time period. When successive accesses are made to different rows of standard dynamic random access memory, a row address select (RAS) cycle must be performed each time the row address changes. On the other hand, if successive accesses are performed on memory locations that fall in the same row, then only column address select (CAS) cycles need to be performed. Performing a CAS cycle is significantly faster in a standard dynamic random access memory than a RAS cycle. Accordingly, when memory unit 114 is realized as a dynamic random access memory and when memory unit 116 is read and written in the fashion illustrated in Figure 44, memory accesses are slow.

Figure 45 illustrates a sequence of reading and writing memory unit 116 in accordance with another embodiment of the present invention wherein memory unit 116 is realized as a dynamic random access memory. Again, the dots denote individual memory locations and the matrix of memory locations is assumed to be wide enough and deep enough to accommodate the Y portion of the sub-band decomposition in a single two-dimensional matrix. In the first time period, the memory location designated R0 is read. In the next time period, the memory location R1 is read, then R2 is read in a subsequent time period, then R3 is read in a subsequent period, and so forth. In this way one row of low pass component HH values is read into the video encoder/decoder chip 112 using only one RAS cycle and multiple CAS cycles. Then, a second row of low pass component HH data values is read as designated in Figure 45 by numerals R160, R161, R162 and so forth. The last low pass component data value to be read in the second row is designated R319. This row is also read into the video encoder/decoder chip 112 using only one RAS cycle and multiple CAS cycles. Figure 15 illustrates that after reading the data values that the resulting octave 1 transformed data values determined by the discrete wavelet transform chip 122 are now present in the line delays designated 1334 and 1340 illustrated in Figure 13. At this point in this embodiment of the present invention, the row convolver and the column convolver of the discrete wavelet transform chip 122 are stopped by freezing all the control signals except that line delays 1334 and 1340 are read in sequential fashion and written to the Y portion of the new portion of memory unit 116 as illustrated in Figure 45. In this fashion, two rows of memory locations which were previously read in time periods 0 through 319 are now overwritten with the resulting octave 1 transformed values in periods 320 through 639. Only one RAS cycle is required to write the transformed data values in time periods 320 through 479. Similarly, only one RAS cycle is required to write transformed data values during time periods 480 through

639. This results in significantly faster accessing of memory unit 116. Because dynamic random access memory can be used to realize memory unit 116 rather than static random access memory, system cost is reduced considerably.

In accordance with this embodiment of the present invention, the output of the output OUT2 of the column convolver of the video encoder/decoder circuit 112 is coupled to the output leads of block 1332 as illustrated in Figure 13. However, in the forward or inverse transform of any other octave, the output leads OUT2 are coupled to the line delay 1340. Accordingly, in an embodiment in accordance with the memory accessing scheme illustrated in Figure 45, a multiplexer (not shown) is provided to couple either the output of line delay 1340 or the output of adder block 1332 to the output leads OUT2 of the column wavelet transform circuit 704 of Figure 13.

Figure 46 illustrates another embodiment in accordance with the present invention. Memory unit 116 contains a new portion and an old portion. Each of the new and old portions contains a sub-band decomposition. Due to the spatial locality of the wavelet sub-band decomposition, each two-by-two block of low pass component data values has a high pass component consisting of three trees of high frequency two-by-two blocks of data values. For example, in a three octave sub-band decomposition, each two-by-two block of low pass component data values and its associated three trees of high pass component data values forms a 16-by-16 area of memory which is illustrated in Figure 46.

In order for memory unit 116 to be realized in dynamic random access memory (DRAM), the static random access memories (SRAMs) 4600, 4601, 4602 and 4603 which are used as line delays in the discrete wavelet transform circuit 122 are used as cache memory to hold one 16-by-16 block in the new portion of memory unit 116 as well as one 16-by-16 block in the old portion of memory unit 116. This allows each 16-by-16 block of dynamic random access memory realizing the new and old portions of memory unit 116 to be accessed using at most sixteen RAS cycles. This allows the video encoder/decoder chip 112 to use dynamic random access memory for memory unit 116 rather than static random access memory, thereby reducing system cost.

Figure 47 illustrates a time line of a sequence of operations performed by the circuit illustrated in Figure 46. In a first time period, old 16-by-16 block 3 is read into SRAM 1 4601. Because there is only one set of data pins on video encoder/decoder chip 112 for accessing memory unit 116, the 16-by-16 block 0 of the new portion of memory unit 116 is read into SRAM 0 4600 in the second time period. Bidirectional multiplexer 4604 is controlled by select inputs 4605 to couple the 16-by-16 block of old data values now present in SRAM 1 4601 to the bidirectional input port old 4606 of the tree processor/ encoder/decoder circuit 124. Similarly, the 16-by-16 new data values present in SRAM 0 4600 are coupled to the input port new 4607 of the tree processor/encoder/decoder circuit 124. Accordingly, the tree processor/ encoder/decoder circuit 124 performs tree processing and encoding in a third time period. During the same third time period, the inverse quantized old 16-by-16 block is rewritten into SRAM 1 4601 through multiplexer 4604. In a fourth time period, old 16-by-16 block 2 is read into SRAM 2 4602. Subsequently, in the fifth time period a 16-by-16 block of new data values is read from memory unit 116 into SRAM 0 4600. The new and old 16-by-16 blocks are again provided to the tree processor/encoder/decoder for processing, the inverse quantized 16-by-16 old block being written into SRAM 2 4602. During the period of time when the tree processor/encoder/decoder circuit 124 is performing tree processing and encoding, the inverse quantized 16-by-16 block in SRAM 1 4601 is written back to 16-by-16 block 3 of the old portion of memory unit 116. Subsequently, in the seventh time period, 16-by-16 block 5 of the old portion of memory unit 116 is read into SRAM 1 4601 and in the eighth time period the 16-by-16 block of new data values 4 in memory unit 116 is read into SRAM 0 4600. In the ninth time period, tree processor/encoder/decoder circuit 124 processes the 16-by-16 new and old blocks 4 and 5 while the 16-by-16 block of inverse quantized data values in SRAM 2 4602 is written to 16-by-16 block 2 in the old portion of memory unit 116. This pipelining technique allows the dynamic random access memory (DRAM) to be accessed during each time period by taking advantage of the time period when the tree processor/encoder/decoder circuit 124 is processing and not reading from memory unit 116. Because all accesses of memory unit 116 are directed to 16-by-16 blocks of memory locations, the number of CAS cycles is maximized. Arrows are provided in Figure 46 between memory unit 116 and video encoder/decoder circuit 112 to illustrate the accessing of various 16-by-16 blocks of the new and old sub-band decompositions during different time periods. However, because video encoder/decoder chip 112 only has one set of data leads through which data values can be read from and written to memory unit 116, the input/output ports on the right sides of dual port static random access memories 4600-4602 are bussed together and coupled to the input/output data pins of the video encoder/decoder chip 112.

In order to avoid the necessity of providing an additional memory to realize first-in-first-out (FIFO) memory 120, SRAM 3 4603, which is used as a line delay in the column convolver of the video encoder/decoder chip 112, is coupled to the tree processor/encoder/decoder circuit 124 to buffer the compressed data stream for encoding and decoding operations between the ISA bus 106 and the video encoder/decoder chip 112. This

sharing of SRAM 3 is possible because the discrete wavelet transform circuit 122 operates in a first time period and the tree processor/encoder-decoder circuit 124 operates in a second time period.

When the tree processor/encoder/decoder circuit 124 is performing the decoding function, the new portion of memory unit 118 is not required and SRAM 0 is unused. The read 0, read 1, and read 4 time periods of the time line illustrated in Figure 47 are therefore omitted during decoding.

Although the present invention has been described by way of the above described specific embodiments, the invention is not limited thereto. Adaptations, modifications, rearrangements and combinations of various features of the specific embodiments may be practiced without departing from the scope of the invention. For example, an integrated circuit chip may be realized which performs compression but not decompression and another integrated circuit chip may be realized which performs decompression but not compression. Any level of integration may be practiced including placing memory units on the same chip with a discrete wavelet transform circuit and a tree processor/encoder-decoder circuit. The invention may be incorporated into consumer items including personal computers, video cassette recorders (VCRs), video cameras, televisions, compact disc (CD) players and/or recorders, and digital tape equipment. The invention may process still image data, video data and/or audio data. Filters other than four coefficient quasi-Daubechies forward transform filters and corresponding four coefficient reconstruction (inverse transform) filters may be used including filters disclosed in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". Various start and end forward transform filters and various corresponding start and end reconstruction (inverse transform) filters may also be used including filters disclosed in copending Patent Cooperation Treaty (PCT) application filed March 30, 1994, entitled "Data Compression and Decompression". Tokens may be encoded or unencoded. Other types of tokens for encoding other information including motion in consecutive video frames may be used. Other types of encoding other than Huffman encoding may be used and different quantization schemes may be employed. The above description of the preferred embodiments is therefore presented merely for illustrative instructional purposes and is not intended to limit the scope of the invention as set forth in the appended claims.

APPENDIX A: VHDL Language Implementation of CONTROL_ENABLE Block 3420

--The state machine to control the address counter#
 --works for 3 octave decomposition in y & 2 in u/v#

use work.DWT_TYPES.all;
 use work.diff_package.all;

entity U_CONTROL_ENABLE is

```
port(
  ck : in bit ;
  reset : in t_reset ;
  new_channel,channel : in t_channel ;
  c_blk : in BIT_VECTOR(1 to 3) ;
  subband : in BIT_VECTOR(1 to 2) ;
  load_channel : in t_load ;
  new_mode : in t_mode ;

  out_1 : out BIT_VECTOR(1 to 3) ;
  out_2 : out t_octave;
  out_3 : out bit;
  out_4 : out bit;
  out_5 : out t_state) ;
```

end U_CONTROL_ENABLE;

architecture behave OF U_CONTROL_ENABLE is

```
signal      state:t_state;
signal      new_state_sig:t_state;
BEGIN
```

state_machine:PROCESS(reset,new_channel,channel,c_blk,subband,load_channel,new_mode,state,new_state_sig)

```
VARIABLE en_blk:BIT_VECTOR(1 to 3) := B"000";
--enable blk_count#
```



```

5
10
15
20
25
30
35
40
45
50
55

variable lpf_block_done:bit := '0';
--enable x_count for LPF#
variable tree_done:bit := '0';
--enable x_count for other subbands#
variable reset_state:t_state;
variable new_state:t_state;
variable octave:t_octave := 0;
--current octave#
variable start_state:t_state;
-- dummy signals for DFI

BEGIN
-- default initial conditions
en_blk:=b"000";
lpf_block_done:= '0';
tree_done:= '0';
octave:= 0;
reset_state:=up0;
new_state:=state;
start_state:=up0;
--set up initial state thro mux on reset, on HH stay in z=0 state#
CASE channel IS
WHEN u|v => start_state:= down1;
WHEN y => start_state:= up0;
END CASE;
CASE reset IS
WHEN rst => reset_state:= start_state;
WHEN OTHERS => reset_state := state;
END CASE;
CASE reset_state IS
WHEN up0 => octave :=2;
en_blk(3):= '1';
CASE c_blk(3) IS
WHEN '1' =>
CASE eubband IS
WHEN B"00" => lpf_block_done := '1';
--clock x_count for LPF y channel#

```

```

5
10
15
20
25
30
35
40
45
50
55

--change state when count done#
WHEN OTHERS => new_state := up1;
END CASE;

CASE new_mode IS
WHEN stop => tree_done := '1';
WHEN OTHERS => null;
END CASE;
WHEN OTHERS => null;
END CASE;

WHEN up1 => octave := 1;
en_blk(2) := '1';
CASE c_blk(2) IS
WHEN '1' => new_state := zz0;

CASE new_mode IS
--in luminance, terminate branch & move to next branch#
WHEN stop => new_state := down1;
WHEN en_blk(3) := '1';
WHEN OTHERS => null;
END CASE;
WHEN OTHERS => null;
END CASE;

WHEN zz0 => octave := 0;
en_blk(1) := '1';
CASE c_blk(1) IS
WHEN '1' => new_state := zz1;
WHEN en_blk(2) := '1';
WHEN OTHERS => null;
END CASE;
WHEN zz1 => octave := 0;
en_blk(1) := '1';
CASE c_blk(1) IS
WHEN '1' => new_state := zz2;
WHEN en_blk(2) := '1';
WHEN OTHERS => null;
END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

    END CASE;
    WHEN zz2 => octave := 0;
        en_blk(1) := '1';
        CASE c_blk(1) IS
            WHEN '1' => new_state := zz3;
            en_blk(2) := '1';
            OTHERS => null;
        WHEN
            END CASE;
    WHEN zz3 => octave := 0;
        en_blk(1) := '1';
        --now decide the next state, on block(1) carry check the other block carries
        --now decide the next state, on block(1) carry check the other block carries
        --
        CASE c_blk(1) IS
            WHEN '1' => new_state := down1;
            en_blk(2) := '1';
            en_blk(3) := '1';
        --roll over to 0#
        --because state zz3 clock 1 pulse#
        WHEN
            END CASE;
    WHEN down1 => octave := 1;
        en_blk(2) := '1';
        CASE c_blk(2) IS
            WHEN '1' => CASE subband IS
                WHEN B"00" => lpf_block_done := '1';
                OTHERS => new_state := zz0;
            END CASE;
            CASE new_mode IS
                WHEN stop => CASE channel IS
                    WHEN u/v => tree_done := '1';

```

```

5
10
15
20
25
30
35
40
45
50
55

--move to next tree#
WHEN y => en_blk(3) := '1';
CASE c_blk(3) IS
  WHEN '1' => tree_done := '1';
  WHEN OTHERS => new_state := down1;
END CASE;

WHEN OTHERS => null;
END CASE;

WHEN OTHERS => null;
END CASE;

WHEN OTHERS => null;
END CASE;

WHEN OTHERS => null;
END CASE;

END CASE;

CASE channel IS
  WHEN u'v => IF c_blk(1)='1' AND c_blk(2)='1' THEN tree_done := '1';
               ELSE null;
               END IF;

  WHEN y => IF c_blk(1)='1' AND c_blk(2)='1' AND c_blk(3)='1' THEN
               tree_done := '1';
               ELSE null;
               END IF;

  END CASE;

--now change to start state if the sequence has finished#
CASE tree_done IS
  --in LPF state doesn't change when block done#
  WHEN '1' => new_state := start_state;
  WHEN OTHERS => null;
END CASE;

--on channel change, use starting state for new channel#
CASE load_channel IS
  --in LPF state doesn't change when block done#
  WHEN write => CASE new_channel IS

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN y => new_state:= up0;
    WHEN u|v => new_state:=down1;
    END CASE;
    OTHERS => null;
  WHEN
  END CASE;

  new_state_sig<=new_state;

  out_1 <= en_blk;
  out_2 <= octave;
  out_3 <= tree_done;
  out_4 <= lpf_block_done;
  out_5 <=reset_state;

  END PROCESS;

  DPl(ck,new_state_sig,state);
  END behave;

  CONFIGURATION CONTROL_ENABLE_CON OF U_CONTROL_ENABLE is
  FOR behave
  END FOR;
  END CONTROL_ENABLE_CON;

```

APPENDIX B: VHDL Language Implementation of MODE_CONTROL Block 3118

--generates the new_mode from the old, and outputs control signals to the tokeniser--

```
use work.DWT_TYPES.all;
use work.diff_package.all;
```

```
entity U_MODE_CONTROL IS
PORT(
```

```
  ck : in bit ;
  reset : in t_reset ;
  intra_inter : in t_intra ;
  lpf_done : in bit ;
  flags : in BIT_VECTOR(1 to 7);
  token_in : in BIT_VECTOR(1 to 2) ;
  octave : in t_octave ;
  state : in t_state ;
  direction : in t_direction ;
  load_mode_in : in t_load ;
  cycle : in t_cycle ;
```

```
  out_1:out t_mode;
  out_2:out t_mode;
  out_3:out BIT_VECTOR(1 to 2) ;
  out_4:out t_diff;
  out_5:out BIT_VECTOR(1 to 2) ;
  out_6:out t_mode);
```

```
end U_MODE_CONTROL;
```

architecture behave of U_MODE_CONTROL is

```

5
10
15
20
25
30
35
40
45
50
55

--new_mode,proposed_mode,current_token,difference,token_length, --
signal nzflag:bit;
signal origin:bit;
signal noflag:bit;
signal ozflag:bit;
signal motion:bit;
signal pro_new_zibit;
signal pro_no_zibit;
signal lpf_done_del:bit;
signal load_modeit_load_vec(1 to 4);
signal load_nextit_load;
signal pre_mode_sig:t_mode;
signal pro_mode_sig:t_mode;
signal new_mode_sig:t_mode;
signal modeit_mode;
signal diff_sig:t_diff;
signal diff_out:t_diff;
signal mode_regs:t_mode_vec(1 to 4);
BEGIN

    nzflag <= flags(1);
    origin <= flags(2);
    noflag <= flags(3);
    ozflag <= flags(4);
    motion <= flags(5);
    pro_new_z <= flags(6);
    pro_no_z <= flags(7);

    DFL(ck,lpf_done,lpf_done_del);

--the proposed value for the mode at that octave, flags etc will change this value as necessary--
--proposed, or inherited mode from previous tree--
MODE_CONTROL:PROCESS( nzflag,origin,noflag,ozflag,motion,pro_new_z,pro_no_z,lpf_done_del,token_in,direction,
    mode_regs ,state,reset,intra_inter,octave)

variable      pro_mode :t_mode;

```

```

5
10
15
20
25
30
35
40
45
50
55

variable      new_mode :t_mode/
variable      token_out :bit_vector(1 to 2);
variable      difference :t_diff;
variable      token_length :bit_vector(1 to 2);
variable      pro_flag :bit;

BEGIN

--initialise variables

CASE reset IS
WHEN rst => CASE intra_inter IS
--reset on frame start, so do lpf--
WHEN intra => pro_mode:= lpf_etill;
WHEN OTHERS => pro_mode:=lpf_send;
END CASE;
WHEN OTHERS => CASE lpf_done_del IS
WHEN '1' => CASE intra_inter IS
WHEN intra => pro_mode:=etill;
WHEN OTHERS => pro_mode:= send;
END CASE;
--store default mode in mode(4)--
WHEN OTHERS => CASE state IS
WHEN down1 => pro_mode:=
WHEN up0 => pro_mode:=
WHEN OTHERS => CASE
mode_regs(3);
--jump sideways in oct 1--
mode_regs(4);
octave IS
WHEN 0 =>pro_mode:= mode_regs(1);
WHEN 1 =>pro_mode:= mode_regs(2);
WHEN 2 =>pro_mode:= mode_regs(3);

```



```

5
10
15
20
25
30
35
40
45
50
55

WHEN 3 => pro_mode := mode_regs(4);
CASE;
END CASE;

END CASE;

--Inherit the previous mode--
new_mode := pro_mode;
token_out := B"00";
difference := nodiff;
token_length := B"00";
pro_flag := '0';

CASE direction IS
WHEN forward =>

CASE pro_mode IS

WHEN lpf_stop|stop => null;
WHEN void => CASE ozflag IS
WHEN '1' => new_mode := stop;
WHEN OTHERS => null;
END CASE;

WHEN void_still => null;
--Intra so must zero out all of tree--

WHEN still_send => token_length := B"01";
IF nzflag='1' OR pro_new_z='1' THEN token_out :=
B"00";

CASE ozflag IS
WHEN '1' => new_mode
:= stop;
WHEN OTHERS =>
new_mode := void;
END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

ELSE token_out := B"10";
    new_mode := still_send;
END IF;

WHEN send => CASE ozflag IS
    WHEN '1' => token_length := B"01";

        IF nzflag = '1' OR pro_new_z='1' THEN

            token_out := B"10";
            new_mode :=

        END IF;

        OTHERS => token_length := B"10";

        IF (NOT(nzflag) = '1' OR motion = '1') AND
        THEN

            CASE origin IS
            WHEN '1' => pro_flag :=

            WHEN OTHERS => pro_flag :=

            END CASE;

            CASE pro_flag IS
            WHEN '1' => token_out

            WHEN OTHERS => CASE

```

```

5
10
15
20
25
30
35
40
45
50
55

WHEN '1' => token_out := B"01";
new_mode:= still_send;
WHEN OTHERS => token_out := B"11";
new_mode:= send;
CASE;
END CASE;
END

ELSE
IF (motion = '1' OR origin
THEN
token_out
ELSE
token_out :=
END IF;
END IF;
END CASE;

WHEN still => token_length := B"01";
IF nzflag = '1' OR pro_new_z = '1'
THEN token_out := B"00";
new_mode:= void_still;
ELSE token_out := B"10";
--zero out tree--

```

```

5
10
15
20
25
30
35
40
45
50
55

new_mode := still;
END IF;

WHEN lpf_still => token_out := B"00";
token_length := B"00";

WHEN lpf_send => difference := diff;
token_length := B"01";

IF noflag = '1' OR pro_no_2 = '1'
THEN
token_out := B"00";
new_mode := lpf_stop;
ELSE
token_out := B"10";
new_mode := lpf_send;
END IF;

END CASE;

--as mode stop but for this block only--

WHEN inverse =>
CASE pro_mode IS
WHEN lpf_stop|stop => null;

WHEN void => CASE ozflag IS
WHEN '1' => new_mode := stop;
WHEN OTHERS => null;
END CASE;

WHEN void_still => null;

WHEN send => CASE ozflag IS
WHEN '1' => token_length := B"01";

CASE token_in(1) IS
WHEN '1' => new_mode :=
WHEN '0' => new_mode :=

still_send;
stop;

--repeat of still-send code--

```

```

5
10
15
20
25
30
35
40
45
50
55

:= diff;
new_mode := send;
:= still_send;
:= void;
:= stop;

END CASE;

WHEN OTHERS => token_length := B"10";
CASE token_in IS
WHEN B"11" => difference

WHEN B"01" => new_mode
WHEN B"10" => new_mode
WHEN B"00" => new_mode
END CASE;

END CASE;

WHEN still_send => token_length := B"01";
CASE token_in(1) IS
WHEN '1' => new_mode := still_send;
WHEN '0' => CASE ozflag IS
WHEN '1' =>

new_mode := stop;
=> new_mode := void;

END CASE;

WHEN still => token_length := B"01";
CASE token_in(1) IS
WHEN '1' => new_mode := still;
WHEN '0' => new_mode := void_still;
END CASE;

WHEN lpf_send => difference := diff;
token_length := B"01";

```

```

5
10
15
20
25
30
35
40
45
50
55

CASE token_in(1) IS
  WHEN '0' => new_mode := lpf_stop;
  WHEN '1' => new_mode := lpf_end;
END CASE;

WHEN lpf_still => null;
END CASE;

--relate variable to corresponding signals
out_2 <= pro_mode;
pro_mode_sig <= pro_mode;
out_3 <= token_out;
out_5 <= token_length;
out_6 <= new_mode;
new_mode_sig <= new_mode;
diff_sig <= difference;

END PROCESS MODE_CONTROL;

out_1 <= mode;
out_4 <= diff_out;

pre_mode_sig <= pro_mode_sig WHEN reset = rst OR lpf_done_del = '1' ELSE
mode;

--save the new mode & difference during a token cycle, when the flag and tokens are valid--
-- on lpf_still & inverse no token cycles so load on skip cycle, just so next_mode is defined

load_next <= write WHEN cycle = token_cycle ELSE
write WHEN cycle = skip_cycle AND pro_mode_sig=lpf_still AND direction = inverse ELSE
read;

```

```

5
10
15
20
25
30
35
40
45
50
55

DFF_INIT(ck,no_rst,load_next,new_mode_sig,mode);
DFF_INIT(ck,no_rst,load_next,diff_sig,diff_out);

--now write the new mode value into the mode stack at end of cycle, for later use --
--dont update modes at tree base from lpf data, on reset next(1) is undefined--
--store base mode in mode(3)& mode(4), base changes after lpf--

load_mode <= (read,read,write,write) WHEN reset=rst OR lpf_done_del= '1' ELSE
              (write,write,read,read) WHEN octave= 1 AND load_mode_in= write ELSE
              (read,write,write,read) WHEN octave = 2 AND load_mode_in=write ELSE
              (read,read,read,read) ;

DFF_INIT(ck,no_rst,load_mode(1),pre_mode_sig,mode_regs(1));
DFF_INIT(ck,no_rst,load_mode(2),pre_mode_sig,mode_regs(2));
DFF_INIT(ck,no_rst,load_mode(3),pre_mode_sig,mode_regs(3));
DFF_INIT(ck,no_rst,load_mode(4),pre_mode_sig,mode_regs(4));

END behave;

CONFIGURATION MODE_CONTROL_CON OF U_MODE_CONTROL is
FOR behave
END FOR;
END MODE_CONTROL_CON;

```

APPENDIX C: VHDL Language Implementation of CONTROL_COUNTER 3124

```

--a counter to control the sequencing ofw, token, huffman cycles--
--decide reset is enabled 1 cycle early, and latched to avoid glitches--
--lpf_stop is a dummy mode to disable the block writeshuffman data--
--cycles for that block--
use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_CONTROL_COUNTER IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  mode,new_mode : in t_mode ;
  direction : in t_direction ;

  out_0 : out t_load;
  out_1 : out t_cycle;
  out_2 : out t_reset;
  out_3 : out bit;
  out_4 : out bit;
  out_5 : out t_load;
  out_6 : out t_cs;
  out_7 : out t_load;
  out_8 : out t_cs) ;

--mode load,cycle,decide reset,read_addr_enable,write_addr_enable,load flags--
--decode write_addr_enable early and latch to avoid feedback loop with pro_mode--
--in MODES CONTROL--
end U_CONTROL_COUNTER;

architecture behave of U_CONTROL_COUNTER IS
  COMPONENT COUNT_SYNC
  GENERIC (n:integer);
  PORT(

```


5

10

15

20

25

30

35

40

45

50

55

```

ck:in bit ;
reset:in t_reset;
en:in bit;
q:out bit_vector(1 to n);
carry:out bit;
end COMPONENT;

signal write_del:bit;
signal write_sig:bit;
signal decide_del:t_reset;
signal decide_sig:t_reset;
signal count_reset:t_reset;
signal count_len:t_length;
signal count_1:BIT_VECTOR(1 to 4);
signal count_2:bit;
signal always_one:bit:= '1';
BEGIN

count_len <= U_TO_I(count_1);

control:PROCESS(ck,count_reset,direction,mode,new_mode,count_len)

VARIABLE
VARIABLE
VARIABLE
VARIABLE
VARIABLE
VARIABLE
VARIABLE
VARIABLE

cycle : t_cycle;
decide_reset : t_reset;
load_mode : t_load;
load_flags : t_load;
cs_new : t_cs;
cs_old : t_cs;
rw_old : t_load;
read_addr_enable : bit;
write_addr_enable : bit;

BEGIN

cycle := skip_cycle;
decide_reset := no_reset;
load_mode := read;
load_flags := read;

```

```

5
10
15
20
25
30
35
40
45
50
55

cs_new := no_sel;
cs_old := sel;
rw_old := read;
read_addr_enable := '0';
write_addr_enable := '0';

CASE direction IS
  WHEN forward =>
    CASE mode IS
      WHEN send|still_send|lpf_send =>
        CASE count_len IS
          0 to 3 => read_addr_enable := '1';
                  cs_new := sel;
          4 => cycle := token_cycle;
              load_flag := write;
              write_addr_enable := '1';
          5 to 7 => write_addr_enable := '1';
                  CASE new_mode IS
                    WHEN stop|lpf_stop =>

cycle := skip_cycle;

rw_old := read;

cs_old := no_sel;
skip_cycle;

rw_old := write;
data_cycle;

rw_old := write;

cycle := skip_cycle;

      WHEN void => cycle :=

      WHEN OTHERS => cycle :=

      END CASE;
    CASE new_mode IS
      WHEN stop|lpf_stop =>

```

```

5
10
15
20
25
30
35
40
45
50
55

rw_old:= read;
cs_old:= no_sel;
skip_cycle;
load_mode:= write;
rw_old:= write;

data_cycle;
load_mode:= write;
rw_old:= write;

:= skip_cycle;
data_cycle;

WHEN void => cycle :=

WHEN OTHERS => cycle :=

END CASE;

WHEN OTHERS => null;
END CASE;

CASE count_len IS
WHEN 0 to 3 => read_addr_enable := '1';
cs_new:= sel;
WHEN 4 => cycle := token_cycle;
write_addr_enable := '1';
load_flag:= write;
WHEN 5 to 7 => rw_old := write;
write_addr_enable := '1';
CASE new_mode IS
WHEN void_still => cycle

WHEN OTHERS => cycle :=

END CASE;

WHEN 8 => decide_reset := reset;

```

```

5
10
15
20
25
30
35
40
45
50
55

    rw_old:= write;
    load_mode:= write;
    CASE new_mode IS
    WHEN void_still => cycle
    WHEN OTHERS => cycle :=
        END CASE;

    WHEN OTHERS => null;
    END CASE;

    WHEN lpf_still => CASE count_len IS
    WHEN 0 to 3 => read_addr_enable := '1';
        cs_new:= sel;
    WHEN 4 => cycle := token_cycle;
        write_addr_enable := '1';
        load_flag:= write;
    WHEN 5 to 7 => cycle := data_cycle;
        rw_old:= write;
        write_addr_enable := '1';
    WHEN 8 => cycle := data_cycle;
        rw_old:= write;
        decide_reset:= ret;
        load_mode:= write;

    WHEN OTHERS => null;
    END CASE;

    WHEN void => CASE count_len IS
    WHEN 0 to 3 => read_addr_enable := '1';
        cs_new:= sel;
    WHEN 4 => load_flag := write;
        cycle:= token_cycle;

    WHEN 5 to 7 => write_addr_enable := '1';
        write_addr_enable := '1';

    --dummy token cycle for mode update--
    --keep counters going--

```

```

5
10
15
CASE new_mode IS
  WHEN stop => rw_old :=
20
25
    WHEN OTHERS => rw_old :=
      END CASE;
    WHEN 8 => decide_reset := rst;
      CASE new_mode IS
        WHEN stop => rw_old :=
30
35
          WHEN OTHERS => load_mode
            END CASE;
          WHEN OTHERS => null;
        END CASE;
        CASE count_len IS
          WHEN 0 => write_addr_enable := '1';
          WHEN 1 to 3 => write_addr_enable := '1';
          WHEN 4 => rw_old := write;
            load_mode := write;
            decide_reset := rst;
          WHEN OTHERS => null;
        END CASE;
40
45
        WHEN OTHERS => null;
        END CASE;
        CASE mode IS
50
55
          WHEN inverse =>

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN send|still_send|lpf_send =>
        CASE count_len IS
            WHEN 0 to 3 => read_addr_enable := '1';
            WHEN 4 => cycle := token_cycle;
                    write_addr_enable := '1';
                    load_flags := write;
            WHEN 5 to 7 => write_addr_enable := '1';
                        CASE new_mode IS
                            WHEN stop|lpf_stop =>

cycle := skip_cycle;

rw_old := read;

cs_old := no_sel;

skip_cycle;

rw_old := write;

data_cycle;

rw_old := write;

        WHEN void => cycle :=

        WHEN OTHERS => cycle :=

            END CASE;
            WHEN 8 => decide_reset := rst;
            CASE new_mode IS
                WHEN stop|lpf_stop =>

        WHEN void => cycle :=

```

```

5      WHEN OTHERS => cycle :=
10
15      END CASE;
20
25      WHEN OTHERS => null;
30      END CASE;
35      WHEN still =>
40      CASE count_len IS
45      WHEN 0 => null ;
50      WHEN 1 => cycle := token_cycle;
55      write_addr_enable := '1';
      WHEN 2 to 4 => rw_old := write;
      write_addr_enable := '1';
      CASE new_mode IS
      WHEN void_still => cycle
      := skip_cycle;
      WHEN OTHERS => cycle :=
      data_cycle;
      END CASE;
      WHEN 5 => rw_old:=write;
      decide_reset:= rat;
      load_mode:= write;
      CASE new_mode IS
      WHEN void_still => cycle
      := skip_cycle;
      WHEN OTHERS => cycle :=
      data_cycle;
      END CASE;
      WHEN OTHERS => null;
      END CASE;
      WHEN lpf_still => CASE count_len IS
      WHEN 0 =>null ;

```

```

--match with previous--
--skip for write.enb delay--

WHEN 1 => write_addr_enable := '1';
WHEN 2 to 4 => cycle := data_cycle;
               rw_old := write;
               write_addr_enable := '1';
WHEN 5 => cycle := data_cycle;
               rw_old := write;
               decide_reset := rat;
               load_mode := write;

WHEN OTHERS => null;
END CASE;
CASE count_len IS
  WHEN 0 to 3 => read_addr_enable := '1';
  WHEN 4 => load_flag := write;
               cycle := token_cycle;

               write_addr_enable := '1';
               write_addr_enable := '1';
               CASE new_mode IS
                 WHEN stop => rw_old :=

                     WHEN OTHERS => rw_old :=

                     BND CASE;
                 WHEN 8 => decide_reset := rat;
                     CASE new_mode IS
                       WHEN stop => rw_old :=

                       WHEN OTHERS => load_mode

read;
cs_old := no_sel;
write;

read;
cs_old := no_sel;
:= write;
rw_old := write;

```



```

5
10
15
20
25
30
35
40
45
50
55

END CASE;

WHEN OTHERS => null;
END CASE;

CASE count_len IS
  WHEN 0 => null;

  WHEN 1 => write_addr_enable := '1';

  WHEN 2 to 4 => write_addr_enable := '1';
  rw_old := write;

  WHEN 5 => rw_old := write;
  load_mode := write;
  decide_reset := ret;

  WHEN OTHERS => null;
END CASE;

WHEN void_still =>

--match with reset--
--dummy as write delayed--

END CASE;

WHEN OTHERS => null;
END CASE;

write_sig <= write_addr_enable;
decide_sig <= decide_reset;

DP7(ck, reset, write_sig, write_del);
out_0 <= load_mode;
out_1 <= cycle;
out_2 <= decide_sig;
out_3 <= read_addr_enable;
out_4 <= write_del;
out_5 <= load_flag;
out_6 <= cs_new;
out_7 <= rw_old;
out_8 <= cs_old;

END PROCESS;

WITH reset SELECT
count_reset <= ret WHEN ret,

```

5

10

15

20

25

30

35

40

45

50

55

decide_sig WHEN OTHERS;

control_cnt: count_sync GENERIC MAP(4) PORT MAP(ck,count_reset,always_one,count_1,count_2);

END behave;

CONFIGURATION CONTROL_COUNTER_CON OF U_CONTROL_COUNTER IS
FOR behave

FOR ALL:count_sync USE ENTITY WORK.count_sync(behave);

END FOR;

END FOR;

END CONTROL_COUNTER_CON;

APPENDIX D: VHDL Language Implementation of Video Encoder/Decoder Integrated Circuit Chip

```
--VHDL Description of Discrete Wavelet Transform Circuit--
```

```
--the string base address calculators--
```

```
use WORK.dwt_types.all;
```

```
use WORK.utils.all;
```

```
use WORK.utils_dwt.all;
```

```
use WORK.dff_package.all;
```

```
entity U_NOMULT is
```

```
PORT(
```

```
    ck : in bit ;
```

```
    reset : in t_reset ;
```

```
    col_end : in bit ;
```

```
    mux_control : in t_mux4 ;
```

```
    incr : in t_memory_addr;
```

```
    oct_add_factor : in t_memory_addr ;
```

```
    base_u, base_v : in BIT_VECTOR(1 to 19);
```

```
    out_1 : out t_memory_addr;
```

```
end U_NOMULT;
```

```
architecture behave of U_NOMULT is
```

```
    signal mux:t_memory_addr;
```

```
    signal next_addr:t_memory_addr;
```

```
    signal dff_out:t_memory_addr;
```

```
    signal addit_memory_addr;
```

```
BEGIN
```

```
    WITH mux_control SELECT
```

```
    next_addr <= add
```

```
    0
```

```
    U_TO_I(base_u)
```

```
    U_TO_I(base_v)
```

```
    WHEN uno,      WHEN dos,
    WHEN tres,     WHEN quatro;
```

```

5
10
15
20
25
30
35
40
45
50
55

BEGIN
    temp <= '0' WHEN reset=rat ELSE
        '1' WHEN j='1' ELSE
            q ;

    DFF1(ck,temp,q);
    out_1 <= q;
    end behave;

    CONFIGURATION JKFF_CON OF JKFF is
    FOR behave
    END FOR;
    END JKFF_CON;

    use WORK.dwt_types.all;
    use WORK.utils.all;
    use WORK.utils_dwt.all;
    use WORK.dff_package.all;

    entity TOGGLE IS
    PORT(
        ck : in bit ;
        reset : in t_reset ;
        j:in bit;

        out_1:out bit);
    end TOGGLE;

    architecture behave of TOGGLE is
    signal temp:bit;
    signal q:bit;
    BEGIN
        temp <= j XOR q;
        DFF(ck,reset,temp,q);
        out_1 <= q;
    end behave;

    DFF(ck,reset,next_addr,dff_out);

    WITH col_end SELECT
        mux <= incr WHEN '0',
            oct_add_factor WHEN '1';

    add<= dff_out + mux;

    --architecture outputs--
    out_1 <= dff_out;

    END behave;

    CONFIGURATION NOMULT_CON OF U_NOMULT is
    FOR behave
    END FOR;
    END NOMULT_CON;

    -- a toggle flip-flop
    use WORK.dwt_types.all;
    use WORK.utils.all;
    use WORK.utils_dwt.all;
    use WORK.dff_package.all;

    entity JKFF IS
    PORT(
        ck : in bit ;
        reset : in t_reset ;
        j:in bit;

        out_1:out bit);
    end JKFF;

    architecture behave of JKFF is
    signal temp:bit;
    signal q:bit;

```

```

5
10
15
20
25
30
35
40
45
50
55

CONFIGURATION TOGGLE_CON OF TOGGLE is
FOR behave
END FOR;
END TOGGLE_CON;

-----
--the read and write address generator, input the initial image & block sizes for octave 0 for the y channel--
use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;

entity U_ADDR_GEN IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  direction : in t_direction ;
  channel : in t_channel ;
  x_p_1 : in BIT_VECTOR(1 to 10) ;
  x3_p_1 : in BIT_VECTOR(1 to 12) ;
  x7_p_1 : in BIT_VECTOR(1 to 13) ;
  octave_row_length : in BIT_VECTOR (1 to ysize) ;
  octave_col_length : in BIT_VECTOR (1 to xsize) ;
  octave_reset : in t_reset ;
  octave : in t_octave ;
  y_done : in bit ;
  uv_done : in bit ;
  octave_finished : in t_load ;
  base_u_base_v : in BIT_VECTOR(1 to 19) ;

  out_1 : out t_input_mux;
  out_2_1 : out t_memory_addr; -- memory port
  out_2_2 : out t_memory_addr;
  out_2_3 : out t_load;

  out_3_1 : out t_load; --dwt in control

```

```

5
10
15
20
25
30
35
out_3_2 : out t_cs;

out_4 : out t_load;      --IDWT data valid
out_5 : out t_load;      --read valid
out_6 : out t_count_control; --row read

out_7_1 : out t_col;
out_7_2 : out t_count_control;
and U_ADDR_GEN;

```

--the current octave and when the block finishes the 3 octave transform--

architecture behave OF U_ADDR_GEN IS

COMPONENT U_MEM_CONTROL

PORT(

```

    ck : in bit ;
    reset : in t_reset ;
    direction : in t_direction ;
    channel : in t_channel ;
    octave : in t_octave ;
    addr_w,addr_r : in t_memory_addr ;
    zero_hh : in t_load ;

```

```

out_1 : out t_input_mux;

```

```

out_2_1 : out t_memory_addr;
out_2_2 : out t_memory_addr;
out_2_3 : out t_load;

```

```

out_3_1 : out t_load;
out_3_2 : out t_cs);
end COMPONENT;

```

COMPONENT JKFF

PORT(

```

    ck : in bit ;

```

```

5
10
15
20
25
30
35
40
45
50
55

reset : in t_reset ;
j:in bit;
    out_1:out bit);
end COMPONENT;

COMPONENT U_ROW_COUNT
PORT(
    ck : in bit ;
    reset : in t_reset ;
    octave_cnt_length : in BIT_VECTOR(1 to ysize) ;
    col_carry: in t_count_control;

    out_1 : out t_row;
    out_2 : out t_count_control);
--count value , and flag for count=0,1,2,row_length-1, row_length
end COMPONENT;

COMPONENT U_COL_COUNT
PORT(
    ck : in bit ;
    reset : in t_reset ;
    octave_cnt_length : in BIT_VECTOR(1 to xsize) ;

    out_1 : out t_col;
    out_2 : out t_count_control);
--count value , and flag for count=0,1,2,col_length-1, col_length
end COMPONENT;
COMPONENT U_MULT
PORT(
    ck : in bit ;
    reset : in t_reset ;
    col_end : in bit ;
    mux_control : in t_mux4 ;
    incr : in t_memory_addr;
    oct_add_factor : in t_memory_addr ;
    base_u_base_v : in BIT_VECTOR(1 to 19));

```

```

out_1 : out t_memory_addr;
end COMPONENT;

signal mem_sel : t_mux4;
signal read_mux : t_mux4;
signal write_mux : t_mux4;
signal incr : t_memory_addr;
signal oct_add_factor : t_memory_addr;
signal add_2_y : BIT_VECTOR(1 to 13);
signal add_2_uv : BIT_VECTOR(1 to 13);
signal add_2 : BIT_VECTOR(1 to 13);
signal addr_col_flag : BIT;
signal write_latency : BIT;
signal read_done : BIT;
signal zero_hh : t_load;
signal zero_hh_bit : bit;
signal read_done_bit : bit;
signal start_write_col : t_load;
signal read_valld : t_load;
signal addr_col_2 : t_count_control;
signal addr_row_2 : t_count_control;
signal addr_row_1 : t_row;
signal addr_col_1 : t_col;
signal all_done : bit;
signal all_one : bit;
signal temp0 : bit;
signal temp1 : bit;
signal temp2 : bit;
signal temp3 : bit;
signal temp4 : bit;
signal temp5 : bit;
signal temp6 : bit;
signal read_addr : t_memory_addr;
signal write_addr : t_memory_addr;
signal mem_control : t_input_mux;

```



```

5
10
15
20
25
30
35
40
45
50
55

signal mem_control_2_1:t_memory_addr;
signal mem_control_2_2:t_memory_addr;
signal mem_control_2_3:t_load;
signal mem_control_3_1:t_load;
signal mem_control_3_2:t_cs;
BEGIN

WITH octave SELECT
incr      <= 1      WHEN 0,
          2 WHEN 1,
          4 WHEN 2,
          8 WHEN 3;

WITH octave SELECT
add_2_y   <= B"0000000000001" WHEN 0,
          B"000" & x_p_1(1 to 8) & B"10" WHEN 1,
          B"0" & x3_p_1(1 to 9) & B"100" WHEN 2,
          x7_p_1(1 to 9) & B"1000" WHEN 3;

WITH octave SELECT
add_2_uv  <= B"00000000000001" WHEN 0,
          B"0000" & x_p_1(1 to 7) & B"10" WHEN 1,
          B"00" & x3_p_1(1 to 8) & B"100" WHEN 2,
          B"0" & x7_p_1(1 to 8) & B"1000" WHEN 3;

WITH channel SELECT
add_2     <= add_2_y WHEN Y,
          add_2_uv WHEN OTHERS;

oct_add_factor <= U_TO_I(add_2);

--signals when write must start delayed 1 tu for use in zero_hh--

--decode to bit--
WITH addr_col_2 SELECT
addr_col_flag <= '1' WHEN count_carry ,

```

```

5
10
15
20
25
30
35
40
45
50
55

        '0' WHEN OTHERS;

write_latency <= '1' WHEN addr_row_1 = 2 AND addr_col_1 = conv2d_latency-1 ELSE '0';

--read input data done--
read_done <= '1' WHEN addr_row_2 = count_carry AND addr_col_flag = '1' ELSE '0';

WITH zero_hh_bit SELECT
zero_hh <= write WHEN '1',
read WHEN '0';

WITH read_done_bit SELECT
read_valid <= write WHEN '1',
read WHEN '0';

DFF(ck,reset,zero_hh,start_write_col); --1 to after zero_hh--

read_mux <= tres WHEN y_done='1' AND uv_done='0' AND octave_finished=write AND channel=y ELSE --base u--
tres WHEN y_done='0' AND uv_done='0' AND octave_finished=write AND channel=u ELSE
quatro WHEN y_done='0' AND uv_done='1' AND octave_finished=write AND channel=u ELSE
quatro WHEN y_done='0' AND uv_done='0' AND octave_finished=write AND channel=v ELSE --base v--
dos WHEN y_done='0' AND octave_finished=write AND channel=y ELSE --base y--
uno;

write_mux <= uno WHEN zero_hh =write ELSE -- keep address 0
dos WHEN channel=y ELSE --base y--
tres WHEN channel=u ELSE --base u--
quatro; --base v--

--note that all the counters have to be reset at the end of an octave, ie on octave_finished--

--the, rowfcol, counts, for, the, read, address--

col_map: U_COL_COUNT PORT MAP(ck,octave_reset,octave_col_length,addr_col_1,addr_col_2);
row_map: U_ROW_COUNT PORT MAP(ck,octave_reset,octave_row_length,addr_col_2,addr_row_1,addr_row_2);

```

```

5
10
15
20
25
30
35
40
45
50
55

all_one <='1';

tog_1:JKFF PORT MAP(ck,octave_reset,write_latency,zero_hh_bit);

tog_2:JKFF PORT MAP(ck,octave_reset,read_done,read_done_bit);

--vgr addresses for sparc mem--

--conv_2d PIPELINE DELAY ON THIS FLAG

DF1(ck,addr_col_flag,temp0);
DF1(ck,temp0,temp1);
DF1(ck,temp1,temp2);
DF1(ck,temp2,temp3);
DF1(ck,temp3,temp4);
DF1(ck,temp4,temp5);
DF1(ck,temp5,temp6);

read_map:U_NOMULT PORT MAP(ck,reset,addr_col_flag,read_mux,incr,oct_add_factor,base_u,base_v,read_addr);

write_map:U_NOMULT PORT MAP(ck,reset,temp6,write_mux,incr,oct_add_factor,base_u,base_v,write_addr);

mem_ctrl_map:U_MEM_CONTROL PORT MAP(ck,reset,direction,channel,octave,write_addr,read_addr,zero_hh,
mem_control_1,mem_control_2_1,mem_control_2_2,mem_control_2_3,mem_control_3_1,mem_control_3_2);
--architecture outputs--

out_1 <=mem_control_1;

out_2_1<= mem_control_2_1;
out_2_2 <= mem_control_2_2;
out_2_3 <= mem_control_2_3;

out_3_1 <= mem_control_3_1;
out_3_2 <= mem_control_3_2;

out_4 <=zero_hh;
out_5 <=read_valid;

```

```

5
10
15
20
25
30
35
40
45
50
55

out_6 <=addr_row_2;
out_7_1<=addr_col_1;
out_7_2<=addr_col_2;

END;

CONFIGURATION ADDR_GEN_CON OF U_ADDR_GEN IS
FOR behave
FOR ALL:U_NOMULT USE ENTITY WORK.U_NOMULT(behave);
END FOR;
FOR ALL:U_MEM_CONTROL USE ENTITY WORK.U_MEM_CONTROL(behave);
END FOR;
FOR ALL:U_COL_COUNT USE ENTITY WORK.U_COL_COUNT(behave);
END FOR;
FOR ALL:U_ROW_COUNT USE ENTITY WORK.U_ROW_COUNT(behave);
END FOR;
FOR ALL:JKFF USE ENTITY WORK.JKFF(behave);
END FOR;
END ADDR_GEN_CON;
--the basic 2d convolver for forward transform, rows first then cols for the forward transform--
-- cols first then rows for the inverse transform
use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;

entity U_CONV_2D IS
PORT(
ck : in bit ;
reset : in t_reset ;
in_in : in t_input ;
direction : in t_direction ;
pdel : in t_scratch_array(1 to 4);
conv_reset : in t_reset ;

```

```

row_flag : in t_count_control ;
addr_col_read_1 : in t_col ;
addr_col_read_2 : in t_count_control ;

```

```

out_1 : out t_input ;
out_2_1 : out t_scratch_array(1 to 4) ;
out_2_2 : out t_col ;
out_2_3 : out t_col ;
out_3 : out t_count_control ;
out_4 : out t_count_control ;
out_5 : out t_count_control ;
end U_CONV_2D ;

```

```

--forward direction outputs in row form --

```

```

-- HH HG HH HG to . --
-- HG GG HG GG to . --
-- HH HG HH HG to . --
-- HG GG HG GG to . --

```

```

--the inverse convolver returns the raster scan format output data--
--the convolver automatically returns a 3 octave transform--

```

```

architecture behave OF U_CONV_2D IS

```

```

    COMPONENT U_CONV_ROW

```

```

    PORT(

```

```

        ck : in bit ;
        reset : in t_reset ;
        direction : in t_direction ;
        in_in : in t_input ;
        col_flag : in t_count_control ;

```

```

        out_1 : out t_input ) ;
    end COMPONENT ;

```

```

    COMPONENT U_CONV_COL

```

```

    PORT(

```

```

5
6
7
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
24
25
26
27
28
29
30
31
32
33
34
35
36
37
38
39
40
41
42
43
44
45
46
47
48
49
50
51
52
53
54
55
56
57
58
59
60
61
62
63
64
65
66
67
68
69
70
71
72
73
74
75
76
77
78
79
80
81
82
83
84
85
86
87
88
89
90
91
92
93
94
95
96
97
98
99
100
101
102
103
104
105
106
107
108
109
110
111
112
113
114
115
116
117
118
119
120
121
122
123
124
125
126
127
128
129
130
131
132
133
134
135
136
137
138
139
140
141
142
143
144
145
146
147
148
149
150
151
152
153
154
155
156
157
158
159
160
161
162
163
164
165
166
167
168
169
170
171
172
173
174
175
176
177
178
179
180
181
182
183
184
185
186
187
188
189
190
191
192
193
194
195
196
197
198
199
200
201
202
203
204
205
206
207
208
209
210
211
212
213
214
215
216
217
218
219
220
221
222
223
224
225
226
227
228
229
230
231
232
233
234
235
236
237
238
239
240
241
242
243
244
245
246
247
248
249
250
251
252
253
254
255
256
257
258
259
260
261
262
263
264
265
266
267
268
269
270
271
272
273
274
275
276
277
278
279
280
281
282
283
284
285
286
287
288
289
290
291
292
293
294
295
296
297
298
299
300
301
302
303
304
305
306
307
308
309
310
311
312
313
314
315
316
317
318
319
320
321
322
323
324
325
326
327
328
329
330
331
332
333
334
335
336
337
338
339
340
341
342
343
344
345
346
347
348
349
350
351
352
353
354
355
356
357
358
359
360
361
362
363
364
365
366
367
368
369
370
371
372
373
374
375
376
377
378
379
380
381
382
383
384
385
386
387
388
389
390
391
392
393
394
395
396
397
398
399
400
401
402
403
404
405
406
407
408
409
410
411
412
413
414
415
416
417
418
419
420
421
422
423
424
425
426
427
428
429
430
431
432
433
434
435
436
437
438
439
440
441
442
443
444
445
446
447
448
449
450
451
452
453
454
455
456
457
458
459
460
461
462
463
464
465
466
467
468
469
470
471
472
473
474
475
476
477
478
479
480
481
482
483
484
485
486
487
488
489
490
491
492
493
494
495
496
497
498
499
500
501
502
503
504
505
506
507
508
509
510
511
512
513
514
515
516
517
518
519
520
521
522
523
524
525
526
527
528
529
530
531
532
533
534
535
536
537
538
539
540
541
542
543
544
545
546
547
548
549
550
551
552
553
554
555
556
557
558
559
560
561
562
563
564
565
566
567
568
569
570
571
572
573
574
575
576
577
578
579
580
581
582
583
584
585
586
587
588
589
590
591
592
593
594
595
596
597
598
599
600
601
602
603
604
605
606
607
608
609
610
611
612
613
614
615
616
617
618
619
620
621
622
623
624
625
626
627
628
629
630
631
632
633
634
635
636
637
638
639
640
641
642
643
644
645
646
647
648
649
650
651
652
653
654
655
656
657
658
659
660
661
662
663
664
665
666
667
668
669
670
671
672
673
674
675
676
677
678
679
680
681
682
683
684
685
686
687
688
689
690
691
692
693
694
695
696
697
698
699
700
701
702
703
704
705
706
707
708
709
710
711
712
713
714
715
716
717
718
719
720
721
722
723
724
725
726
727
728
729
730
731
732
733
734
735
736
737
738
739
740
741
742
743
744
745
746
747
748
749
750
751
752
753
754
755
756
757
758
759
760
761
762
763
764
765
766
767
768
769
770
771
772
773
774
775
776
777
778
779
780
781
782
783
784
785
786
787
788
789
790
791
792
793
794
795
796
797
798
799
800
801
802
803
804
805
806
807
808
809
810
811
812
813
814
815
816
817
818
819
820
821
822
823
824
825
826
827
828
829
830
831
832
833
834
835
836
837
838
839
840
841
842
843
844
845
846
847
848
849
850
851
852
853
854
855
856
857
858
859
860
861
862
863
864
865
866
867
868
869
870
871
872
873
874
875
876
877
878
879
880
881
882
883
884
885
886
887
888
889
890
891
892
893
894
895
896
897
898
899
900
901
902
903
904
905
906
907
908
909
910
911
912
913
914
915
916
917
918
919
920
921
922
923
924
925
926
927
928
929
930
931
932
933
934
935
936
937
938
939
940
941
942
943
944
945
946
947
948
949
950
951
952
953
954
955
956
957
958
959
960
961
962
963
964
965
966
967
968
969
970
971
972
973
974
975
976
977
978
979
980
981
982
983
984
985
986
987
988
989
990
991
992
993
994
995
996
997
998
999
1000

```

```

5
10
15
20
25
30
35
40
45
50
55

signal addr_temp4:t_count_control;
signal addr_temp3:t_count_control;
signal del_conv_col:t_input;
signal del_conv_row:t_input;
signal del_conv_init_input;
signal row_init_input;
signal conv_row:t_input;
signal conv_col:t_input;
signal col_init_input;
signal del_init_input;
signal pdel_out:t_scratch_array(1 to 4);
signal wr_addr:t_col;
BEGIN

--reset must be delayed for row convolver depending on direction of transform
DF1(ck,conv_reset,temp1);
DF1(ck,temp1,conv_reset_inv);

WITH direction SELECT
row_reset <= conv_reset WHEN forward,
conv_reset_inv WHEN inverse; --pipeline delays in col_conv--

--reset must be delayed for col convolver depending on direction of transform
DF1(ck,conv_reset_inv,temp2);
DF1(ck,temp2,col_reset_forw);

WITH direction SELECT
col_reset <= col_reset_forw WHEN forward,
conv_reset WHEN inverse; --pipeline delays in row_conv--

-- counter flags must be delayed for col convolver depending on pipelining
DF1(ck,addr_col_read_2,addr_temp1);
DF1(ck,addr_temp1,addr_col_rd_del);

WITH direction SELECT
addr_temp2 <= addr_col_read_2 WHEN forward,

```

```

addr_col_rd_del WHEN inverse;

```

```

DF1(ck,addr_temp2,col_flag);

```

```

-- counter flags must be delayed for row convolver depending on pipelining

```

```

DF1(ck,row_flag,row_temp0);

```

```

DF1(ck,row_temp0,row_temp1);

```

```

DF1(ck,row_temp1,row_temp2);

```

```

DF1(ck,row_temp2,row_temp3);

```

```

WITH direction SELECT

```

```

row_temp4 <= row_temp3 WHEN forward,

```

```

row_flag WHEN inverse;

```

```

DF1(ck,row_temp4,row_control);

```

```

--pipeline delays for col counter, count value

```

```

DF1(ck,addr_col_read_1,col_temp0);

```

```

DF1(ck,col_temp0,col_temp1);

```

```

DF1(ck,col_temp1,col_temp2);

```

```

DF1(ck,col_temp2,col_temp3);

```

```

WITH direction SELECT

```

```

col_count_1 <= col_temp3 WHEN forward,

```

```

addr_col_read_1 WHEN inverse;

```

```

-- similar for carry flag of col counter

```

```

DF1(ck,addr_col_rd_del,addr_temp3);

```

```

DF1(ck,addr_temp3,addr_temp4);

```

```

WITH direction SELECT

```

```

col_count_2 <= addr_temp4 WHEN forward,

```

```

addr_col_read_2 WHEN inverse;

```

```

--pipeline delays for the convolver values and input value--

```

```

DF1(ck,conv_col,del_conv_col);

```



```

5
10
15
20
25
30
35
40
45
50
55

    DF1(ck,conv_row,del_conv_row);

    DF1(ck,in_in,del_in);

    WITH direction SELECT
    row_in <= del_in WHEN forward,
        del_conv_col WHEN inverse;

    row_map: U_CONV_ROW PORT MAP (ck,row_reset,direction,row_in,col_flag,conv_row);

    WITH direction SELECT
    col_in <= del_conv_row WHEN forward,
        del_in WHEN inverse;

    col_map: U_CONV_COL PORT MAP(ck,col_reset,direction,col_in,pdel,row_control,col_count_1,col_count_2,
        conv_col,pdel_out,wr_addr);

    --architecture outputs
    WITH direction SELECT
    out_1 <= del_conv_col WHEN forward,
        del_conv_row WHEN inverse;

    out_2_1 <= pdel_out;
    out_2_2 <= wr_addr;
    out_2_3 <= col_count_1;

    out_3 <= row_control;
    out_4 <= col_count_2;
    out_5 <= col_flag;

    end behave;

    CONFIGURATION CONV_2D_CON OF U_CONV_2D IS
    FOR behave
    FOR ALL:U_CONV_COL      USE ENTITY WORK.U_CONV_COL(behave);
    END FOR;
    FOR ALL:U_CONV_ROW      USE ENTITY WORK.U_CONV_ROW(behave);

```

5

10

15

20

25

30

35

40

45

50

55

```

    END FOR;
  END FOR;
  END CONV_2D_CONV;

  -- 1d col convolver, with control --
  use WORK.dwt_types.all;
  use WORK.utils.all;
  use WORK.utils_dwt.all;
  use WORK.dff_package.all;
  -- a 12 line by line resettable counter for the state machines, out->one on rst--
  --carry active on last element of row--
  entity U_COUNTCOL_2 IS
    PORT(
      ck : in bit ;
      reset : in t_reset ;
      carry: in t_count_control;

      out_1 : out t_count_2 ) ;
    end U_COUNTCOL_2;

  architecture behave OF U_COUNTCOL_2 IS
    signal countdel:t_count_2;
    signal coutout:t_count_2;
  BEGIN
    PROCESS(ck,reset,carry)
    BEGIN
      IF reset = rst THEN coutout <= one;
      ELSIF countdel=one AND carry = count_carry THEN coutout <= two ;
      ELSIF countdel=two AND carry = count_carry THEN coutout <= one;
      ELSE null;
      END IF;
      DF1(ck,coutout,countdel);
      --architecture outputs--
      out_1 <= countdel;
    END PROCESS;
  
```

```

5
10
15
20
25
30
35
40
45
50
55
END behave;
CONFIGURATION COUNTCOL_2_CON OF U_COUNTCOL_2 IS
FOR behave
END FOR;
END COUNTCOL_2_CON;

```

```

use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;
entity U_CONV_COL IS
PORT(
    ck : in bit ;
    reset : in t_reset ;
    direction : in t_direction ;
    in_in : in t_input ;
    pdel : in t_scratch_array(1 to 4) ;
    row_flag : in t_count_control ;
    col_count_1 : in t_col;
    col_count_2 : in t_count_control;

    out_1 : out t_input;
    out_2 : out t_scratch_array(1 to 4);
    out_3 : out t_col;
    end U_CONV_COL;

```

architecture behave OF U_CONV_COL IS

```

--input is data in and, pdel, out from line-delay memories--
-- out is (G,H), and line delay out port. The row counter is started 1 cycle later to allow for--
--pipeline delay between MULTIPLIER and this unit --

```

```

COMPONENT U_COUNTCOL_2
PORT(
    ck : in bit ;

```

```

5
10
15
20
25
30
35
40
45
50
55

    reset : in t_reset ;
    carry: in t_count_control;
    out_1 : out t_count_2 ) ;
    end COMPONENT;

COMPONENT U_ROUND_BITS
PORT(
    in_in :in t_scratch;
    sel:in t_round;

    out_1:out t_input);
    end COMPONENT;

COMPONENT U_MULT_ADD
PORT(
    reset : in t_reset ;
    in_in : in t_input ;
    andsel : in t_and_array(1 to 3) ;
    centermuxsel : in t_mux_array(1 to 2) ;
    muxsel : in t_mux4_array(1 to 3) ;
    muxandsel : in t_and_array(1 to 3) ;
    addsel : in t_add_array(1 to 4) ;
    direction : in t_direction ;
    pdel : in t_scratch_array(1 to 4) ;

    out_1 : out t_scratch_array(1 to 4) );
    end COMPONENT;

    signal row_control:t_count_control;
    signal row_control_del:t_count_control;
    signal col_carry:t_count_control;
    signal reset_row:t_reset;
    signal shift_const:t_round;

    signal andsel:t_and_array(1 to 3);
    signal muxandsel:t_and_array(1 to 3);

```

```

5
10
15
20
25
30
35
40
45
50
55

signal adsel:t_add_array(1 to 4);
signal count:t_count_2;
signal count_delay:t_count_2;
signal centermuxsel:t_mux_array(1 to 2);
signal muxsel:t_mux4_array(1 to 3);
signal mult_add:t_scratch_array(1 to 4);
signal pdel_init:t_scratch_array(1 to 4);
signal pdel_out:t_scratch_array(1 to 4);
signal pdel1_delit_scratch;
signal gh_out:t_scratch;
signal rb_out:t_scratch;

signal col_count_temp:t_col;
signal wr_addr:t_col;
signal rd_addr:t_col;
signal gh_select:t_mux;
BEGIN
--the code for the convolver--

DFF1(ck,reset,reset_row);

--starts row counter 1 cycle after frame start--
--we want the row counter to be 1 cycle behind the col counter for the delay for the--
--pipelined line delay memory--

DFF(ck,reset,col_count_2,col_carry);

--these need to be synchronised to keep the row counter aligned with the data stream--
--also the delay on col_count deglitches the col carryout--

row_control <= row_flag;      --signal for row<=0;1;2;3; last row; etc--

andsel(1) <=  pass WHEN direction = forward AND count=one ELSE
              zero WHEN direction = forward AND count=two ELSE
              pass WHEN direction=inverse AND count=two ELSE
              zero ;

```

```

5
10
15
20
25
30
35
40
45
50
55

WITH row_control AS SELECT
andseel(2) <= zero WHEN count_0 ,
pass WHEN OTHERS;

andseel(3) <= zero WHEN direction=forward AND row_control = count_0 ELSE
pass;

--now the add/sub control for the convolver adders--
WITH count SELECT
addseel <= (add,add,add,subt) WHEN one ,
(add,subt,add,add) WHEN two ;

--now the mux control--
centermuxsel <= (left,right) WHEN (direction = forward AND count = one) OR (direction = inverse AND count = two) ELSE
(right,left);

--the addmuxsel signal--
muxandseel(1 to 2) <= (pass, andseel(2)) WHEN direction = inverse ELSE
(andseel(2), pass) ;
muxandseel(3) <= zero WHEN direction = inverse AND row_control=count_1 ELSE
pass WHEN direction = inverse ELSE
andseel(2);

muxsel(1) <= dos WHEN direction = inverse AND row_control=count_0 ELSE
quatro WHEN direction = inverse AND row_control=count_1 ELSE
tres WHEN direction = inverse AND row_control= count_lml ELSE
dos WHEN direction = inverse ELSE
uno;
muxsel(2) <= tres WHEN direction = inverse AND row_control=count_0 ELSE
dos WHEN direction = inverse AND row_control= count_carry ELSE
uno WHEN direction = inverse ELSE
dos WHEN direction = forward AND row_control=count_0 ELSE
tres WHEN direction = forward AND row_control= count_carry ELSE
uno;
muxsel(3) <= uno WHEN direction = inverse ELSE

```

```

5
10
15
20
25
30
35
40
45
50
55

    tres WHEN direction = forward AND row_control=count_0 ELSE
    quatro WHEN direction = forward AND row_control= count_carry ELSE
    fives;

COUNT_MAP: U_COUNTCOL_2 PORT MAP(ck,reset_row,col_carry, count);

-- set up the r/w address for the line delay memory
--need 2 delays between wr and rd addr

Df1(ck,col_count_1,col_count_temp);
Df1(ck,col_count_temp,wr_addr);

rd_addr <= col_count_1;

--in the control signals to the mult_add block--
MULT_ADD_MAP: U_MULT_ADD PORT
MAP(reset,in_in,andsel,centermuxsel,muxsel,addsel,direction,pdel_out,mult_add);

--delay to catch the write address
Df1(ck,mult_add(1),pdel_in(1));
Df1(ck,mult_add(2),pdel_in(2));
Df1(ck,mult_add(3),pdel_in(3));
Df1(ck,mult_add(4),pdel_in(4));

--read delay to match MULT delay
Df1(ck,pdel(1),pdel_out(1));
Df1(ck,pdel(2),pdel_out(2));
Df1(ck,pdel(3),pdel_out(3));
Df1(ck,pdel(4),pdel_out(4));

Df1(ck,count,count_del);

sh_select <= right WHEN (direction = inverse AND count_del = one) OR (direction = forward AND count_del = two) ELSE
left;

Df1(ck,pdel_out(1), pdell_del);

```

```

5
10
15
20
25
30
35
40
45
50
55

gh_out <= MUX_2(pd_el_in(4),pd_el1_del,gh_select);
DFI(ck,row_control,row_control_del);

shift_const <= shift3 WHEN direction = inverse AND (row_control_del=count_1 OR row_control_del=count_2) ELSE
    'shift4 WHEN direction = inverse ELSE
        shift5;

RB_MAP: U_ROUND_BITS PORT MAP(gh_out,shift_const,rb_out);
--architecture outputs--
out_1 <= rb_out;
out_2 <= pd_el_in;
out_3 <= wr_addr;

END behave;

CONFIGURATION CONV_COL_CON OF U_CONV_COL IS
FOR behave
    FOR ALL:U_ROUND_BITS          USE ENTITY WORK.U_ROUND_BITS(behave);
    END FOR;
    FOR ALL:U_COUNTCOL_2          USE ENTITY WORK.U_COUNTCOL_2(behave);
    END FOR;
    FOR ALL:U_MULT_ADD            USE ENTITY WORK.U_MULT_ADD(behave);
    END FOR;
END CONV_COL_CON;

-- a 12 line by line resettable counter for the state machines, out->one on reset--
use WORK.dwt_types.all;
use WORK.utile.all;
use WORK.utile_dwt.all;
use WORK.dff_package.all;
entity U_COUNT_2 IS
PORT(
    ck : in bit ;
    reset : in t_reset ;

```



```

5
10
15
20
25
30
35
40
45
50
55

out_1 : out t_count_2 ) ;
end U_COUNT_2;

architecture behave of U_COUNT_2 is
    signal countdel:t_count_2;
    signal countout:t_count_2;
begin
    countout <= one when reset = rat or countdel = two else
        two ;
    DFF(ck,countout,countdel);
    --architecture outputs--
    out_1 <= countdel;
end;

CONFIGURATION COUNT_2_CON OF U_COUNT_2 is
    FOR behave
    END FOR;
END COUNT_2_CON;

--the 1d convolver, with control and coeff extend--

use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.dff_package.all;
use WORK.utils_dwt.all;

entity U_CONV_ROW is
    PORT(
        ck : in bit ;
        reset : in t_reset ;
        direction : in t_direction ;
        in_in : in t_input ;
        col_flag : in t_count_control ;

```

```

5
10
15
20
25
30
35
40
45
50
55

out_1 : out t_input ) ;
end U_CONV_ROW;

--
architecture behave of U_CONV_ROW is
-- out is (G,H). The row counter is started 1 cycle later to allow for--
--pipeline delay between MULTIPLIER and this unit --
--the strings give the col & row lengths for this octave--
COMPONENT U_ROUND_BITS
PORT(
in_in : in t_scratch;
sel:in t_round;

out_1:out t_input);
and COMPONENT;

COMPONENT U_COUNT_2
PORT(
ck : in bit ;
reset : in t_reset ;

out_1 : out t_count_2 ) ;
and COMPONENT;

COMPONENT U_MULT_ADD
PORT(
reset : in t_reset ;
in_in : in t_input ;
andsel : in t_and_array(1 to 3) ;
centermuxsel : in t_mux_array(1 to 2) ;
muxsel : in t_mux4_array(1 to 3) ;
muxandsel : in t_and_array(1 to 3) ;
addsel : in t_add_array(1 to 4) ;
direction : in t_direction ;
pdel : in t_scratch_array(1 to 4) ;

```

```

5
10
15
20
25
30
35
40
45
50
55

out_1 : out_t_scratch_array(1 to 4) );
and_COMPONENT;

signal reset_col:it_reset;
signal col_control:it_count_control;
signal temp:it_and;
signal andsel:it_and_array(1 to 3);
signal muxandsel:it_and_array(1 to 3);
signal addsel:it_add_array(1 to 4);
signal count:it_count_2;
signal centermuxsel:it_mux_array(1 to 2);
signal muxsel:it_mux4_array(1 to 3);
signal mult_add:it_scratch_array(1 to 4);
signal pdel:it_scratch_array(1 to 4);
signal pdel1:del:it_scratch;
signal rb_out:it_scratch;
signal gh_out:it_scratch;
signal rb_select:it_round;
signal gh_select:it_mux;

BEGIN

--the code for the convolver--
-- now the state machines to control the convolver--
--First the and gates--

Df1(ck,reset,reset_col);
--starts row counter 1 cycle after frame start--
--makes up for the pipeline delay in MULT--

--LATENCY DEPENDENT!--

col_control <= col_flag;
--flag when col_count=0;1;2;col_length;etc--

andsel(1) <= pass WHEN direction = forward AND count=one ELSE
zero WHEN direction = forward AND count=two ELSE
pass WHEN direction=inverse AND count=two ELSE
zero ;

```

```

5
10
15
20
25
30
35
40
45
50
55

WITH col_control* SELECT
andseel(2) <= zero WHEN count_0 ,
      pass WHEN OTHERS;

andseel(3) <= zero WHEN direction=forward AND col_control = count_0 ELSE
      pass;

--now the add/sub control for the convolver adders--
WITH count SELECT
addseel <= (add,add,add,subt) WHEN one ,
      (add,subt,add,add) WHEN two ;

--now the mux control--
centermuxseel <= (left,right) WHEN (direction = forward AND count = one) OR (direction = inverse AND count = two) ELSE
      (right,left);

--the addmuxseel signal--
muxandseel(1 to 2) <= (pass, andseel(2)) WHEN direction = inverse ELSE
      (andseel(2),pass) ;
muxandseel(3) <= zero WHEN direction = inverse AND col_control=count_1 ELSE
      pass WHEN direction = inverse ELSE
      andseel(2);

muxseel(1) <= dos WHEN direction = inverse AND col_control=count_0 ELSE
      quatro WHEN direction = inverse AND col_control=count_1 ELSE
      tres WHEN direction = inverse AND col_control= count_lm1 ELSE
      dos WHEN direction = inverse ELSE
      uno;
muxseel(2) <= tres WHEN direction = inverse AND col_control=count_0 ELSE
      dos WHEN direction = inverse AND col_control= count_carry ELSE
      uno WHEN direction = inverse ELSE
      dos WHEN direction = forward AND col_control=count_0 ELSE
      tres WHEN direction = forward AND col_control= count_carry ELSE
      uno;
muxseel(3) <= uno WHEN direction = inverse ELSE

```

```

5
10
15
20
25
30
35
40
45
50
55

      tres WHEN direction = forward AND col_control=count_0 ELSE
      quatro WHEN direction = forward AND col_control= count_carry ELSE
      do;

COUNT_MAP: U_COUNT_2 PORT MAP(ck,reset_col, count);

--join the control signals to the mult_add block--
MULT_ADD_MAP: U_MULT_ADD PORT
MAP(reset,in_in,andael,centermuxsel,muxsel,muxandsel,addsel,direction,pdel,mult_add);

--pipeline delay for mult-add,unit--

DPI(ck,mult_add(1),pdel(1));
DPI(ck,mult_add(2),pdel(2));
DPI(ck,mult_add(3),pdel(3));
DPI(ck,mult_add(4),pdel(4));

gh_select <= left WHEN (direction = inverse AND count =one) OR (direction = forward AND count =two) ELSE
right;

DPI(ck,pdel(1), pdel1_del);

gh_out <= MUX_2(pdel(4),pdel1_del,gh_select);

rb_select <= shift3 WHEN direction = inverse AND (col_control=count_2 OR col_control=count_3) ELSE
shift4 WHEN direction = inverse ELSE
shift5;

RB_MAP: U_ROUND_BITS PORT MAP(gh_out,rb_select,rb_out);
--architecture outputs--
out_1 <= rb_out;

END behave;

CONFIGURATION CONV_ROW_CON OF U_CONV_ROW is
FOR behave
FOR ALL:U_ROUND_BITS USE ENTITY WORK.U_ROUND_BITS(behavior);

```

```

5
10
15
20
25
30
35
40
45
50
55

    END FOR;
    FOR ALL:U_COUNT_2
        USE ENTITY WORK.U_COUNT_2(behave);
    END FOR;
    FOR ALL:U_MULT_ADD
        USE ENTITY WORK.U_MULT_ADD(behave);
    END FOR;
END FOR;
END CONV_ROW_CON;
--The basic toggle flip-flop plus and gate for a synchronous counter
--input t is the toggle ,outputs are q and tc (toggle for next counter)
--stage
-- reset is synchronous, is active on final count
use work.DWT_TYPES.all;
use work.dff_package.all;

entity BASIC_COUNT is
    PORT(
        ck:in bit ;reset:in t_reset;en:in bit;q:out bit;carry:out bit);
    end BASIC_COUNT;

architecture behave OF BASIC_COUNT is
    signal dlat:bit;
    signal in_dff:bit;
    signal reset_bit:bit;
BEGIN
    WITH reset SELECT
        reset_bit <= '0' WHEN rst,
        '1' WHEN no_rst;
    in_dff<=(dlat XOR en) AND reset_bit;
    DFF(ck,in_dff,dlat);
    carry<=dlat AND en;
    q<=dlat;
END behave;

configuration basic_count_con of basic_count is

```

```

5
10
15
20
25
30
35
40
45
50
55

    FOR behave
    END for;
    and basic_count;
    -- The n-bit macro counter generator, en is the enable, the outputs
    --are mab(bit 1).....lab,carry.This is the same order as ELLA strings are stored#

    use work.DWT_TYPES.all;

    entity COUNT_SYNC is
    GENERIC (n:integer);
    PORT(
        ck:in bit ;
        reset:in t_reset;
        en:in bit;
        q:out bit_vector(1 to n);
        carry:out bit);
    end COUNT_SYNC;

    architecture behave OF COUNT_SYNC is

    COMPONENT basic_count
    PORT(
        ck:in bit ,reset:in t_reset,en:in bit;q:out bit;carry:out bit);
    end COMPONENT;

    signal enable:bit_vector(1 to n+1);
    BEGIN
        enable(n+1)<=en;
        cl: for i in n downto 1 generate
            bc:  basic_count PORT MAP(ck,reset,enable(i+1),q(i),enable(i));
            and generate;
        carry<=enable(1);
        end behave;

    --configuration for simulation

```

```

5
10
15
20
25
30
35
40
45
50
55

CONFIGURATION COUNT_SYNC_CON OF COUNT_SYNC IS
FOR behave
  FOR ALL:basic_count USE ENTITY WORK.basic_count(behave);
  END FOR;
END FOR;
END COUNT_SYNC_CON;

use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;

entity U_COL_COUNT IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  octave_cnt_length : in BIT_VECTOR(1 to xsize) ;

  out_1 : out t_col;
  out_2 : out t_count_control;
  --count value , and flag for count=0,1,2,col_length-1, col_length
  end U_COL_COUNT;

architecture behave OF U_COL_COUNT IS
COMPONENT COUNT_SYNC
  GENERIC (n:integer);
  PORT(
    ck:in bit ;
    reset:in t_reset;
    en:in bit;
    q:out bit_vector(1 to n);
    carry:out bit);
  end COMPONENT;

```



```

5
10
15
20
25
30
35
40
45
50
55

signal count_control:t_count_control;
signal count_reset:t_reset;
signal count_flag:bit;
signal all_one:bit;
signal count_str:BIT_VECTOR(1 to xsize);
signal countit_col,

BEGIN
count <= U_TO_I(count_str);

count_control <= count_0 WHEN count= 0 ELSE
count_1 WHEN count = 1 ELSE
count_2 WHEN count = 2 ELSE
count_3 WHEN count = 3 ELSE
count_lml WHEN count = (U_TO_I(octave_cnt_length) -1) ELSE
count_carry WHEN count = U_TO_I(octave_cnt_length) ELSE
count_rst;

count_reset <= rst WHEN reset =rst ELSE
rst WHEN count_control = count_carry ELSE
no_rst;

all_one <= '1';

count_map:COUNT_SYNC_GENERIC MAP(size) PORT MAP(ct,count_reset,all_one,count_str,count_flag);--count always enabled

--architecture outputs--
out_1 <= count;
out_2 <= count_control;
END behave;

CONFIGURATION COL_COUNT_CON OF U_COL_COUNT IS
FOR behave
FOR ALL:COUNT_SYNC USE CONFIGURATION WORK.count_sync_con;
END FOR;
END FOR;
END COL_COUNT_CON;

```

```

5
10
15
20
25
30
35
40
45
50
55

use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;

entity U_ROW_COUNT is
PORT(
  ck : in bit ;
  reset : in t_reset ;
  octave_cnt_length : in BIT_VECTOR(1 to ysize) ;
  col_carry: in t_count_control;

  out_1 : out t_row;
  out_2 : out t_count_control);
  --count value , and flag for count=0,1,2,row_length-1, row_length
end U_ROW_COUNT;

architecture behave of U_ROW_COUNT is
  COMPONENT COUNT_SYNC
  GENERIC (n:integer);
  PORT(
    ck:in bit ;
    reset:in t_reset;
    en:in bit;
    q:out bit_vector(1 to n);
    carry:out bit);
  end COMPONENT;

  signal count_control:t_count_control;
  signal count_reset:t_reset;
  signal count_flag:bit;
  signal count_en:bit;
  signal count_str:BIT_VECTOR(1 to ysize);
  signal count:t_row;

```

```

5
10
15
20
25
30
35
40
45
50
55

BEGIN
count <= U_TO_I(count_str);

--count_control <= count_0 WHEN reset= rst ELSE
count_control <= count_0 WHEN count= 0 ELSE
count_1 WHEN count = 1 ELSE
count_2 WHEN count = 2 ELSE
count_3 WHEN count = 3 ELSE
count_lm1 WHEN count = (U_TO_I(octave_cnt_length) -1) ELSE
count_carry WHEN count = U_TO_I(octave_cnt_length) ELSE
count_rst;

count_reset <= rst WHEN reset =rst ELSE
rst WHEN count_control = count_carry AND col_carry = count_carry ELSE
no_rst;

count_en <= '1' WHEN col_carry = count_carry ELSE '0';

count_map:COUNT_SYNC GENERIC MAP(size) PORT MAP(ct,count_en,count_ar,count_flag);--count always enabled

--architecture outputs--
out_1 <= count;
out_2 <= count_control;
END behave;

CONFIGURATION ROW_COUNT_CON OF U_ROW_COUNT IS
FOR behave
FOR ALL:COUNT_SYNC USE CONFIGURATION WORK.count_sync_con;
END FOR;
END FOR;
END ROW_COUNT_CON;
-- create the rising edge function, and a model of a active high DFF.

use work.DWT_TYPES.all;
use work.utils.all;

package dff_package is

```

```

5
10
15
20
25
30
35
40
45
50
55

FUNCTION rising_edge (SIGNAL a:bit) return bool;

PROCEDURE DP1(
    SIGNAL ck:in bit;SIGNAL di:in integer;SIGNAL q:out integer);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in t_state;SIGNAL q:out t_state);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in t_count_control;SIGNAL q:out t_count_control);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in t_count_2;SIGNAL q:out t_count_2);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in t_reset;SIGNAL q:out t_reset);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in bit;SIGNAL q:out bit);

    PROCEDURE DP1(
        SIGNAL ck:in bit;SIGNAL di:in t_load;SIGNAL q:out t_load);

    PROCEDURE DP1(CONSTANT n:in integer;
        SIGNAL ck:in bit;SIGNAL di:in bit_vector;SIGNAL q:out bit_vector);

    PROCEDURE DPF(
        SIGNAL ck:in bit;reset:in t_reset;SIGNAL di:in integer;SIGNAL q:out integer);

    PROCEDURE DPF(
        SIGNAL ck:in bit;reset:in t_reset;SIGNAL di:in t_reset;SIGNAL q:out t_reset);

    PROCEDURE DPF(
        SIGNAL ck:in bit;reset:in t_reset;SIGNAL di:in t_count_2;SIGNAL q:out t_count_2);

    PROCEDURE DPF(
        SIGNAL ck:in bit;reset:in t_reset;SIGNAL di:in t_count_control;SIGNAL q:out t_count_control);

```

```

5
10
15
20
25
30
35
40
45
50
55

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in bit;SIGNAL q:out bit);

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_load;SIGNAL q:out t_load);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in integer;SIGNAL q:out integer);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_channel;SIGNAL q:out t_channel);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_diff;SIGNAL q:out t_diff);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_mode;SIGNAL q:out t_mode);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in bit;SIGNAL q:out bit);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR);

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low);

PROCEDURE LATCH(
  load:in t_load;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector);

PROCEDURE LATCH(
  load:in t_load;SIGNAL d:in bit;SIGNAL q:out bit);

end dff_package;

```

```

package body dff_package is
  FUNCTION rising_edge (SIGNAL s:bit) return bool is
  BEGIN
    IF (s'event) AND (s='1') AND (s'last_value = '0') THEN return t;
    ELSE return f;
    END IF;
  END rising_edge;

  --THE DFF flip-flops, NO RESET-----
  PROCEDURE DFF1(
    SIGNAL ck:in bit;SIGNAL di:in integer;SIGNAL q:out integer) IS
  BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
  END DFF1;

  PROCEDURE DFF1(CONSTANT n:integer;
    SIGNAL ck:in bit;SIGNAL di:in bit_vector;SIGNAL q:out bit_vector) IS
  BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
  END DFF1;

  PROCEDURE DFF1(
    SIGNAL ck:in bit;SIGNAL di:in t_state;SIGNAL q:out t_state) IS
  BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
  END DFF1;

  PROCEDURE DFF1(
    SIGNAL ck:in bit;SIGNAL di:in t_load;SIGNAL q:out t_load) IS

```

```

5
10
15
20
25
30
35
40
45
50
55

BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END Df1;

PROCEDURE Df1(
  SIGNAL ck:in bit;SIGNAL d:in t_reset;SIGNAL q:out t_reset) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END Df1;

PROCEDURE Df1(
  SIGNAL ck:in bit;SIGNAL d:in t_count_2;SIGNAL q:out t_count_2) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END Df1;

PROCEDURE Df1(
  SIGNAL ck:in bit;SIGNAL d:in bit;SIGNAL q:out bit) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END Df1;

PROCEDURE Df1(
  SIGNAL ck:in bit;SIGNAL d:in t_count_control;SIGNAL q:out t_count_control) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;

```

```

5
10
15
20
25
30
35
40
45
50
55

END IF;
END DFF1;

--THE DFF flip-flops, with RESET-----
PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in integer;SIGNAL q:out integer) IS
BEGIN
  IF reset=rat THEN q<= 0;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
  --IF(rising_edge(ck) = t ) THEN IF reset=rat THEN q<= 0; ELSE q<=d;END IF;
  ELSE null;
  END IF;
END DFF;

```

```

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_reset;SIGNAL q:out t_reset) IS
BEGIN
  IF reset=rat THEN q<= rat;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END DFF;

```

```

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in bit;SIGNAL q:out bit) IS
BEGIN
  IF reset=rat THEN q<= '0';
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END DFF;

PROCEDURE DFF(

```



```

5
10
15
20
25
30
35
40
45
50
55

SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_load;SIGNAL q:out t_load) IS
BEGIN
  IF reset=rat THEN q<= read;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
  END IF;
END DFF;

PROCEDURE DFF(
SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_count_2;SIGNAL q:out t_count_2) IS
BEGIN
  IF reset=rat THEN q<= one;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
  END IF;
END DFF;

PROCEDURE DFF(
SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_count_control;SIGNAL q:out t_count_control) IS
BEGIN
  IF reset=rat THEN q<= count_0;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
  END IF;
END DFF;

--- THE DFF_INIT FLIP-FLOPS

PROCEDURE DFF_INIT(
SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in Integer;SIGNAL q:out Integer) IS
BEGIN
  IF reset=rat THEN q<= 0;
  ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
  END IF;
END IF;
END DFF_INIT;

```

5

10

15

20

25

30

35

40

45

50

55

```

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in bit;SIGNAL q:out bit) IS
BEGIN
  IF reset=rat THEN q<= '0';
  ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
  END IF;
  END DFF_INIT;

```

```

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low) IS
BEGIN
  IF reset=rat THEN q<= low;
  ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
  END IF;
  END DFF_INIT;

```

```

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_channel;SIGNAL q:out t_channel) IS
BEGIN
  IF reset=rat THEN q<= y;
  ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
  END IF;
  END DFF_INIT;

```

```

PROCEDURE DFF_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_mode;SIGNAL q:out t_mode) IS
BEGIN
  IF reset=rat THEN q<= still;
  ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;

```

```

5
10
15
20
25
30
35
40
45
50
55

    ELSE null;
      END IF;
    END IF;
  END DFF_INIT;

  PROCEDURE DFF_INIT(
    SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_diff;SIGNAL q:out t_diff) IS
  BEGIN
    IF reset=rst THEN q<= nodiff;
    ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
      ELSE null;
    END IF;
  END IF;
  END DFF_INIT;

  PROCEDURE DFF_INIT(
    SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR) IS
  BEGIN
    IF reset=rst THEN q<= ZERO(d'length);
    ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
      ELSE null;
    END IF;
  END IF;
  END DFF_INIT;

  PROCEDURE LATCH(
    load:in t_load;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector) IS
  BEGIN
    IF load=write THEN q<=d;
      ELSE null;
    END IF;
  END LATCH;

  PROCEDURE LATCH(
    load:in t_load;SIGNAL d:in bit;SIGNAL q:out bit) IS
  BEGIN

```

```

5
10
15
20
25
30
35
40
45
50
55

IF load=write THEN qc=d;
ELSE null;
END IF;
END LATCH;

end behave;

END dff_package;
--the discrete wavelet transform multi-octave/2d transform with edge compensation--
--when ext & csl are both low latch the setup params from the nubs(active low), as follows--
--adl[1 to 4] select function--
-- 0000 load max_octaves,colour,inversebar--
-- 0001 load yimage--
-- 0010 load ximage--
--jump table values--
-- 0011 load ximage+1--
-- 0100 load 3ximage+3--
-- 0101 load 7ximage+7--
-- 0110 load base u addr--
-- 0111 load base v addr--
--adl[21 to 22] max_octaves--
--adl[23] luminance/crominancebar active low, 0 is luminance, 1 is colour--
--adl[24] forward/inversebar active low, 0 is forward, 1 is inverse--
--adl[5 to 24] data (bit 24 lsb)--

use WORK.dwt_types.all;
use WORK.utile.all;
use WORK.utile_dwt.all;
use WORK.dff_package.all;

entity U_DWT is
PORT(
  ck : in bit ;
  reset : in t_reset ;
  in_in : in t_input ;
  extwritel,csl: in bit ;

```

```

5
10
15
20
25
30
    ad1 : in BIT_VECTOR(1 to 24) ;
    mem : in t_input ;
    pdel_in : in t_scratch_array(1 to 4);

    out_1 : out t_input;

    out_2 : out t_load_array(1 to 3);

    out_3 : out t_load_array(1 to 3);

    out_4_1 : out t_memory_addr;  -- memory port
    out_4_2 : out t_memory_addr;
    out_4_3 : out t_load;

    out_5_1 : out t_scratch_array(1 to 4);  --line delay port
    out_5_2 : out t_col;
    out_5_3 : out t_col;
    end U_DWT;

```

architecture behave of U_DWT is

COMPONENT JKPF

PORT(

ck : in bit ;

reset : in t_reset ;

j in bit;

out_l:out bit);

end COMPONENT;

COMPONENT U_CONV_2D

PORT(

ck : in bit ;

reset : in t_reset ;

in_in : in t_input ;

direction : in t_direction ;

pdcl : in t_scratch_array(1 to 4);

```

conv_reset : in t_reset ;
row_flag : in t_count_control ;
addr_col_read_1 : in t_col ;
addr_col_read_2 : in t_count_control ;

```

```

out_1 : out t_input ;
out_2_1 : out t_scratch_array(1 to 4) ;
out_2_2 : out t_col ;
out_2_3 : out t_col ;
out_3 : out t_count_control ;
out_4 : out t_count_control ;
out_5 : out t_count_control ;
end COMPONENT ;

```

```

COMPONENT U_ADDR_GEN
PORT(

```

```

    ck : in bit ;
    reset : in t_reset ;
    direction : in t_direction ;
    channel : in t_channel ;
    x_p_1 : in BIT_VECTOR(1 to 10) ;
    x3_p_1 : in BIT_VECTOR(1 to 12) ;
    x7_p_1 : in BIT_VECTOR(1 to 13) ;
    octave_row_length : in BIT_VECTOR (1 to ysize) ;
    octave_col_length : in BIT_VECTOR (1 to xsize) ;
    octave_reset : in t_reset ;
    octave : in t_octave ;
    y_done : in bit ;
    uv_done : in bit ;
    octave_finished : in t_load ;
    base_u_base_v : in BIT_VECTOR(1 to 19) ;

```

```

out_1 : out t_input_mux;      --input data from memory/external

out_2_1 : out t_memory_addr;  -- memory port
out_2_2 : out t_memory_addr;
out_2_3 : out t_load;

```

```

5
10
15
20
25
30
35
40
45
50
55

out_3_1 : out_t_load; --dwt in control
out_3_2 : out_t_cs;

out_4 : out_t_load; --IDWT data valid
out_5 : out_t_load; --read valid
out_6 : out_t_count_control; --row read

out_7_1 : out_t_col;
out_7_2 : out_t_count_control;
end COMPONENT;

signal max_oct:t_octave;
signal max_oct_str:BIT_VECTOR(1 to 2);
signal col_length:BIT_VECTOR(1 to 10);
signal row_length:BIT_VECTOR(1 to 9);
signal channel_factor_st:BIT;
signal channel_factor:t_channel_factor;
signal direction:t_direction;
signal dir:bit;
signal convcol_row:t_count_control;
signal convcol_col:t_count_control;
signal convrow_col:t_count_control;
signal conv_2d_1:t_input;
signal conv_2d_2_1:t_scratch_array(1 to 4);
signal conv_2d_2_2:t_col;
signal conv_2d_2_3:t_col;
signal conv_2d_3:t_count_control;
signal conv_2d_4:t_count_control;
signal conv_2d_5:t_count_control;

signal octave:t_octave;
signal channel:t_channel;
signal octave_finished:t_load;
signal load_octave:t_load;
signal max_oct_1:t_octave;
signal y_done:bit;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal uv_done1bit;
signal all_one1bit;

signal octave_sel1:t_mux4;
signal octave_row_length:BIT_VECTOR(1 to ysize);

signal conv_reset:t_reset;
signal octave_col_length:BIT_VECTOR(1 to xsize);

signal input_mux:t_input_mux;
signal addr_gen_1:t_input_mux;
signal addr_gen_2_1:t_memory_addr;
signal addr_gen_2_2:t_memory_addr;
signal addr_gen_2_3:t_load;
signal addr_gen_3_1:t_load;
signal addr_gen_3_2:t_cs;
signal addr_gen_4:t_load;
signal addr_gen_5:t_load;
signal addr_gen_6:t_count_control;
signal addr_gen_7_1:t_col;
signal addr_gen_7_2:t_count_control;
signal mem_rw:t_load;
signal mem_r:t_memory_addr;
signal mem_w:t_memory_addr;

signal q1:bit;
signal inverse_out:t_load_array(1 to 3);
signal forward_in:t_load_array(1 to 3);

signal decode_int:natural;
signal decode:BIT_VECTOR(1 to 8);
signal x_p_1:BIT_VECTOR(1 to 10);
signal x3_p_1:BIT_VECTOR(1 to 12);
signal x7_p_1:BIT_VECTOR(1 to 13);
signal base_u:BIT_VECTOR(1 to 19);
signal base_v:BIT_VECTOR(1 to 19);
signal ad14_2:BIT_VECTOR(1 to 3);

```



```

5
10
15
20
25
30
35
40
45
50
55

signal load_regs:BIT_VECTOR(1 to 8);
signal conv_init_input;
signal row_bitbit;
signal row_carry_ff:bit;
signal initial_octave:it_octave;
signal initial_channel:it_channel;
signal max_octave:it:BIT_VECTOR(1 to 2);

BEGIN
--must delay the write control to match the data output of conv_2d, ie by conv2d_latency--

--set up the control params--
max_oct <= U_TO_I(max_octave_et);

WITH channel_factor_et SELECT
channel_factor <= luminance WHEN '0',
color WHEN '1';

WITH dir SELECT
direction <= forward WHEN '0',
inverse WHEN '1';

--set up the octave params--
convcol_row <= conv_2d_3;
convcol_col <= conv_2d_4;
convrow_col <= conv_2d_5;
--signals that conv_col, for forward, or conv_row, for inverse, has finished that octave--
--and selects the next octave value and the sub-image sizes--
--row then col, gives write latency
octave_finished <= write WHEN direction = forward AND row_carry_ff = '1' AND convcol_row = count_2 AND convcol_col = count_2
ELSE
--extra row as col then row
write WHEN direction = inverse AND row_carry_ff = '1' AND convrow_row = count_2 AND convrow_col = count_3
ELSE

```

```

5
10
15
20
25
30
35
40
45
50
55

      read;

--max octaves for u|v--
WITH max_oct SELECT
max_oct_1 <= 0 WHEN
1 WHEN
2 WHEN
3 ;

y_done <= '1' WHEN channel = y AND direction = forward AND octave = max_oct
ELSE
'1' WHEN channel = y AND direction = inverse AND octave = 0 ELSE
'0';

uv_done <= '1' WHEN channel = u AND direction = forward AND octave = max_oct_1 ELSE
'1' WHEN channel = v AND direction = forward AND octave = max_oct_1 ELSE
'1' WHEN channel = u AND direction = inverse AND octave = 0 ELSE
'1' WHEN channel = v AND direction = inverse AND octave = 0 ELSE
'0';

PROCESS(octave,channel,ck,load_octave)
variable new_oct it_octave;
variable new_channel it_channel;
BEGIN
new_oct := octave;
new_channel := channel;

-- first describe the progression of the octaves for a max_oct decomposition
CASE direction IS
WHEN forward => CASE octave IS
WHEN 0 => new_oct := 1;
WHEN 1 => new_oct := 2 ;
WHEN 2|3 => new_oct := 3 ;
END CASE;
IF y_done = '1' OR uv_done = '1' THEN new_oct := 0; ELSE null;
END IF;
WHEN inverse => CASE octave IS

```

```

5
10
15
20
25
30
35
40
45
50
55

WHEN 3 => new_oct := 2;
WHEN 2 => new_oct := 1;
WHEN 1|0 => new_oct := 0;
END CASE;

CASE channel IS
  WHEN y => CASE octave IS
    WHEN 0 => CASE channel_factor IS
      WHEN luminance => new_oct := max_oct; --watch for colour
      WHEN OTHERS => new_oct := max_oct_1;
    END CASE;
    WHEN OTHERS => null;
  END CASE;
  WHEN u
    => CASE octave IS
      WHEN 0 => new_oct := max_oct_1;
      WHEN OTHERS => null;
    END CASE;
  WHEN v
    => CASE octave IS
      WHEN 0 => new_oct := max_oct; --move to y
      WHEN OTHERS => null;
    END CASE;
  END CASE;

END CASE;
--the progression of channels is first y then u then v
CASE channel_factor IS
  WHEN luminance => new_channel := y;
  WHEN color => IF channel = y AND y_done = '1' THEN new_channel := u; ELSE null; END IF;

  IF channel = u AND uv_done = '1' THEN new_channel := v;
  ELSEIF channel = v AND uv_done = '1' THEN new_channel := y;
  ELSE null;
  END IF;

END CASE;

-- set initial values for octave and channel after reset
CASE reset IS
  WHEN no_reset => initial_octave := new_oct;

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN rat => IF direction = forward THEN initial_octave:=0;
        ELSIF direction = inverse AND channel=y THEN initial_octave:=max_oct;
        ELSIF direction = inverse AND (channel=u OR channel=v) THEN initial_octave:=max_oct_1;
        END IF;
    END CASE;

CASE reset IS
    WHEN no_rat => initial_channel:=new_channel;
    WHEN rat => initial_channel:=y;
    END CASE;

-- the DPF's for the state machine

DPF_INIT(ck,no_rat,load_octave,initial_octave,octave);
DPF_INIT(ck,no_rat,load_octave,initial_channel,channel);

END PROCESS;

--the block size divides by 2 every octave--
--the u,v image starts 1/4 size--
octave_sel <= uno WHEN octave =0 AND channel= y ELSE
    dos WHEN (octave =1 AND channel= y) OR (octave =0 AND (channel= u OR channel =v)) ELSE
    tres WHEN (octave =2 AND channel= y) OR (octave =1 AND (channel= u OR channel =v)) ELSE
    quatro ;

WITH octave_sel SELECT
    octave_row_length <= row_length WHEN uno ,
        B"0" & row_length(1 to ysize-1) WHEN dos,
        B"00" & row_length(1 to ysize-2) WHEN tres,
        B"000" & row_length(1 to ysize-3) WHEN quatro;

WITH octave_sel SELECT
    octave_col_length <= col_length WHEN uno ,
        B"0" & col_length(1 to xsize-1) WHEN dos,
        B"00" & col_length(1 to xsize-2) WHEN tres,
        B"000" & col_length(1 to xsize-3) WHEN quatro;

```

```

5
10
15
20
25
30
35
40
45
50
55

--load next octave, either on system reset, or write finished--
WITH reset SELECT
load_octave <=
  write WHEN rst,
  octave_finished WHEN OTHERS;

--reset the convolvers at the end of an octave, ready for the next octave--
--latch pulses to clean it, note 2 reset pulses at frame start--
--FOR SYNC RESET DONT NEED TO LATCH PULSE
--cant glitch as resetoctave_finished dont change at similar times--

conv_reset <= rst WHEN reset = rst ELSE
  rst WHEN octave_finished = write ELSE
  no_rst;

--latch control data off nubus
gl <= '1' WHEN extwritel = '1' AND csl = '1'
  ELSE '0';

mem_w <= addr_gen_2_1; --write addresses--
mem_r <= addr_gen_2_2; --read addresses--
mem_rw <= addr_gen_2_3;

inverse_out <= (write,read,read) WHEN direction=inverse AND octave=0 AND channel=y AND addr_gen_4=write ELSE
  (read,write,read) WHEN direction=inverse AND octave=0 AND channel=u AND addr_gen_4=write ELSE
  (read,read,write) WHEN direction=inverse AND octave=0 AND channel=v AND addr_gen_4=write ELSE
  (read,read,read);

forward_in <= (read,write,write) WHEN direction=forward AND octave=0 AND channel=y AND addr_gen_5=read ELSE
  (write,read,write) WHEN direction=forward AND octave=0 AND channel=u AND addr_gen_5=read ELSE
  (write,write,read) WHEN direction=forward AND octave=0 AND channel=v AND addr_gen_5=read ELSE
  (write,write,write);

--the control section latch values when read from the NUBUS
--a 3x8 decoder, active high outputs selects the load signal for the appropriate register

```

```

5
10
15
20
25
30
35
40
45
50
55

adl4_2 <= (adl(2),adl(3),adl(4));

decode_int <= '0' 1 WHEN adl4_2 = B"000" ELSE
                  2 WHEN adl4_2 = B"001" ELSE
                  4 WHEN adl4_2 = B"010" ELSE
                  8 WHEN adl4_2 = B"011" ELSE
                  16 WHEN adl4_2 = B"100" ELSE
                  32 WHEN adl4_2 = B"101" ELSE
                  64 WHEN adl4_2 = B"110" ELSE
                  128 ;

I_TO_S(decode_int, decode);

load_regs <= ALL_SAME(8,gl) AND decode;

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(8)),adl(21 to 22),max_octave_st);

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(8)),adl(23),channel_factor_st);
DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(8)),adl(24),dir);

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(7)),adl(15 to 24),col_length);
DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(6)),adl(16 to 24),row_length);
DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(5)),adl(15 to 24),x_p_1);

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(4)),adl(13 to 24),x3_p_1);

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(3)),adl(12 to 24),x7_p_1);

DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(2)),adl(6 to 24),base_u);
DFF_INIT(ck,no_rst,BIT_LOAD(load_regs(1)),adl(6 to 24),base_v);

--sets a flag when row counter moves onto next frame
WITH convcol_row SELECT
row_bit <= '1' WHEN count_carry,
          '0' WHEN OTHERS,

all_one <= '1';

```

```

5
10
15
20
25
30
35
40
45
50
55

log_map:JKFF PORT MAP(ck,conv_reset,row_bit,row_carry_ff);

addr_map:U_ADDRGEN PORT
MAP(ck,reset,direction,channel,x_p_1,x3_p_1,x7_p_1,octave_row_length,octave_col_length,
  conv_reset,octave_y_done,uv_done,octave_finished,base_u,base_v,
  addr_gen_1, addr_gen_2_1, addr_gen_2_2, addr_gen_2_3, addr_gen_3_1, addr_gen_3_2, addr_gen_4,
  addr_gen_5, addr_gen_6, addr_gen_7_1, addr_gen_7_2);

WITH addr_gen_1 SELECT
conv_in <=in_in WHEN dwt_in,
      mem WHEN mem_in;

conv_map:U_CONV_2D PORT MAP(ck,reset,conv_in,direction,pdel_in,
conv_reset,addr_gen_6,addr_gen_7_1,addr_gen_7_2,
      conv_2d_1, conv_2d_2_1, conv_2d_2_2, conv_2d_2_3, conv_2d_3, conv_2d_4, conv_2d_5);

--architecture outputs--
out_1 <= conv_2d_1;
out_2 <= inverse_out;
out_3 <= forward_in;
out_4_1 <= addr_gen_2_1;
out_4_2 <=addr_gen_2_2;
out_4_3 <=addr_gen_2_3;
out_5_1 <=conv_2d_2_1;
out_5_2 <=conv_2d_2_2;
out_5_3 <=conv_2d_2_3;

END;

CONFIGURATION DWT_CON OF U_DWT IS

```

```

5
10
15
20
25
30
35
40
45
50
55

FOR behave
  FOR ALL:U_CONV_2D      USE ENTITY WORK.U_CONV_2D(behav);
  END FOR;
  FOR ALL:U_ADDR_GEN      USE ENTITY WORK.U_ADDR_GEN(behav);
  END FOR;
  FOR ALL:JKFF      USE ENTITY WORK.JKFF(behav);
  END FOR;

  END FOR;
  END DWT_CON;
  package dwt_types is
  --constant values
  constant result_exp:Integer:= 14;
  --length of result arith
  constant input_exp:Integer:= 10;
  --length of 1D convolver input/output--
  constant qmax :Integer:= 7;
  --maximum shift value for quantisation constant--
  constant result_range :Integer:= 2 ** (result_exp-1);
  constant input_range :Integer:= 2 ** (input_exp-1);
  constant max_octave:Integer:= 3;
  --no of octaves:Integer:=max_octave +1; can not be less in this example--
  constant no_octave:Integer:= max_octave+1;
  constant xsize :Integer:= 10;
  --no of bits for ximage--
  constant ysize :Integer:= 9;
  --no of bits for yimage--
  constant ximage:Integer:= 319;
  --the xdimension -1 of the image; ie no of cols--
  constant yimage:Integer:= 239 ;
  --the ydimension -1 of the image; ie no of rows--

  --int types--
  subtype t_result is integer range -result_range to result_range-1;
  subtype t_input is integer range -input_range to input_range-1;
  subtype t_length is integer range 0 to 15;
  subtype t_inp is integer range 0 to 1023;

```



```

subtype t_blk is integer range 0 to 3;
subtype t_sub is integer range 0 to 3;
subtype t_cgd is integer range 0 to ximage;
subtype t_row is integer range 0 to yimage;
subtype t_carry is integer range 0 to 1;
subtype t_quant is integer range 0 to qmax;
--address for resultdwt memory; is 1 frame--
subtype t_memory_addr is integer range 0 to (2 ** max_octave) * (yimage+1) + (ximage+1)) - 1;
subtype t_octave is integer range 0 to max_octave;

--bit string and boolean types types--
type bool is (f,t);
type flag is (error, ok);
--control signals--
type t_reset is (rst,no_rst);
type t_load is (write,read);
type t_load_vec is ARRAY (NATURAL RANGE <>) of t_load;
--r/wbar control--
TYPE t_mem IS (random,old_mem,new_mem);
type t_cs is (no_sel,sel);
type t_updown is (down,up);
--up/down counter control--
type t_diff is (diff,nodiff);
--diff or not in quantiser--
type t_intra is (intra,inter);
--convolver mux & and types--
type t_mux is (left,right);
type t_mux3 is (l,c,r);
type t_mux4 is (uno,doe,tree,quatro);
type t_add is (add,subt);
type t_direction is (forward,inverse);
--counter types--
type t_count_2 is (one,two);
--state types--
type t_mode is (void,void_still,stop,send,still,still_send,lpf_send,lpf_still,lpf_stop);
type t_mode_vec is ARRAY (NATURAL RANGE <>) of t_mode;
type t_cycle is (token_cycle,data_cycle,skip_cycle);

```

```

5
10
15
20
25
30
35
40
45
50
55

type t_state is (up0,up1,zz0,zz1,zz2,zz3,down1);
--type t_state is (start,up0,up1,zz0,zz1,zz2,zz3,down1);
type t_decode is (load_low,load_high);
type t_high_low is (low,high);
type t_fifo is (ok_fifo,error_fifo);
--types for the octave control unit--
type t_channel is (y,u,v);
type t_channel_factor is (luminance,color);
--types for the control of memory ports--
--type t_sparcport (t_sparc_addr,t_sparc_addr,t_load,t_cs);

-- TYPES FOR DWT CHIP

CONSTANT scratch_exp:Integer:=16; --length of scratch arith--
CONSTANT conv2d_latency:Integer:=7; --the 2d convolver latency
constant scratch_range :Integer:= 2 ** (scratch_exp-1);
subtype t_scratch is integer range -scratch_range to scratch_range-1;

type t_scratch_array is array(NATURAL range <>) of t_scratch;

type t_load_array is array(NATURAL range <>) of t_load;

type t_and is (zero,pass);

type t_and_array is array(NATURAL range <>) of t_and;
type t_add_array is array(NATURAL range <>) of t_add;
type t_mux_array is array(NATURAL range <>) of t_mux;
type t_mux4_array is array(NATURAL range <>) of t_mux4;

type t_count_control is (count_0,count_1,count_2,count_3,count_rst,count_carry,count_lm1);
type t_round is (shift3,shift4,shift5);
type t_input_mux is (dwt_in,mem_in);

```

```

5
10
15
20
25
30
35
40
45
50
55

FUNCTION U_TO_I(bits: in bit_vector) RETURN natural;
FUNCTION S_TO_I(bits: in bit_vector) RETURN integer;
PROCEDURE I_TO_S(int: in integer; SIGNAL bits: out bit_vector);
end dwt_types;

package body dwt_types is

FUNCTION U_TO_I(bits: bit_vector) RETURN natural IS
    variable result: natural:=0;
BEGIN
    FOR i IN bits'range LOOP
        result:=result*2 + bit'pos(bits(i));
    END LOOP;
    RETURN result;
END U_TO_I;

FUNCTION S_TO_I(bits: bit_vector) RETURN integer IS
    variable temp: bit_vector(bits'range);
    variable result: integer:=0;
BEGIN
    IF bits(bits'left) = '1' THEN
        temp:=NOT bits;
    ELSE
        temp:=bits;
    END IF;
    FOR i IN bits'range LOOP
        result:=result*2 + bit'pos(temp(i));
    END LOOP;
    IF bits(bits'left) = '1' THEN
        result:=(-result)-1;
    END IF;
    RETURN result;
END S_TO_I;

```

PROCEDURE I_TO_S(int:in Integer; SIGNAL bits:out bit_vector) IS
 variable result:bit_vector(bits'range);
 variable temp: Integer;

BEGIN

IF int < 0 THEN

temp:=-(int+1);

ELSE temp:=int;

END IF;

FOR i IN bits'reverse_range LOOP

result(i):= bit'val(temp rem 2);

temp:=temp/2;

END LOOP;

IF int<0 THEN

result:=NOT result;

result(bits'left):='1';

END IF;

bits<=result;

END I_TO_S;

FUNCTION INT_TO_S(n:natural;SIGNAL int:in Integer) RETURN bit_vector IS

variable result:bit_vector(1 to n);

variable temp: Integer;

BEGIN

IF int < 0 THEN

temp:=-(int+1);

ELSE temp:=int;

END IF;

FOR i IN n downto 1 LOOP

result(i):= bit'val(temp rem 2);

temp:=temp/2;

END LOOP;

-- check to see if Integer fits in n bits

5

10

15

20

25

30

35

40

45

50

55

```

ASSERT (temp=0)
REPORT "Int TO BIG FOR n BITS"
SEVERITY FAILURE;

IF int<0 THEN
    result:=NOT result;
    result(1):='1';
END IF;

RETURN result;
END INT_TO_S;

end dwf_types;
--a model of an ELLA compatible RAM

use work.DWT_TYPES.all;

entity ella_ram is
    PORT(
        in_data:in t_input;
        wr_addr:in t_memory_addr;
        rd_addr:in t_memory_addr;
        rw:in t_load;

        out_data:out t_input;
    end ella_ram;

    architecture behave of ella_ram is
    BEGIN

        ram:process
            type mem is array(natural range <>) of t_input;
            variable memory:mem(0 to 2000);
            --variable memory:mem(0 to (2 ** max_octave)*(ximage+1)*(yimage+1)-1);

```

```

5
10
15
20
25
30
35
40
45
50
55

BEGIN
wait on rw, wr_addr, rd_addr ;
--IF rw'event AND rw = write THEN memory(wr_addr):=in_data ;
IF rw = write THEN memory(wr_addr):=in_data ;
ELSE null;
END IF;

out_data <= memory(rd_addr);
END PROCESS;
END behave;

CONFIGURATION ELLA_RAM_CON OF ELLA_RAM IS
FOR behave
END FOR;
END ELLA_RAM_CON;

-- ram for scratch memories
use work.DWT_TYPES.all;

entity scratch_ram is
PORT(
in_data:in t_scratch;
wr_addr:in t_memory_addr;
rd_addr:in t_memory_addr;
rw:in t_load;

out_data:out t_scratch;
end scratch_ram;

architecture behave of scratch_ram is
BEGIN

ram:process
variable memory:t_scratch_array(0 to 1023);
--variable memory:mem(0 to (2**max_octave)*(ximage+1)*(yimage+1)-1);

```

```

5
10
15
20
25
30
35
40
45
50
55

BEGIN
wait on rw, wr_addr, rd_addr ;
--
--IP rw'event AND rw = write THEN memory(wr_addr):=in_data ;
IF rw = write THEN memory(wr_addr):=in_data ;
ELSE null;
END IF;

out_data <= memory(rd_addr);
END PROCESS;
END behave;

CONFIGURATION SCRATCH_RAM_CON OF SCRATCH_RAM IS
FOR behave
END FOR;
END SCRATCH_RAM_CON;

--the mem control unit for the DWT chip, outputs the memport values for the sparcs, and dwt--
--Inputs datain from these 2 ports and mux's it to the 2d convolver.--
use WORK.dwt_types.all;
use WORK.utils.all;
use WORK.utils_dwt.all;
use WORK.dff_package.all;

entity U_MEM_CONTROL IS
PORT(
ck : in bit ;
reset : in t_reset ;
direction : in t_direction ;
channel : in t_channel ;
octave : in t_octave ;
addr_w,addr_r : in t_memory_addr ;
zero_hh : in t_load ;

out_1 : out t_input_mux;
out_2_1 : out t_memory_addr;

```

```

out_2_2 : out t_memory_addr;
out_2_3 : out t_load;
out_3_1 : out t_load;
out_3_2 : out t_cs;
end U_MEM_CONTROL;

```

architecture behave OF U_MEM_CONTROL IS.

```
BEGIN
```

```
--the comb. logic for the control of the i/o ports of the chip--
```

```
PROCESS(direction,octave,zero_hh)
```

```

variable   rw_sparc it_load;
variable   rw_dwt it_load;
variable   cs_dwt it_cs;
variable   input_mux it_input_mux;
variable   zero_hh_bit;bit;
BEGIN

```

```

    rw_sparc := read;
    rw_dwt := read;
    cs_dwt := no_sel;
    input_mux := mem_in;
    zero_hh_bit:='0';

```

```
IF direction = forward AND octave=0 THEN
```

```

    cs_dwt := sel;
    input_mux:= dwt_in;

```

```

ELSIF direction = inverse AND octave=0 AND zero_hh =write THEN
    rw_dwt := write;
    cs_dwt:= sel;

```

```
ELSE null;
```



```

5
10
15
20
25
30
35
40
45
50
55

END IP;

--rw_sparc = write when ck=1 and zero_hh=write, otherwise = read--
CASE zero_hh IS
  WHEN write => zero_hh_bit:= '1';
  WHEN OTHERS => zero_hh_bit:= '0';
END CASE;

rw_sparc := zero_hh;

out_1 <= input_mux;
out_2_3 <= rw_sparc;

out_3_1 <= rw_dwt;
out_3_2 <= cs_dwt;

END PROCESS;
out_2_1 <=addr_w;
out_2_2 <=addr_r;

END;

CONFIGURATION MEM_CONTROL_CON OF U_MEM_CONTROL IS
FOR behave
END FOR;
END MEM_CONTROL_CON;
-- the basic 1d convolver without the control unit--
use work.DWT_TYPES.all;
use work.utils_dwt.all;
entity U_MULT_ADD IS
PORT(
  reset : in t_reset ;
  in_in : in t_input ;
  andsel : in t_and_array(1 to 3) ;
  centermuxsel : in t_mux_array(1 to 2) ;

```

```

5
10
15
20
25
30
35
40
45
50
55

    muxsel : in t_mux4_array(1 to 3) ;
    muxandsel : in t_and_array(1 to 3) ;
    adtsel : in t_add_array(1 to 4) ;
    direction : in t_direction ;
    pdel : in t_scratch_array(1 to 4) ;

    out_1 : out t_scratch_array(1 to 4) ;

FUNCTION AND_2 (in1:t_scratch;sel:t_and) RETURN t_scratch IS
BEGIN
    CASE sel IS
        WHEN pass => RETURN in1;
        WHEN zero => RETURN 0;
    END CASE;
END;

end U_MULT_ADD;

architecture behave OF U_MULT_ADD IS

    COMPONENT U_MULTIPLIER_ST
    PORT(
        in_in : in t_input ;

        out_1 : out t_scratch_array(1 to 7) ) ;
    end COMPONENT;

    signal x3it_scratch;
    signal x2it_scratch;
    signal x8it_scratch;
    signal x5it_scratch;
    signal x11it_scratch;
    signal x19it_scratch;
    signal x30it_scratch;

    signal multit_scratch_array(1 to 7);

```

```

5
10
15
20
25
30
35
40
45
50
55

signal mux1:t_scratch;
signal mux2:t_scratch;
signal mux3:t_scratch;
signal centermux:t_scratch_array(1 to 2);
signal and1:t_scratch;
signal and2:t_scratch;
signal and3:t_scratch;
signal and4:t_scratch;
signal add1init_scratch;
signal add3init_scratch;
signal add4init_scratch;
signal add_out:t_scratch_array(1 to 4);

BEGIN

--the multiplier outputs--
x3 <= mult(1);
x5 <= mult(2);
x11 <= mult(3);
x19 <= mult(4);
x2 <= mult(5);
x8 <= mult(6);
x30 <= mult(7);

--the mux outputs--
mux1 <= MUX_4(x11,x5,x8,x2,muxsel(1));
mux2 <= MUX_4(x19,x30,x8,0,muxsel(2));
mux3 <= MUX_4(x11,x5,x8,x2,muxsel(3));
centermux <= (MUX_2(pdcl(1),pdcl(3),centermuxsel(1)),
               MUX_2(pdcl(2),pdcl(4),centermuxsel(2)) );

-- the AND gates zero the adder inputs every 2nd row--
--the and gate outputs--
and1 <= AND_2(pdcl(2),andsel(1));

```

```

5
10
15
20
25
30
35
40
45
50
55

and2 <= AND_2(pdol(3), andsel(1));
and3 <= AND_2(centermux(1), andsel(2));
and4 <= AND_2(centermux(2), andsel(3));

add1in <= AND_2(mux1, muxandsel(1));
add3in <= AND_2(mux3, muxandsel(2));
add4in <= AND_2(x3, muxandsel(3));

MULT_MAP: U_MULTIPLIER_ST PORT MAP(in_in, mult);

add_out(1) <= ADD_SUB(and1, add1in, addsel(1));
add_out(2) <= ADD_SUB(and3, mux2, addsel(2));
add_out(3) <= ADD_SUB(and4, add3in, addsel(3));
add_out(4) <= ADD_SUB(and2, add4in, addsel(4));

--architecture outputs--
out_1 <= add_out;

END;

CONFIGURATION MULT_ADD_CON OF U_MULT_ADD IS
FOR behave
    FOR ALL U_MULTIPLIER_ST USE ENTITY WORK.U_MULTIPLIER_ST (behave);
    END FOR;
END FOR;

END MULT_ADD_CON;

-- the basic multiplier unit of the convolver --
use WORK.dwt_types.all;
entity U_MULTIPLIER_ST IS
PORT(
    in_in : in t_input ;

    out_1 : out t_scratch_array(1 to 7) ;
    end U_MULTIPLIER_ST;

architecture behave OF U_MULTIPLIER_ST IS
    signal in_s:BIT_VECTOR(1 to input_exp);

```

```

5
10
15
20
25
30
35
40
45
50
55

signal x2_st:BIT_VECTOR(1 to input_exp+1);
signal x8_st:BIT_VECTOR(1 to input_exp+3);
signal x4_st:BIT_VECTOR(1 to input_exp+2);
signal x16_st:BIT_VECTOR(1 to input_exp+4);
signal x2_it_scratch:=0;
signal x3_it_scratch:=0;
signal x5_it_scratch:=0;
signal x8_it_scratch:=0;
signal x11_it_scratch:=0;
signal x19_it_scratch:=0;
signal x30_it_scratch:=0;
BEGIN
--the multiplier outputs, fast adder code commented out--
I_TO_S(ln_in,ln_e);
x2_st <= ln_e & B"0";
x2 <= S_TO_I(x2_st);
x8_st <= ln_e & B"000";
x8 <= S_TO_I(x8_st);
x3 <= ln_in + x2;
x4_st <= ln_e & B"00";
x5 <= ln_in + S_TO_I(x4_st);
x11 <= x3 + S_TO_I(x8_st);
x16_st <= ln_e & B"0000";
x19 <= x3 + S_TO_I(x16_st);
x30 <= x11 + x19;
--architecture outputs--
out_1 <= ( x3,x5,x11,x19,x2,x8,x30);

```

```

5
10
15
20
25
30
35
40
45
50
55

END;

CONFIGURATION MULTIPLIER_ST_CON OF U_MULTIPLIER_ST IS
FOR behave
END FOR;
END MULTIPLIER_ST_CON;
use WORK.dwt_types.all;
entity U_ROUND_BITS IS
PORT(
    in_in : in t_scratch ;
    sel : in t_round ;
    out_1 : out t_input);
end U_ROUND_BITS;

architecture behave OF U_ROUND_BITS IS
    signal sel : BIT_VECTOR(1 to scratch_exp);
    signal shift : BIT_VECTOR(1 to scratch_exp);
    signal meb : BIT;
    signal cs : BIT;
    signal cs_int : t_carry;
    signal sel1 : BIT;
    signal sum17 : integer;
    signal sum17_str : BIT_VECTOR(1 to scratch_exp+1);
    signal sum : BIT_VECTOR(1 to scratch_exp);
    signal out_final : BIT_VECTOR(1 to input_exp);
BEGIN

--THIS ASSUMES THAT THE INPUT_EXP = 10!!!!--
--sel chooses a round factor of 3, 4, 5--
--the leb is the right hand of the string,--
--the index 1 of the string is the left hand end, tie the meb--
--so on add ops bit 1 is the carryout--

I_TO_S(in_in, sel);

meb <= sel(1);

```

```

--needs to be a 16 bit output for the adder--
WITH sel SELECT
  shift <=
    mab & mab & mab & mab & s1(1 to scratch_exp-3) WHEN shift3,
    mab & mab & mab & mab & s1(1 to (scratch_exp-4)) WHEN shift4,
    mab & mab & mab & mab & s1(1 to scratch_exp-5) WHEN shift5,
    --the carry to round, 1/2 value is rounded towards 0--
    '0' WHEN sel=shift4 AND mab='0' AND s1(scratch_exp-3 to scratch_exp)=b"1000" ELSE --round down on
    1/2 value
    s1(scratch_exp-3) WHEN sel=shift4 ELSE -- neg. no
    '0' WHEN sel=shift3 AND mab='0' AND s1(scratch_exp-2 to scratch_exp) = b"100" ELSE
    s1(scratch_exp-2) WHEN sel=shift3 ELSE
    '0' WHEN sel=shift5 AND mab='0' AND s1(scratch_exp-4 to scratch_exp)= b"10000" ELSE
    s1(scratch_exp-4);
ce_int <= 1 WHEN ce='1' ELSE
  0;
sum17 <= ce_int + s_TO_I(shift);
I_TO_S(sum17,sum17_str);
sum <= sum17_str(2 to scratch_exp+1);
-- 1 signifies the rounded value is in range, 0 that it must be saturated
--these are the 5 mab's from the 13 bit word
sel <= '1' WHEN sel=shift3 AND (sum(4 to 7) = b"1111" OR sum(4 to 7) = b"0000") ELSE --value in range
  '0' WHEN sel=shift3 ELSE
  --these are the 3 mab's from the 12 bit word left after taking out the 4 sign extension bits
  '1' WHEN sel=shift4 AND (sum(5 to 7) = b"111" OR sum(5 to 7) = b"000") ELSE --value in range
  '0' WHEN sel=shift4 ELSE
  --these are the 2 mab's from the 11 bit word
  '1' WHEN sel=shift5 AND (sum(6 to 7) = b"11" OR sum(6 to 7) = b"00") ELSE --value in range

```

```

5
10
15
20
25
30
35
40
45
50
55

'0';

out_final <= b"01111111" WHEN ssel = '0' AND sum(1) = '0' ELSE -- saturate to 511
b"1000000001" WHEN ssel = '0' AND sum(1) = '1' ELSE -- saturate to -511 SEE QUANT FOR REASON
sum(7 to scratch_exp);

--architecture outputs--
out_1 <= S_TO_1(out_final);

END behave;

CONFIGURATION ROUND_BITS_CON OF U_ROUND_BITS IS
FOR behave
END FOR;
END ROUND_BITS_CON;
use work.DWT_TYPES.all;

-- returns a signal with n copies of the zero
package utils is

FUNCTION ZERO ( CONSTANT n:NATURAL) RETURN BIT_VECTOR;

-- returns a signal with n copies of the input bit
FUNCTION ALL_SAME (CONSTANT n:NATURAL; s:bit) RETURN BIT_VECTOR;

-- reverses the bit order
FUNCTION REV (CONSTANT n:natural;in_in:BIT_VECTOR) RETURN BIT_VECTOR;

end utils;

package body utils is
FUNCTION ALL_SAME (CONSTANT n:NATURAL; s:bit) RETURN BIT_VECTOR IS
variable out_b:BIT_VECTOR(1 to n);
BEGIN
for i IN 1 to n LOOP
out_b(i):= s;

```



```

5
10
15
20
25
30
35
40
45
50
55

END LOOP;
RETURN out_b;
END ALL_SAME; %

FUNCTION ZERO (CONSTANT n:natural) RETURN BIT_VECTOR IS
variable out_b:BIT_VECTOR(1 to n);
BEGIN
for i IN 1 to n LOOP
out_b(i):='0';
END LOOP;
RETURN out_b;
END ZERO;

FUNCTION REV (CONSTANT n:natural;in_in:BIT_VECTOR) RETURN BIT_VECTOR IS
variable temp:BIT_VECTOR(1 to n);
BEGIN
for i IN 1 to n LOOP
temp(i):=in_in(n-i+1);
END LOOP;
RETURN temp;
END;

END utils;

use work.dwt_types.all;
use work.utils.all;
use work.dff_package.all;

-- returns a signal with n copies of the zero
package utils_dwt is
FUNCTION MUX_4 (in1:t_scratch;in2:t_scratch;in3:t_scratch;in4:t_scratch;sel:t_mux4) RETURN t_scratch;

FUNCTION MUX_2 (in1:t_scratch;in2:t_scratch;sel:t_mux) RETURN t_scratch;

FUNCTION ADD_SUB (in1:t_scratch;in2:t_scratch;addsel:t_add) RETURN t_scratch;

FUNCTION BIT_LOAD(in1:bit) RETURN t_load;
```

```

5
10
15
20
25
30
35
40
45
50
55

end utils_dwt;

package body utils_dwt is
FUNCTION MUX_4 (in1:t_scratch;in2:t_scratch;in3:t_scratch;in4:t_scratch;sel:t_mux4) RETURN t_scratch IS
BEGIN
CASE sel IS
WHEN uno => RETURN in1;
WHEN dos => RETURN in2;
WHEN tres => RETURN in3;
WHEN quatro => RETURN in4;
END CASE;
END;

FUNCTION MUX_2 (in1:t_scratch;in2:t_scratch;sel:t_mux) RETURN t_scratch IS
BEGIN
CASE sel IS
WHEN left => RETURN in1;
WHEN right => RETURN in2;
END CASE;
END;

FUNCTION ADD_SUB (in1:t_scratch;in2:t_scratch;addsel:t_add) RETURN t_scratch IS
BEGIN
CASE addsel IS
WHEN add => RETURN in1 + in2;
WHEN sub => RETURN in1 - in2;
END CASE;
END;

FUNCTION BIT_LOAD(in1:bit) RETURN t_load IS
BEGIN
CASE in1 IS
WHEN '1' => RETURN write;
WHEN OTHERS => RETURN read;
END CASE;
END;

END utils_dwt;

```

5

10

15

20

25

30

35

40

45

50

55

```

--VHDL Description of Tree Processor/Encoder-Decoder Circuit--
--The state machine to control the address counters#
--only works for 3 octave decomposition in y,2 in u/v#
--these are the addr gens for the x & y addresses of a pixel given the octave#
--subblk no. for each octave. Each x&y address is of the form
-- x = count(5 bits)(blk(3) to blk(octave+1))(s) (octave 0's)
-- y = count(5 bits)(blk(3) to blk(octave+1))(s) (octave 0's)
--this makes up the 9 bit address for CIF images
--the blk & s counters are vertical 2 bit with the lab in the x coord
--and carry out on 3, last counter is both horis and vertical counter
--read_enable enable the block count for the read address, but not the
--carry-outs for the mode change, this is done on the write addr cycle
--by write_enable, so same address values generated on read & write cycles#

```

```

use work.DWT_TYPES.all;
use work.dff_package.all;

```

```

entity U_ADDR_GEN is

```

```

port(
  ck : in bit ;
  reset : in t_reset ;
  new_channel , channel : in t_channel ;
  load_channel : in t_load ;
  sub_count : in BIT_VECTOR(1 to 2) ;
  col_length : in BIT_VECTOR(1 to xszize) ;
  row_length : in BIT_VECTOR(1 to ysize) ;
  ximage_string : in BIT_VECTOR(1 to xszize) ;
  yimage_string : in BIT_VECTOR(1 to ysize) ;
  yimage_string_3 : in BIT_VECTOR(1 to 11) ;
  read_enable, write_enable : in bit ;
  new_mode : in t_mode ;

  out_1 : out t_memory_addr;
  out_2 : out t_octave;
  out_3 : out bit;

```

```

5
10
15
20
25
30
35
40
45
50
55

    out_4 : out_bit;
    out_5 : out_bit;
    out_6 : out_t_state);
end U_ADDR_GEN;

architecture behave of U_ADDR_GEN is

    component U_CONTROL_ENABLE
    port(
        ck : in bit;
        reset : in t_reset;
        new_channel, channel : in t_channel;
        c_blk : in BIT_VECTOR(1 to 3);
        subband : in BIT_VECTOR(1 to 2);
        load_channel : in t_load;
        new_mode : in t_mode;

        out_1 : out_BIT_VECTOR(1 to 3);
        out_2 : out t_octave;
        out_3 : out_bit;
        out_4 : out_bit;
        out_5 : out t_state);

    end component;

    component COUNTER
    generic (ncount:integer);
    port(
        ck:in bit;
        reset:in t_reset;
        en:in bit;
        x_lpf:in bit_vector(1 to ncount);
        q:out bit_vector(1 to ncount);
        carry:out bit);
    end component;

```

```

5
10
15
20
25
30
35
40
45
50
55

COMPONENT BLK_SUB_COUNT
PORT(
  ck:in bit; reset:in t_reset; en,cln_en,cout_en:in bit; q:out bit_vector(1 to 2); carry:out bit);
end COMPONENT;

  signal rw_enable:bit;
  signal y_lpf:BIT_VECTOR(1 to ysize-4);
  signal x_lpf:BIT_VECTOR(1 to xsize-4);
  signal tree_done:bit:='0';
  signal lpf_done:bit:='0';
  signal lpf_block_done:bit:='0';
  signal sub_en:bit;
  signal y_en:bit;
  signal x_en:bit;
  signal blk_en:BIT_VECTOR(1 to 3):="B'000";
  signal octave:t_octave:=0;
  signal control_4:it_state:=down1;
  signal x_count_1:BIT_VECTOR(1 to xsize-4);
  signal x_count_2:bit;
  signal y_count_1:BIT_VECTOR(1 to ysize-4);
  signal y_count_2:bit;
  signal blk_count_2:BIT_VECTOR(1 to 3):="B'000";
  signal blk_count_1_1:BIT_VECTOR(1 to 2):="B'00";
  signal blk_count_1_2:bit;
  signal blk_count_2_1:BIT_VECTOR(1 to 2):="B'00";
  signal blk_count_2_2:bit;
  signal blk_count_3_1:BIT_VECTOR(1 to 2):="B'00";
  signal blk_count_3_2:bit;
  signal x_msb_out:BIT_VECTOR(1 to xsize-3);
  signal x_lsb_out:BIT_VECTOR(1 to 3);
  signal y_msb_out:BIT_VECTOR(1 to ysize-3);
  signal y_lsb_out:BIT_VECTOR(1 to 3);
  signal x_addr:BIT_VECTOR(1 to xsize);
  signal y_addr:BIT_VECTOR(1 to ysize);
  signal base_rows:BIT_VECTOR(1 to 11);
  signal mult_fac:BIT_VECTOR(1 to xsize);

```

```

5
10
15
20
25
30
35
40
45
50
55

signal int_addr:integer:=0;
signal temp:integer:=0;
signal address_x:memory_addr;
signal address_y:t_memory_addr;
signal address_y:t_memory_addr;

BEGIN

--size of lpf/2 -1, for y,u|v. 2 because count in pairs of lpf values
--lpf same size for all channels

y_lpf <= row_length(1 to ysize-4);
x_lpf <= col_length(1 to xsize-4);

x_en<= '1' WHEN tree_done='1' OR lpf_block_done= '1' ELSE
        '0';

--clk y_count when all blocks done for subs 1-3, or when final blk done for lpf
y_en<= '1' WHEN sub_count = 8'00" AND lpf_block_done='1' AND x_count_2='1' ELSE
        '1' WHEN sub_count /= 8'00" AND tree_done= '1' AND x_count_2='1' ELSE
        '0';

--enable the sub band counter#
sub_en<= '1' WHEN y_count_2='1' AND y_en='1' ELSE
        '0';

lpf_done <= sub_en WHEN sub_count = 8'00" ELSE '0';

WITH channel SELECT
x_msb_out<= x_count_1 & blk_count_3_1(2) WHEN y,
--always the msb bits#
        B'0" & x_count_1 WHEN u|v ;

WITH channel SELECT
y_msb_out<= y_count_1 & blk_count_3_1(1) WHEN y,
        B'0" & y_count_1 WHEN u|v ;

```

```

5
10
15
20
25
30
35
40
45
50
55

WITH octave SELECT
--bit 2 is lab#
  x_lab_out<= blk_count_2_1(2) & blk_count_1_1(2) & sub_count(2) WHEN 0 ,
               blk_count_2_1(2) & sub_count(2) & '0' WHEN 1,
               sub_count(2) & '0' & '0' WHEN 2,
               b'000' WHEN OTHERS;

WITH octave SELECT
--bit 1 is mab#
  y_lab_out<= blk_count_2_1(1) & blk_count_1_1(1) & sub_count(1) WHEN 0 ,
               blk_count_2_1(1) & sub_count(1) & '0' WHEN 1,
               sub_count(1) & '0' & '0' WHEN 2,
               b'000' WHEN OTHERS;

x_addr <= x_mab_out & x_lab_out;
y_addr <= y_mab_out & y_lab_out;

WITH channel SELECT
base_rows<=b'000000000000' WHEN y,
          b'0' & yimage_string(1 to ysize) & b'0' WHEN u,
          yimage_string_3 WHEN v;

--base address for no of rows for y,u & v memory areas#

WITH channel SELECT
mult_fac<=ximage_string WHEN y,
          b'0' & ximage_string(1 to xsize-1) WHEN u|v;

address_x<= U_TO_I(x_addr);
address_y<= U_TO_I(y_addr);
address <= U_TO_I(x_addr) + ( U_TO_I(y_addr) + U_TO_I(base_rows)) * U_TO_I(mult_fac) );

```



```

5
10
15
20
25
30
35
40
45
50
55

blk_count_2 <= blk_count_1_2 & blk_count_2_2 & blk_count_3_2;
rw_enable <= read_enable OR write_enable;
cnt1: COUNTER GENERIC MAP(xsize-4) PORT MAP(ck,reset,x_en,x_lpf,x_count_1,x_count_2);
cnt2: COUNTER GENERIC MAP(ysize-4) PORT MAP(ck,reset,y_en,y_lpf,y_count_1,y_count_2);
--use new_channel so on channel change control state picks up correct values
cnt_en: U_CONTROL_ENABLE PORT MAP(ck,reset,new_channel,channel,blk_count_2,
    sub_count,load_channel,new_mode, blk_en,octave,tree_done,lpf_block_done,control_4);
bsub_1: BLK_SUB_COUNT PORT MAP(ck,reset,blk_en(1),rv_enable,write_enable,blk_count_1_1,blk_count_1_2);
bsub_2: BLK_SUB_COUNT PORT MAP(ck,reset,blk_en(2),rv_enable,write_enable,blk_count_2_1,blk_count_2_2);
bsub_3: BLK_SUB_COUNT PORT MAP(ck,reset,blk_en(3),rv_enable,write_enable,blk_count_3_1,blk_count_3_2);
--procedure outputs#
out_1 <= address;
out_2 <= octave;
out_3 <= sub_en;
out_4 <= tree_done;
out_5 <= lpf_done;
out_6 <= control_4;

end behave;

CONFIGURATION ADDR_GEN_CON OF U_ADDR_GEN IS
FOR behave
    FOR ALL : BLK_SUB_COUNT USE CONFIGURATION WORK.BLK_SUB_CON;
END FOR;
    FOR ALL : COUNTER USE CONFIGURATION WORK.COUNTER_CON;
END FOR;
    FOR cnt_en : U_CONTROL_ENABLE USE CONFIGURATION WORK.CONTROL_ENABLE_CON;
END FOR;
END FOR;
END ADDR_GEN_CON;

--a counter to control the sequencing ofw, token, huffman cycles--
--decide reset is enabled 1 cycle early, and latched to avoid glitches--
--lpf_stop is a dummy mode to disable the block writeshuffman data--
--cycles for that block--

```

5

10

15

20

25

30

35

40

45

50

55

```

use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_CONTROL_COUNTER IS
PORT(
    ck : in bit ;
    reset : in t_reset ;
    mode,new_mode : in t_mode ;
    direction : in t_direction ;

    out_0 : out t_load;
    out_1 : out t_cycle;
    out_2 : out t_reset;
    out_3 : out bit;
    out_4 : out bit;
    out_5 : out t_load;
    out_6 : out t_cs;
    out_7 : out t_load;
    out_8 : out t_cs) ;

    --mode load,cycle,decide reset,read_addr_enable,write_addr_enable,load flags--
    --decode write_addr_enable early and latch to avoid feedback loop with pro_mode--
    --in MODE_CONTROL--
end U_CONTROL_COUNTER;

```

```

architecture behave of U_CONTROL_COUNTER IS
    COMPONENT COUNT_SYNC
    GENERIC (n:integer);
    PORT(
        ck:in bit ;
        reset:in t_reset;
        en:in bit;
        q:out bit_vector(1 to n);
        carry:out bit);
    end COMPONENT;

```



```

5
10
15
20
25
30
35
40
45
50
55

WHEN forward =>
    CASE mode IS
        WHEN send|still_send|lpf_send =>
            WHEN 0 to 3 => read_addr_enable := '1';
            cs_new := sel;
            WHEN 4 => cycle := token_cycle;
            load_flags := write;
            write_addr_enable := '1';

            WHEN 5 to 7 => write_addr_enable := '1';
            CASE new_mode IS
                WHEN stop|lpf_stop => cycle := skip_cycle;
                rw_old := read;
                cs_old := no_sel;
                WHEN void => cycle := skip_cycle;
                rw_old := write;
                WHEN OTHERS => cycle := data_cycle;
                rw_old := write;
            END CASE;
            WHEN 8 => decide_reset := rat;
            CASE new_mode IS
                WHEN stop|lpf_stop => cycle := skip_cycle;
                rw_old := read;
                cs_old := no_sel;
                WHEN void => cycle := skip_cycle;
                load_mode := write;
                rw_old := write;

                WHEN OTHERS => cycle := data_cycle;
                load_mode := write;
                rw_old := write;
            END CASE;
            WHEN OTHERS => null;
            END CASE;

        WHEN still =>
            CASE count_len IS
                WHEN 0 to 3 => read_addr_enable := '1';
                cs_new := sel;
            END CASE;
    END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

WHEN 4 => cycle := token_cycle;
        write_addr_enable := '1';
        load_flags := write;
WHEN 5 to 7 => rw_old := write;
        write_addr_enable := '1';
        CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
        END CASE;

WHEN 8 => decide_reset := rat;
        rw_old := write;
        load_mode := write;
        CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
        END CASE;

WHEN OTHERS => null;
END CASE;

WHEN lpf_still => CASE count_len IS
    WHEN 0 to 3 => read_addr_enable := '1';
        ca_new := sel;
    WHEN 4 => cycle := token_cycle;
        write_addr_enable := '1';
        load_flags := write;
    WHEN 5 to 7 => cycle := data_cycle;
        rw_old := write;
        write_addr_enable := '1';
    WHEN 8 => cycle := data_cycle;
        rw_old := write;
        decide_reset := rat;
        load_mode := write;
    WHEN OTHERS => null;
END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN void =>
        CASE count_len IS
            WHEN 0 to 3 => read_addr_enable := '1';
                cs_new := sel;
            WHEN 4 => load_flags := write;
                cycle := token_cycle;

            WHEN 5 to 7 => write_addr_enable := '1';
                write_addr_enable := '1';

            CASE new_mode IS
                WHEN stop => rw_old := read;
                    cs_old := no_sel;
                WHEN OTHERS => rw_old := write;
            END CASE;

            WHEN 8 => decide_reset := rst;
            CASE new_mode IS
                WHEN stop => rw_old := read;
                    cs_old := no_sel;
                WHEN OTHERS => load_mode := write;
                    rw_old := write;
            END CASE;

            WHEN OTHERS => null;
        END CASE;

    WHEN void_still =>
        CASE count_len IS
            WHEN 0 => write_addr_enable := '1';

            WHEN 1 to 3 => write_addr_enable := '1';
                rw_old := write;
            WHEN 4 => rw_old := write;
                load_mode := write;
                decide_reset := rst;
            WHEN OTHERS => null;
        END CASE;

    WHEN OTHERS => null;
END CASE;

--dummy token cycle for mode update--
--keep counters going--
--allow for delay--

```

```

35      WHEN inverse =>
36          CASE mode IS
37              WHEN send|still_send|lpf_send =>
38                  CASE count_len IS
39                      WHEN 0 to 3 => read_addr_enable := '1';
40                      WHEN 4 => cycle := token_cycle;
41                          write_addr_enable := '1';
42                          load_flags := write;
43                      WHEN 5 to 7 => write_addr_enable := '1';
44                      CASE new_mode IS
45                          WHEN stop|lpf_stop => cycle := skip_cycle;
46                              rw_old := read;
47                              cs_old := no_sel;
48                          WHEN void => cycle := skip_cycle;
49                              rw_old := write;
50                          WHEN OTHERS => cycle := data_cycle;
51                              rw_old := write;
52                      END CASE;
53                      WHEN 8 => decide_reset := rat;
54                      CASE new_mode IS
55                          WHEN stop|lpf_stop => cycle := skip_cycle;
56                              rw_old := read;
57                              cs_old := no_sel;
58                          WHEN void => cycle := skip_cycle;
59                              load_mode := write;
60                              rw_old := write;
61                      WHEN OTHERS => cycle := data_cycle;
62                          load_mode := write;
63                          rw_old := write;
64                      END CASE;
65                  WHEN OTHERS => null;
66              END CASE;
67          CASE count_len IS
68              WHEN 0 => null ;
69              WHEN 1 => cycle := token_cycle;
70          END CASE;
71      --skip to allow reset in huffman--

```

```

5
10
15
20
25
30
35
40
45
50
55

    write_addr_enable := '1';
    WHEN 2 to 4 => rw_old := write;
    write_addr_enable := '1';
    CASE new_mode IS
    WHEN void_still => cycle := skip_cycle;
    WHEN OTHERS => cycle := data_cycle;
    END CASE;

    WHEN 5 => rw_old := write;
    decide_reset := rat;
    load_mode := write;
    CASE new_mode IS
    WHEN void_still => cycle := skip_cycle;
    WHEN OTHERS => cycle := data_cycle;
    END CASE;

    WHEN OTHERS => null;
    END CASE;

    WHEN lpf_still => CASE count_len IS
    WHEN 0 => null;
    WHEN 1 => write_addr_enable := '1';
    WHEN 2 to 4 => cycle := data_cycle;
    rw_old := write;
    write_addr_enable := '1';
    WHEN 5 => cycle := data_cycle;
    rw_old := write;
    decide_reset := rat;
    load_mode := write;
    WHEN OTHERS => null;
    END CASE;
    CASE count_len IS
    WHEN 0 to 3 => read_addr_enable := '1';
    WHEN 4 => load_flags := write;
    cycle := token_cycle;
    write_addr_enable := '1';

--dummy token cycle for mode update--

--match with previous--
--skip for write enb delay--

```



```

5
10
15
20
25
30
35
40
45
50
55

WHEN 5 to 7 => write_addr_enable := '1';
CASE new_mode IS
WHEN stop => rw_old := read;
               ce_old := no_sel;
WHEN OTHERS => rw_old := write;
END CASE;

WHEN 8 => decide_reset := rst;
CASE new_mode IS
WHEN stop => rw_old := read;
               ce_old := no_sel;
WHEN OTHERS => load_mode := write;
               rw_old := write;
END CASE;

WHEN OTHERS => null;
END CASE;

CASE count_len IS
WHEN 0 => null;

WHEN 1 => write_addr_enable := '1';
WHEN 2 to 4 => write_addr_enable := '1';
               rw_old := write;
WHEN 5 => rw_old := write;
               load_mode := write;
               decide_reset := rst;
WHEN OTHERS => null;
END CASE;

WHEN void_still =>

--match with rest--
--dummy as write delayed--

END CASE;

write_sig <= write_addr_enable;
decide_sig <= decide_reset;

DFF(ck,reset,write_sig,write_del);
out_0 <= load_mode;
out_1 <= cycle;

WHEN OTHERS => null;
END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

out_2 <= decide_sig;
out_3 <= read_addr_enable;
out_4 <= write_del;
out_5 <= load_flags;
out_6 <= cs_new;
out_7 <= rw_old;
out_8 <= cs_old;

END PROCESS;

WITH reset SELECT
count_reset <= rst WHEN rst,
               decide_sig WHEN OTHERS;

control_cnt: count_sync GENERIC MAP(4) PORT MAP(ck,count_reset,always_one,count_1,count_2);

END behave;

CONFIGURATION CONTROL_COUNTER_CON OF U_CONTROL_COUNTER IS
FOR behave
    FOR ALL:count_sync USE ENTITY WORK.count_sync(behave);
    END FOR;
END FOR;

END CONTROL_COUNTER_CON;

--THE STATE machine to control the address counters#
--only works for 3 octave decomposition in y & 2 in u|v#

use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_CONTROL_ENABLE is

port(
ck : in bit ;
reset : in t_reset ;
new_channel,channel : in t_channel ;
c_blk : in BIT_VECTOR(1 to 3) ;

```

```

5
10
15
20
25
30
35
40
45
50
55

subband : in BIT_VECTOR(1 to 2) ;
load_channel : in t_load ;
new_mode : in t_mode ;

out_1 : out BIT_VECTOR(1 to 3);
out_2 : out t_octave;
out_3 : out bit;
out_4 : out bit;
out_5 : out t_state ;

end U_CONTROL_ENABLE;

architecture behave of U_CONTROL_ENABLE is
    signal      state:t_state;
    signal      new_state_sig:t_state;
begin

    state_machine:process(reset,new_channel,channel,c_blk,subband,load_channel,new_mode,state,new_state_sig)

        VARIABLE en_blk:BIT_VECTOR(1 to 3) := 8"000";
        --enable blk_count#
        variable      lpf_block_done:bit := '0';
        --enable x_count for LPP#
        variable      tree_done:bit := '0';
        --enable x_count for other subbands#
        variable      reset_state:t_state;
        variable      new_state:t_state ;
        variable      octave:t_octave := 0;
        --current octave#
        variable      start_state:t_state;
        -- dummy signals for DP1

    BEGIN

        -- default initial conditions
        en_blk:=b"000";
        lpf_block_done:= '0';
        tree_done:= '0';

```

```

5
10
15
20
25
30
35
40
45
50
55

octave:= 0;
reset_state:=up0;
new_state:=state;
start_state:=up0;
--set up initial state thro mux on reset, on HH stay in zz0 state#
CASE channel IS
WHEN
WHEN
END CASE;
CASE reset IS
WHEN rst => reset_state:= start_state;
WHEN OTHERS
=> reset_state := state;
END CASE;

CASE reset_state IS
WHEN up0 => octave :=2;
en_blk(3):='1';
CASE c_blk(3) IS
WHEN '1' =>
CASE subband IS
WHEN B'00" => lpf_block_done := '1';
WHEN OTHERS => new_state := up1;
END CASE;
--clock x_count for LPP y channel#
--change state when count done#

--in luminance & done with that tree#
CASE new_mode IS
WHEN stop => tree_done := '1';
WHEN OTHERS => null;
END CASE;
WHEN OTHERS => null;
END CASE;
WHEN up1 => octave :=1;
en_blk(2):='1';
CASE c_blk(2) IS
WHEN '1' => new_state := zz0;

```



```

CASE channel IS
WHEN u/v => IF c_blk(1)='1' AND c_blk(2)='1' THEN tree_done := '1';
            ELSE null;
            END IF;

WHEN y => IF c_blk(1)='1' AND c_blk(2)='1' AND c_blk(3)='1' THEN
            tree_done := '1';
            ELSE null;
            END IF;

END CASE;

--now change to start state if the sequence has finished#
CASE tree_done IS
--in LPP state doesn't change when block done#
WHEN '1' => new_state := start_state;
WHEN OTHERS => null;
END CASE;

--on channel change, use starting state for new channel#
CASE load_channel IS
--in LPP state doesn't change when block done#
WHEN write => CASE new_channel IS
                WHEN y => new_state:= up0;
                WHEN u/v => new_state:=down1;
                END CASE;
WHEN OTHERS => null;
END CASE;

new_state_sig<=new_state;

out_1 <= en_blk;
out_2 <= octave;
out_3 <= tree_done;
out_4 <= lpf_block_done;
out_5 <= reset_state;

END PROCESS;

```

```

5
10
15
20
25
30
35
40
45
50
55

DF1(ck,new_state_eig,etate);
END behave;

--
CONFIGURATION CONTROL_ENABLE_CON OF U_CONTROL_ENABLE is
FOR behave
END FOR;
END CONTROL_ENABLE_CON;

--The basic toggle flip-flop plus and gate for a synchronous counter
--input t is the toggle ,outputs are q and tc (toggle for next counter)
--stage
-- reset is asynchronous, is active on final count
-- use work.DWT_TYPES.all;
-- use work.dff_package.all;

entity BASIC_COUNT is
PORT(
ck:in bit ;reset:in t_reset;en:in bit;q:out bit;carry:out bit);
end BASIC_COUNT;

architecture behave OF BASIC_COUNT is
signal dlat:bit;
signal in_dff:bit;
signal reset_bit:bit;
BEGIN
WITH reset SELECT
reset_bit <= '0' WHEN reset,
'1' WHEN no_reset;
in_dff<=(dlat XOR en) AND reset_bit;
DF1(ck,in_dff,dlat);
carry<=dlat AND en;
q<=dlat;
END behave;

```



```

5
10
15
20
25
30
35
40
45
50
55

configuration basic_count_con of basic_count is
    FOR behave
        END FOR;
    end basic_count_con;

    -- The n-bit macro counter generator, en is the enable, the outputs
    -- are msb(bit 1)....lsb,carry.This is the same order as ELLA strings are stored;

    use work.DWT_TYPES.all;

    entity COUNT_SYNC is
        GENERIC (n:integer);
        PORT(
            ck:in bit ;
            reset:in t_reset;
            en:in bit;
            q:out bit_vector(1 to n);
            carry:out bit);
        end COUNT_SYNC;

    architecture behave OF COUNT_SYNC is

        COMPONENT basic_count
        PORT(
            ck:in bit ;reset:in t_reset;en:in bit;q:out bit;carry:out bit);
        end COMPONENT;

        signal enable:bit_vector(1 to n+1);
        BEGIN
            enable(n+1)<=en;
            cl: for i in n downto 1 generate
                bc: basic_count PORT MAP(ck,reset,enable(i+1),q(i),enable(i));
            end generate;
            carry<=enable(1);
        end behave;

```

```

5
10
15
20
25
30
35
40
45
50
55

--configuration for simulation
CONFIGURATION COUNT_SYNC_CON OF COUNT_SYNC is
FOR behave
    FOR ALL:basic_count USE ENTITY WORK.basic_count(behave);
    END FOR;
END FOR;
END COUNT_SYNC_CON;

--the basic x/y counter, carry out 1 cycle before final count given by x_lpf/y_lpf#
use work.DWT_TYPES.all;
use work.dff_package.all;

entity COUNTER is
    GENERIC (ncount:integer);
    PORT(
        ck:in bit ;
        reset:in t_reset;
        en:in bit;
        x_lpf:in bit_vector(1 to ncount);
        q:out bit_vector(1 to ncount);
        carry:out bit);
    end COUNTER;

architecture behave OF COUNTER is

    COMPONENT count_sync
    GENERIC (n:integer);
    PORT(
        ck:in bit ;reset:in t_reset;en:in bit;q:out bit_vector(1 to ncount);carry:out bit);
    end COMPONENT;

    signal cnt_reset:t_reset;
    signal final_count:bit;
    signal final_cnt_d:bit;
    signal q_sync:bit_vector(1 to ncount);
    signal carry_sync:bit;
BEGIN

```

```

cnt_sy: count_sync GENERIC MAP(ncount) PORT MAP(ck,cnt_rstet,en,q_sync,carry_sync);
q<=q_sync;
carry<=final_count;

final_count <= '1' WHEN q_sync=x_lpf AND en = '1' ELSE '0';

cnt_rstet <= rst WHEN rstet=rst ELSE
    rst WHEN final_count = '1' ELSE
    no_rst;

END behave;

CONFIGURATION COUNTER_CON OF COUNTER IS
FOR behave
FOR ALL:count_sync USE CONFIGURATION WORK.count_sync_con;
END FOR;
END FOR;
END COUNTER_CON;

--the blk, or sub-band counters, carry out on 3, cout_en enables the carry out, & cin_en AND en enables the
count# use work.DWT_TYPES.all;

entity BLK_SUB_COUNT is
PORT(
ck:in bit ;reset:in t_rstet,en,cin_en,cout_en:in bit;q:out bit_vector(1 to 2);carry:out bit);
end BLK_SUB_COUNT;

architecture behave OF BLK_SUB_COUNT is
COMPONENT count_sync
GENERIC (n:integer);
PORT(
    ck:in bit ;reset:in t_rstet,en:in bit;q:out bit_vector(1 to 2);carry:out bit);
end COMPONENT;

signal q_sync:bit_vector(1 to 2);

```

```

5
10
15
20
25
30
35
40
45
50
55

signal carry_sync:bit;
signal enable:bit;
    xc;
BEGIN
    enable <= en AND cin_en;
    q<=q_sync;
    carry<= '1' WHEN q_sync = b"11" AND cout_en = '1' ELSE '0';
        b_cnt: count_sync GENERIC MAP(2) PORT MAP(ck,reset,enable,q_sync,carry_sync);
END behave;

CONFIGURATION BLK_SUB_CON OF BLK_SUB_COUNT IS
FOR behave
    FOR b_cnt : count_sync USE CONFIGURATION WORK.count_sync_con;
        END FOR;
    END FOR;
END BLK_SUB_CON;
--the L1 norm comparison constant flag values--
--adding 4 absolute data values so result can grow by 2 bits--
--5 cycle sequence, a reset cycle with no data input, followed--
--by 4 data cycles--
    use work.DWT_TYPES.all;
    use work.dff_package.all;
    use work.utils.all;

entity U_LINORM IS
GENERIC(n:integer);
PORT(
    ck : in bit ;
    reset : in t_reset ;
    in_s : in BIT_VECTOR(1 to n) ;
    out_1 : out BIT_VECTOR(1 to n+2) );

```

```

5
10
15
20
25
30
35
40
45
50
55

end U_LlNORM;

architecture behave of U_LlNORM is
    --
    signal mab:BIT_VECTOR(1 to n) ;
    signal add_in1:BIT_VECTOR(1 to n) ;
    signal rst_mux:BIT_VECTOR(1 to n+4) ;
    signal in2:BIT_VECTOR(1 to n+4) ;
    signal add_out:BIT_VECTOR(1 to n+5) ;
    signal carryit_carry;
    signal adder :integer;

BEGIN

    mab <= ALL_SAME(n,in_e(1));
    add_in1 <= (in_e XOR mab);

    WITH reset SELECT
    rst_mux <= ZERO(n+4) WHEN rst ,
               add_out(2 to n+5) WHEN OTHERS;

    --carryin bit to adder
    carry <= 1 WHEN in_e(1)='1' ELSE
              0;

    adder <= S_TO_I(add_in1) + S_TO_I(in2) + carry;
    I_TO_S(adder,add_out);

    DP1(n+4,ck,rst_mux,in2);
    --procedure outputs--
    out_1 <= in2(3 to n+4);

END;
```

```

5
10
15
20
25
30
35
40
45
50
55

CONFIGURATION U_LINORM_CON OF U_LINORM IS
FOR behave
END FOR;
END U_LINORM_CON;
--the block to decide if all its inputs are all 0--
use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_ALL_ZERO IS
PORT(
    ck : in bit ;
    reset : in t_reset ;
    in_in : in t_input ;
    out_1 : out bit );
end U_ALL_ZERO;

architecture behave of U_ALL_ZERO IS

    signal out_b:bit;
    signal in_eq_0:bit;
    signal all_eq_0:bit;
    BEGIN

        in_eq_0 <= '1' WHEN in_in = 0 ELSE
            '0';

        --1 if reset high; & OR with previous flag--

        all_eq_0 <= in_eq_0 WHEN reset = rst ELSE
            '0' WHEN out_b='0' ELSE
                in_eq_0;

        DFF1(ck,all_eq_0,out_b);

```

```

5
10
15
20
25
30
35
40
45
50
55

--procedure outputs--
out_1 <= out_bit;
END;

CONFIGURATION U_ALL_ZERO_CON OF U_ALL_ZERO IS
FOR behave
END FOR;
END U_ALL_ZERO_CON;

use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utils.all;

entity U_ABS_NORM IS
GENERIC(n:positive);
PORT(
    ck : in bit ;
    reset : in t_reset ;
    qshift : in BIT_VECTOR(1 to result_exp-2) ;
    in_in : in t_input:=0;

    out_1 : out BIT_VECTOR(1 to n+2);
    out_2 : out bit;

    end U_ABS_NORM;

architecture behave OF U_ABS_NORM IS

    signal addr_str:BIT_VECTOR(1 to n+5);
    signal rst_mux:BIT_VECTOR(1 to n+4);
    signal in2:BIT_VECTOR(1 to n+4);
    signal add_s:BIT_VECTOR(1 to n+4);
    signal addr:integer:=0;
    signal abs_in:integer:=0;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal in_small:bit;
signal all_small:bit;
signal out_b:bit;
BEGIN
    abs_in <= abs(ln_in);
    adder <= abs_in + S_TO_I(ln2);
    I_TO_S(adder,adder_str);
    add_s <= adder_str(2 to (n+5));

    WITH reset SELECT
        rst_mux <= ZERO(n+4) WHEN rst ,
                add_s WHEN OTHERS;

    in_small <= '1' WHEN abs_in <= U_TO_I(qshift) ELSE
                '0';

    --1 if reset high, & OR with previous flag--
    all_small <= '1' WHEN reset= rst ELSE
                '0' WHEN in_small = '0' ELSE
                out_b;

    DFI(n+4,ck,rst_mux,ln2);
    DFI(ck,all_small,out_b);
    --procedure outputs--
    out_1 <= ln2(3 to n+4);
    out_2 <= out_b;

    END;

    CONFIGURATION U_ABS_NORM_CON OF U_ABS_NORM IS
    FOR behave

```



```

5
10
15
20
25
30
35
40
45
50
55

END FOR;
END U_ABS_NORM_CON;

--the decide fn block--
use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utils.all;

entity U_DECIDE IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  q_int : in t_result ;
  nw,old : in t_input ;
  threshold,comparison : in t_result ;
  octa : in t_octave ;
  load_flags : in t_load ;

  out_1 : out BIT_VECTOR(1 to 7) );
end U_DECIDE;

architecture behave of U_DECIDE IS

  COMPONENT U_LINORM
    GENERIC(n:integer);
  PORT(
    ck : in bit ;
    reset : in t_reset ;
    in_s : in BIT_VECTOR(1 to n) ;

    out_1 : out BIT_VECTOR(1 to n+2) );
  end COMPONENT;

  COMPONENT U_ABS_NORM
    GENERIC(n:positive);

```

```

5
10
15
20
25
30
35
40
45
50
55

PORT(
  ck : in bit ;
  reset : in t_reset ;
  qshift : in BIT_VECTOR(1 to result_exp-2) ;
  in_in : in t_input ;

  out_1 : out BIT_VECTOR(1 to n+2) ;
  out_2 : out bit ;
  end COMPONENT ;

  --nzflag,origin,noflag,ozflag,motion,pro_new_z,pro_no_z--

  signal nz_plus_oz:BIT_VECTOR(1 to input_exp+3);
  signal shift_add:BIT_VECTOR(1 to input_exp+3);
  signal nw_str:BIT_VECTOR(1 to input_exp);
  signal old_str:BIT_VECTOR(1 to input_exp);
  signal q_int_str :BIT_VECTOR(1 to result_exp);
  signal n_o_str:BIT_VECTOR(1 to input_exp+1);
  signal nz_1:BIT_VECTOR(1 to input_exp+2);
  signal oz_1:BIT_VECTOR(1 to input_exp+2);
  signal no_1:BIT_VECTOR(1 to input_exp+3);
  signal qshift:BIT_VECTOR(1 to result_exp-2);
  signal flags:BIT_VECTOR(1 to 7);
  signal decide_flags:BIT_VECTOR(1 to 7);
  signal n_o: integer;
  signal nz: natural:=0;
  signal oz: natural:=0;
  signal no: natural:=0;
  signal nzflag: bit;
  signal ozflag: bit;
  signal noflag: bit;
  signal origin: bit;
  signal motion: bit;
  signal new_z: bit;
  signal no_z: bit;
  signal nz_2: bit;
  signal no_2: bit;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal octs_del: t_octave;
BEGIN
    I_TO_S(q_int,q_int_str);
    qshift <= q_int_str(1 to result_exp-2);
    --divide by 4 as test is on coeff values not block values--
    n_o <= nw - old;
    --new-old; use from quant--

    -- convert to string for LINORM
    I_TO_S(n_o,n_o_str);
    I_TO_S(nw,nw_str);
    I_TO_S(old,old_str);

    --convert to unsigned integer
    nz <= U_TO_I(nz_1);
    oz <= U_TO_I(oz_1);
    no <= U_TO_I(no_1);

    nzflag <= '1' WHEN nz <= threshold ELSE
        '0';
    noflag <= '1' WHEN no <= comparison ELSE
        '0';
    ozflag <= '1' WHEN oz = 0 ELSE
        '0';
    origin <= '1' WHEN nz <= no ELSE
        '0';
    I_TO_S(nz + oz,nz_plus_oz);
    new_z <= nz_2;
    no_z <= no_2;

    --delay octs to match pipelin delay--
    DFI(ck,octs,octs_del);

```

```

5
10
15
20
25
30
35
40
45
50
55

--keep 13 bits here to match no; keep msb's--
--delay octs to match pipelin delay--
WITH octa_del SELECT

    shift_add <= nz_plus_oz(1 to input_exp+3) WHEN 0,
                B"0"& nz_plus_oz(1 to input_exp+2) WHEN 1,
                B"00"& nz_plus_oz(1 to input_exp+1) WHEN 2,
                B"000"& nz_plus_oz(1 to input_exp) WHEN 3;

    motion <= '1' WHEN U_TO_I(shift_add) <= no ELSE
              '0';

    decide_flags <= nzflag&origin&noflag&ozflag&motion&new_z&no_z;

    abs_1: U_ABS_NORM GENERIC MAP(input_exp) PORT MAP(ck,reset,qahift,nw,nz_1,nz_2);

    LATCH(7,load_flags,decide_flags,flags);

    abs_2: U_ABS_NORM GENERIC MAP(input_exp+1) PORT MAP(ck,reset,qahift,n_o,no_1,no_2);

    ll: U_LINORM GENERIC MAP(input_exp) PORT MAP(ck,reset,old_str,oz_1);

--procedure outputs--

    out_1 <= flags ;

END;

CONFIGURATION U_DECIDE_CON OF U_DECIDE IS
FOR behave
    FOR ALL: U_ABS_NORM USE ENTITY WORK.U_ABS_NORM(behave);
    END FOR;
    FOR ALL: U_LINORM USE ENTITY WORK.U_LINORM(behave);
    END FOR;
END FOR;
END U_DECIDE_CON;
-- create the rising edge function, and a model of a active high DFF.

```

```

5
10
15
20
25
30
35
40
45
50
55

use work.DWT_TYPES.all;
use work.utils.all;
package dff_package is
FUNCTION rising_edge (SIGNAL s:bit) return bool;
PROCEDURE DFF_LOAD(
SIGNAL ck:in bit;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR);
PROCEDURE DFF_LOAD(
SIGNAL ck:in bit;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low);
PROCEDURE DFF(
SIGNAL ck:in bit;SIGNAL d:in integer;SIGNAL q:out integer);
PROCEDURE DFF(
SIGNAL ck:in bit;SIGNAL d:in t_direction;SIGNAL q:out t_direction);
PROCEDURE DFF(
SIGNAL ck:in bit;SIGNAL d:in t_state;SIGNAL q:out t_state);
PROCEDURE DFF(
SIGNAL ck:in bit;SIGNAL d:in t_reset;SIGNAL q:out t_reset);
PROCEDURE DFF(
SIGNAL ck:in bit;SIGNAL d:in bit;SIGNAL q:out bit);
PROCEDURE DFF(CONSTANT n:in integer;
SIGNAL ck:in bit;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector);
PROCEDURE DFF(
SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in integer;SIGNAL q:out integer);

```

```

5
10
15
20
25
30
35
40
45
50
55

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_reset;SIGNAL q:out t_reset);
  SIGNAL q:out t_reset;
PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in bit;SIGNAL q:out bit);
PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_load;SIGNAL q:out t_load);

PROCEDURE DPP_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in integer;SIGNAL q:out integer);
PROCEDURE DPP_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_channel;SIGNAL q:out t_channel);
PROCEDURE DPP_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_diff;SIGNAL q:out t_diff);
PROCEDURE DPP_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_mode;SIGNAL q:out t_mode);
PROCEDURE DPP_INIT(CONSTANT n:natural;
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR);
PROCEDURE DPP_INIT(
  SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low);

PROCEDURE LATCH(CONSTANT n:in integer;
  load:in t_load;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector);
end dff_package;

package body dff_package is

```

```

5
10
15
20
25
30
35
40
45
50
55

FUNCTION rising_edge (SIGNAL s:bit) return bool is
BEGIN
    IF (s'event) AND (s='1') AND (s'last_value = '0') THEN return t;
    ELSE return f;
    END IF;
END rising_edge;

--THE Df1 flip-flops, NO RESET-----
PROCEDURE Df1(
    SIGNAL ck:in bit;SIGNAL d:in integer;SIGNAL q:out integer) IS
BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
END Df1;

PROCEDURE Df1(CONSTANT n:in integer;
    SIGNAL ck:in bit;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector) IS
BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
END Df1;

PROCEDURE Df1(
    SIGNAL ck:in bit;SIGNAL d:in t_state;SIGNAL q:out t_state) IS
BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
END Df1;

PROCEDURE Df1(
    SIGNAL ck:in bit;SIGNAL d:in t_direction;SIGNAL q:out t_direction) IS
BEGIN
    IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;

```

```

5
10
15
20
25
30
35
40
45
50
55

END IF;
END DF1;

PROCEDURE DF1(
  SIGNAL ck:in bit;SIGNAL d:in t_reset;SIGNAL q:out t_reset) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END DF1;

PROCEDURE DF1(
  SIGNAL ck:in bit;SIGNAL d:in bit;SIGNAL q:out bit) IS
BEGIN
  IF(rising_edge(ck) = t ) THEN q<=d;
  ELSE null;
  END IF;
END DF1;

--THE DFF flip-flops, with RESET-----
PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in integer;SIGNAL q:out integer) IS
BEGIN
  IF reset=rat THEN q<= 0;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;
  --IF(rising_edge(ck) = t ) THEN IF reset=rat THEN q<= 0; ELSE q<=d ;END IF;
  ELSE null;
  END IF;
END DFF;

PROCEDURE DFF(
  SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_reset;SIGNAL q:out t_reset) IS
BEGIN
  IF reset=rat THEN q<= rat;
  ELSIF(rising_edge(ck) = t ) THEN q<=d;

```


5

10

15

20

25

30

35

40

45

50

55

```

        ELSE null;
      END IF;
    END DFF;

    PROCEDURE DFF(
      SIGNAL tk:in bit;reset:in t_reset;SIGNAL d:in bit;SIGNAL q:out bit) IS
    BEGIN
      IF reset=rat THEN q<= '0';
      ELSIF(rising_edge(ck) = t ) THEN q<=d;
      ELSE null;
      END IF;
    END DFF;

    PROCEDURE DFF(
      SIGNAL ck:in bit;reset:in t_reset;SIGNAL d:in t_load;SIGNAL q:out t_load) IS
    BEGIN
      IF reset=rat THEN q<= read;
      ELSIF(rising_edge(ck) = t ) THEN q<=d;
      ELSE null;
      END IF;
    END DFF;

    ---      THE DFF_INIT FLIP-FLOPS

    PROCEDURE DFF_INIT(
      SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in Integer;SIGNAL q:out Integer) IS
    BEGIN
      IF reset=rat THEN q<= 0;
      ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
      ELSE null;
      END IF;
    END IF;
    END DFF_INIT;

    PROCEDURE DFF_INIT(
      SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low) IS
    BEGIN

```

```

5
10
15
20
25
30
35
40
45
50
55

IF reset=rat THEN q<= low;
ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
    ELSE null;
    END IF;
END IF;
END DFF_INIT;

PROCEDURE DFF_INIT(
    SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_channel;SIGNAL q:out t_channel) IS
BEGIN
    IF reset=rat THEN q<= y;
    ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
        ELSE null;
        END IF;
    END IF;
END DFF_INIT;

PROCEDURE DFF_INIT(
    SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_mode;SIGNAL q:out t_mode) IS
BEGIN
    IF reset=rat THEN q<= still;
    ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
        ELSE null;
        END IF;
    END IF;
END DFF_INIT;

PROCEDURE DFF_INIT(
    SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in t_diff;SIGNAL q:out t_diff) IS
BEGIN
    IF reset=rat THEN q<= nodiff;
    ELSIF load=write THEN IF(rising_edge(ck) = t ) THEN q<=d;
        ELSE null;
        END IF;
    END IF;
END DFF_INIT;

```

```

5
10
15
20
25
30
35
40
45
50
55

PROCEDURE DFF_INIT(CONSTANT n:natural;
SIGNAL ck:in bit;reset:in t_reset;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR) IS
BEGIN
    IF reset=1 THEN q<= ZERO(d'length);
    ELSIF load=1 THEN IF(rising_edge(ck) = 1) THEN q<=d;
    ELSE null;
    END IF;
    END IF;
    END DFF_INIT;

PROCEDURE DFF_LOAD(
SIGNAL ck:in bit;load:in t_load;SIGNAL d:in BIT_VECTOR;SIGNAL q:out BIT_VECTOR) IS
BEGIN
    IF load=1 THEN IF(rising_edge(ck) = 1) THEN q<=d;
    ELSE null;
    END IF;
    END IF;
    END DFF_LOAD;

PROCEDURE DFF_LOAD(
SIGNAL ck:in bit;load:in t_load;SIGNAL d:in t_high_low;SIGNAL q:out t_high_low) IS
BEGIN
    IF load=1 THEN IF(rising_edge(ck) = 1) THEN q<=d;
    ELSE null;
    END IF;
    END IF;
    END DFF_LOAD;

PROCEDURE LATCH(CONSTANT n:n integer;
load:in t_load;SIGNAL d:in bit_vector;SIGNAL q:out bit_vector) IS
BEGIN
    IF load=1 THEN q<=d;
    ELSE null;
    END IF;
    END LATCH;

```

```

5
10
15
20
25
30
35
40
45
50
55

END dff_package;

package dwt_types is
--constant values
constant result_exp:integer:= 14;
--length of result arith
constant input_exp:integer:= 10;
--length of 1D convolver input/output#
constant qmax :integer:= 7;
--maximum shift value for quantisation constant#
constant result_range :integer:= 2 ** (result_exp-1);
constant input_range :integer:= 2 ** (input_exp-1);
constant max_octave:integer:= 3;
--no of octaves:integer:=max_octave+1; can not be less in this example#
constant no_octave:integer:= max_octave+1;
constant xsize :integer:= 10;
--no of bits for ximage#
constant ysize :integer:= 9;
--no of bits for yimage#
constant ximage:integer:= 319;
--the xdimension -1 of the image; ie no of cols#
constant yimage:integer:= 239 ,
--the ydimension -1 of the image; ie no of rows#

--int types#
subtype t_result is integer range -result_range to result_range-1;
subtype t_input is integer range -input_range to input_range-1;
subtype t_length is integer range 0 to 15;
subtype t_inp is integer range 0 to 1023;
subtype t_blk is integer range 0 to 3;
subtype t_sub is integer range 0 to 3;
subtype t_col is integer range 0 to ximage;
subtype t_row is integer range 0 to yimage;
subtype t_carry is integer range 0 to 1;
subtype t_quant is integer range 0 to qmax;
--address for result@dwt memory; ie 1 frame#
subtype t_memory_addr is integer range 0 to ( 3*(ximage+1)*(yimage+1)/2 -1) ;

```

```

5
10
15
20
25
30
35
40
45
50
55

subtype t_octave is integer range 0 to max_octave;

--bit string and boolean types types#
type bool is (f,t);
type flag is (error , ok);
--control signals#
type t_reset is (rst,no_rst);
type t_load is (write,read);
type t_load_vec is ARRAY (NATURAL RANGE <>) of t_load;
--r/wbar control#
type t_mem is (random,old_mem,new_mem);
type t_cs is (no_sel,sel);
type t_updown is (down,up);
--up/down counter control#
type t_diff is (diff,nodiff);
--diff or not in quantiser#
type t_intra is (intra,inter);
--convolver mux t and types#
type t_mux is (left,right);
type t_mux3 is (l,c,r);
type t_mux4 is (uno,dos,tres,quatro);
type t_add is (add,subt);
type t_direction is (forward,inverse);
--counter types#
type t_count_control is (count_rst,count_carry);
type t_count_2 is (one,two);
--state types#
type t_mode is (void,void_still,stop,send,still,still_send,lpf_send,lpf_still,lpf_stop);
type t_mode_vec is ARRAY (NATURAL RANGE <>) of t_mode;
type t_cycle is (token_cycle,data_cycle,skip_cycle);
type t_state is (up0,up1,zz0,zz1,zz2,zz3,down1);
--type t_state is (start,up0,up1,zz0,zz1,zz2,zz3,down1);
type t_decode is (load_low,load_high);
type t_high_low is (low,high);
type t_huffman is (pass,huffman);
type t_fifo is (ok_fifo,error_fifo);
--types for the octave control unit#

```

```

5
10
15
20
25
30
35
40
45
50
55

type t_channel is (y,u,v);
type t_channel_factor is (luminance,color);
--types for the control of memory ports#
--type t_sparcport (t_sparc_addr,t_sparc_addr,t_load,t_ca);

FUNCTION U_TO_I(bits: in bit_vector) RETURN natural;
FUNCTION S_TO_I(bits: in bit_vector) RETURN integer;
PROCEDURE I_TO_S(int:in integer; SIGNAL bits;out bit_vector);
end dwt_types;

package body dwt_types is

FUNCTION U_TO_I(bits:bit_vector) RETURN natural IS
variable result: natural:=0;
BEGIN
FOR i IN bits'range LOOP
result:=result*2 + bit'pos(bits(i));
END LOOP;
RETURN result;
END U_TO_I;

FUNCTION S_TO_I(bits:bit_vector) RETURN integer IS
variable temp:bit_vector(bits'range);
variable result: integer:=0;
BEGIN
IF bits(bits'left) = '1' THEN
temp:=NOT bits;
ELSE
temp:=bits;
END IF;
FOR i IN bits'range LOOP

```

```

5
10
15
20
25
30
35
40
45
50
55

result:=result*2 + bit_pos(temp(i));
END LOOP;
IF bits(bits'left) = '1' THEN
    result:=(-result)-1;
END IF;
RETURN result;
END S_TO_I;

PROCEDURE I_TO_S(int:in Integer; SIGNAL bits:out bit_vector) IS
    variable result:bit_vector(bits'range);
    variable temp: Integer;
BEGIN
    IF int < 0 THEN
        temp:=-(int+1);
    ELSE
        temp:=int;
    END IF;
    FOR i IN bits'reverse_range LOOP
        result(i):= bit_val(temp rem 2);
        temp:=temp/2;
    END LOOP;
    IF int<0 THEN
        result:=NOT result;
        result(bits'left):='1';
    END IF;
    bits<=result;
END I_TO_S;

FUNCTION INT_TO_S(n:natural;SIGNAL int:in Integer) RETURN bit_vector IS
    variable result:bit_vector(1 to n);
    variable temp: Integer;
BEGIN
    IF int < 0 THEN
        temp:=-(int+1);

```

```

5
10
15
20
25
30
35
40
45
50
55

ELSE temp:=int;
END IF;
FOR i IN n downto 1 LOOP
  result(i):= bit_val(temp rem 2);
  temp:=temp/2;
END LOOP;

-- check to see if integer fits in n bits
ASSERT (temp=0)
REPORT "int TO BIG FOR n BITS"
SEVERITY FAILURE;

IF int<0 THEN
  result:=NOT result;
  result(1):='1';
END IF;

RETURN result;
END INT_TO_S;

end dwt_types;
--the length of the huffman encoded word#

use work.DWT_TYPES.all;
entity U_LENGTH IS
  PORT(
    mag_out: in t_input ;
    out_1 : out BIT_VECTOR(1 to 5) );
  end U_LENGTH;

architecture behave of U_LENGTH IS
  BEGIN
  WITH mag_out SELECT

```



```

--length of input coded word#
out_1 <= B"00001" WHEN 0 ,
        B"00011" WHEN 1 ,
        B"00100" WHEN 2 ,
        B"00101" WHEN 3 ,
        B"00110" WHEN 4 ,
        B"00111" WHEN 5 ,
        B"01000" WHEN 6 ,
        B"01100" WHEN 7 to 21 ,
        B"10000" WHEN OTHERS;

end behave;
CONFIGURATION U_LENGTH_CON OF U_LENGTH IS
FOR behave
END FOR;
END U_LENGTH_CON;

--the buffer for the FIFO--
--the length of the Huffman encoded word--

use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utils.all;

entity FIFO_BUFFER IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  direction : in t_direction ;
  cycle : in t_cycle ;
  mode : in t_mode ;
  value,mag_out_huff : in t_input ;
  fifo_in : in BIT_VECTOR(1 to 16) ;
  fifo_full,fifo_empty : in t_fifo ;
  shift : in BIT_VECTOR(1 to 32) ;
  token_length : in BIT_VECTOR(1 to 2) ;
  flush_buffer : in bit ;
  lpf_quant : in t_quant ;

```

```
-- fifo_out, s, fifo_read fifo_write
```

```
out_1 : out BIT_VECTOR(1 to 16);
out_2 : out BIT_VECTOR(1 to 16);
out_3 : out BIT_VECTOR(1 to 16);
out_4 : out BIT_VECTOR(1 to 5);
out_5 : out t_load;
out_6 : out t_load;
```

```
end FIFO_BUFFER;
```

architecture behave OF FIFO_BUFFER IS

```
COMPONENT U_PULSE
```

```
PORT(
```

```
ck : in bit ;
```

```
reset : in t_reset ;
```

```
in_in : in t_load ;
```

```
out_1 : out t_load );
```

```
end COMPONENT;
```

```
COMPONENT U_LENGTH
```

```
PORT(
```

```
mag_out: in t_input ;
```

```
out_1 : out BIT_VECTOR(1 to 5) );
```

```
end COMPONENT;
```

```
subtype t_lpf_length is integer range 0 to 10;
```

```
subtype t_s_length is integer range 0 to 47;
```

```
signal huff : t_input;
```

```
signal length_lpf_bits : BIT_VECTOR(1 to 5);
```

```
signal length: BIT_VECTOR(1 to 5);
```

```
signal length_huff_bits: BIT_VECTOR(1 to 5);
```

```
signal new_s: BIT_VECTOR(1 to 5);
```

```

5
10
15
20
25
30
35
40
45
50
55

signal select_e : BIT_VECTOR(1 to 5);
signal s : BIT_VECTOR(1 to 5);
signal NEW_S_RESET : BIT_VECTOR(1 to 5);
signal new_s_int : t_e_length;
signal new_s_Sbits : t_e_length;
signal length_lpf : t_lpf_length:=0;
signal high_low_flag : t_high_low;
signal high_low : t_high_low;
signal flush_buffer_del : bit;
signal fifo_not_full : t_load;
signal data_ready : t_load;
signal fifo_write : t_load;
signal load_low_word : t_load;
signal load_low : t_load;
signal load_low_strobe : t_load;
signal reset_s : t_reset;
signal reset_2del : t_reset;
signal s_reset : t_reset;
signal load_high : t_load;
signal fifo_read : t_load;
signal write_low : t_load;
signal write_high : t_load;
signal load_high_strobe : t_load;
signal load_high_word : t_load;
signal load_s : t_load;
signal high_out_sel : t_mux;
signal low_out_sel : t_mux;
signal dir_sel : t_mux;
signal high_in : BIT_VECTOR(1 to 16);
signal low_out : BIT_VECTOR(1 to 16);
signal high_out : BIT_VECTOR(1 to 16);
signal low_word : BIT_VECTOR(1 to 16);
signal high_word : BIT_VECTOR(1 to 16);

signal low_in : BIT_VECTOR(1 to 16);
BEGIN

```

```

5
10
15
20
25
30
35
40
45
50
55

WITH direction SELECT
huff <= value WHEN forward ,
    << mag_out_huff WHEN inverse;

length_lpf <= input_exp - lpf_quant;

I_TO_S(length_lpf,length_lpf_bits);

lnth:U_LENGTH PORT MAP(huff,length_huff_bits);

length <= B"000" & token_length WHEN cycle = token_cycle ELSE
    B"00000" WHEN cycle = skip_cycle ELSE

--on LPF_STILL length fixed, given by input_exp-shift const--
    length_lpf_bits WHEN mode=lpf_still AND cycle =data_cycle ELSE length_huff_bits;

WITH direction SELECT
select_s <= B"0" & s(2 to 5) WHEN forward ,
    s WHEN OTHERS;

new_s_int <= U_TO_I(select_s) + U_TO_I(length); --6 bits--

I_TO_S(new_s_int,new_s);

new_s_5bits <= U_TO_I(new_s); -- 0<= new_s_5bits<=31

--if new s pointer >16--
--on inverse passed first 16 bits; active from (16,31) --
high_low_flag <= high WHEN new_s_5bits >= 16 ELSE low ;

--forward--
WITH fifo_full SELECT
fifo_not_full <= write WHEN ok_fifo ,
    read WHEN OTHERS;

```

```

5
10
15
20
25
30
35
40
45
50
55

Df1(ck,flush_buffer,flush_buffer_del);

--type change--
fifo_write <= write WHEN high_low = high ELSE
write WHEN flush_buffer='1' OR flush_buffer_del = '1' ELSE
read;

--flush buffer when frame finished, needs 2 cycles to clear--

--from inverse--

WITH fifo_empty SELECT
data_ready <= write WHEN ok_fifo,
read WHEN OTHERS;

load_low_word <= write WHEN high_low_flag = high AND data_ready=write ELSE read;
load_high_word <= write WHEN high_low_flag = low AND data_ready=write ELSE read;

p1ee_low:U_PULSE PORT MAP(ck,reset,load_low_word,load_low_strobe);
p1ee_high:U_PULSE PORT MAP(ck,reset,load_high_word,load_high_strobe);

--load low on reset to start things--

WITH reset_s SELECT
load_low <= write WHEN rst,
load_low_strobe WHEN no_rst; --load low word--

--delay reset for s and load_high--
Df1(ck,reset,reset_s);
Df1(ck,reset_s,reset_2del);

WITH reset_2del SELECT
load_high <= write WHEN rst, --load high next--
load_high_strobe WHEN no_rst; --load high word--

```

```

5
10
15
20
25
30
35
40
45
50
55

--read control for data_in FIFO--
fifo_read <= read WHEN load_low = write OR load_high = write ELSE write;

--control signals--

WITH direction SELECT
write_low <= fifo_not_full WHEN forward,
load_low WHEN OTHERS;

WITH direction SELECT
write_high <= fifo_not_full WHEN forward,
load_high WHEN OTHERS;

high_out_sel <= left WHEN direction = forward AND high_low = high ELSE
right WHEN direction = forward AND high_low = low ELSE
right WHEN direction = inverse AND U_TO_I(s) >= 16 ELSE
left;

low_out_sel <= right WHEN direction = forward AND high_low = high ELSE
left WHEN direction = forward AND high_low = low ELSE
right WHEN direction = inverse AND U_TO_I(s) >= 16 ELSE
left;

DFI_LOAD(ck, write_low, low_in, low_word);

DFI_LOAD(ck, write_high, high_in, high_word);

DFI_LOAD(ck, fifo_not_full, high_low_flag, high_low);

WITH direction SELECT
s_reset <= reset WHEN forward,
reset_s WHEN inverse;
-- reset_2del WHEN inverse; --make synchronous

WITH direction SELECT
load_s <= fifo_not_full WHEN forward,

```

5

10

15

20

25

30

35

40

45

50

55

```

data_ready WHEN inverse;

WITH a_reset SELECT
new_a_reset <= new_a WHEN no_reset,
               b"00000" WHEN reset;

--DPP_INIT(5,ck,a_reset,load_s, new_a,s);
DPP_INIT(5,ck,no_reset,load_s, new_a_reset,s);

WITH direction SELECT
dir_sel <= left WHEN forward,
        right WHEN inverse;

WITH direction SELECT
high_in <= shift(17 to 32) WHEN forward,
        fifo_in           WHEN inverse;

WITH direction SELECT
low_in <= shift(1 to 16) WHEN forward,
        fifo_in          WHEN inverse;

WITH high_out_sel SELECT
high_out <= high_word WHEN left,
        low_word      WHEN right;

WITH low_out_sel SELECT
low_out <= low_word WHEN left,
        high_word    WHEN right;

--architecture outputs--
out_1 <= low_word;
out_2 <= low_out;
out_3 <= high_out;
out_4 <= s;
out_5 <= fifo_read;
out_6 <= fifo_write;

```

```

5
10
15
20
25
30
35
40
45
50
55

END;

CONFIGURATION FIFO_BUFFER_CON OF FIFO_BUFFER IS
FOR behave
  FOR ALL: U_PULSE USE ENTITY WORK.U_PULSE(behave);
  END FOR;
  FOR ALL: U_LENGTH USE ENTITY WORK.U_LENGTH(behave);
  END FOR;
END FOR;
END FIFO_BUFFER_CON;
--the HUFFMAN encode & decode functions--

--a pulse generator, glitch free--
use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utils.all;

entity U_PULSE IS
PORT(
  ck : in bit ;
  reset : in t_reset ;
  in_in : in t_load ;

  out_l : out t_load );
end U_PULSE;

architecture behave OF U_PULSE IS
signal in_del: t_load;
BEGIN
  --DFF(ck,reset,in_in,in_del);
  DFF(ck,no_reset,in_in,in_del);
  out_l <= write WHEN in_in = write AND in_del = read ELSE
    read;
END behave;

```



```

5
10
15
20
25
30
35
40
45
50
55

CONFIGURATION U_PULSE_CON OF U_PULSE IS
FOR behave
END FOR;
END U_PULSE_CON;

entity U_MX16 IS
PORT(
  in_in : in BIT_VECTOR(1 to 16) ;
  sel   : in BIT_VECTOR(1 to 4) ;

  out_1 : out bit ) ;
end U_MX16;

architecture behave OF U_MX16 IS
BEGIN
  WITH sel SELECT
    out_1 <= in_in(1) WHEN B"0000" ,
              in_in(2) WHEN B"0001" ,
              in_in(3) WHEN B"0010" ,
              in_in(4) WHEN B"0011" ,
              in_in(5) WHEN B"0100" ,
              in_in(6) WHEN B"0101" ,
              in_in(7) WHEN B"0110" ,
              in_in(8) WHEN B"0111" ,
              in_in(9) WHEN B"1000" ,
              in_in(10) WHEN B"1001" ,
              in_in(11) WHEN B"1010" ,
              in_in(12) WHEN B"1011" ,
              in_in(13) WHEN B"1100" ,
              in_in(14) WHEN B"1101" ,
              in_in(15) WHEN B"1110" ,
              in_in(16) WHEN B"1111" ;

end behave;

CONFIGURATION U_MX16_CON OF U_MX16 IS
FOR behave
END FOR;
END U_MX16_CON;

```

```

5
10
15
20
25
30
35
40
-- a barrel shifter outs 16 bits from 32 input at position given by pointer s
-- bit 1 out from muxing in bits(1 to 16), bit 2 out from muxing in bits(2 to 17)
--etc, bit 16 out from muxing in bits(16 to 31)

entity U_SHIFT32_16 IS
PORT(
    buffer_in : in BIT_VECTOR(1 to 32) ;
    s : in BIT_VECTOR(1 to 5) ;
    out_1 : out BIT_VECTOR(1 to 16);
end U_SHIFT32_16;

--left justified value, s shift const--

architecture behave of U_SHIFT32_16 IS

    COMPONENT U_MX16
    PORT(
        in_in : in BIT_VECTOR(1 to 16) ;
        sel: in BIT_VECTOR(1 to 4) ;

        out_1: out bit) ;
    end COMPONENT;

    signal shift:BIT_VECTOR(1 to 5);
    signal shift_in:BIT_VECTOR(1 to 4);
    signal temp:BIT;
    signal temp_in:BIT_VECTOR(1 to 16);
    BEGIN

        shift <= s AND B"01111";
        --architecture outputs--

        mux16: for i in 1 to 16 generate
            mux: U_MX16 PORT MAP(buffer_in(1 to i+15),shift(2 to 5),out_1(i));
            and generate;

```

--input values rotated so always shift<16--

```

5
10
15
20
25
30
35
40
45
50
55

END behave;

CONFIGURATION U_SHIFT32_16_CON OF U_SHIFT32_16 IS
FOR behave
  FOR ALL:U_MX16      USE ENTITY WORK.U_MX16(behave);
  END FOR;
END FOR;
END U_SHIFT32_16_CON;

use work.dwt_types.all;
use work.dff_package.all;
use work.utils.all;

entity U_SHIFT16X16_32 IS
PORT(
  o,n : in BIT_VECTOR(1 to 16) ;
  sel : in BIT_VECTOR(1 to 4) ;
  out_1 : out BIT_VECTOR(1 to 32) );
end U_SHIFT16X16_32;

architecture behave OF U_SHIFT16X16_32 IS
SUBTYPE t_int_selector is integer range 0 to 15;
signal sel_in:t_int_selector;
BEGIN
  sel_in <= U_TO_I(sel);

  WITH sel_in SELECT
  out_1 <= n & ZERO(16) WHEN 0,
         o(1) & n & ZERO(15) WHEN 1,
         o(1 to 2) & n & ZERO(14) WHEN 2,
         o(1 to 3) & n & ZERO(13) WHEN 3,
         o(1 to 4) & n & ZERO(12) WHEN 4,
         o(1 to 5) & n & ZERO(11) WHEN 5,

```

```

5
10
15
20
25
30
35
40
45
50
55

o(1 to 6) & n & ZERO(10) WHEN 6,
o(1 to 7) & n & ZERO(9) WHEN 7,
o(1 to 8) & n & ZERO(8) WHEN 8,
o(1 to 9) & n & ZERO(7) WHEN 9,
o(1 to 10) & n & ZERO(6) WHEN 10,
o(1 to 11) & n & ZERO(5) WHEN 11,
o(1 to 12) & n & ZERO(4) WHEN 12,
o(1 to 13) & n & ZERO(3) WHEN 13,
o(1 to 14) & n & ZERO(2) WHEN 14,
o(1 to 15) & n & '0' WHEN 15;

END;
```

```

CONFIGURATION U_SHIFT16X16_32_CON OF U_SHIFT16X16_32 IS
FOR behave
END FOR;
END U_SHIFT16X16_32_CON;
```

```

use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utile.all;
```

```

entity U_HUFFMAN_DECODE IS
PORT(
mode : in t_mode ;
token_length_in : in BIT_VECTOR(1 to 2) ;
buffer_in : in BIT_VECTOR(1 to 32) ;
s : in BIT_VECTOR(1 to 5) ;
```

```

out_1 : out bit;
out_2 : out t_input;
out_3 : out BIT_VECTOR(1 to 2) );
```

```

end U_HUFFMAN_DECODE;
```

```

architecture behave OF U_HUFFMAN_DECODE IS
```

```

5
10
15
20
25
30
35
40
45
50
55

COMPONENT U_SHIFT32_16
PORT(
    buffer_in : in BIT_VECTOR(1 to 32) ;
    a : in BIT_VECTOR(1 to 5) ;
    out_1 : out BIT_VECTOR(1 to 16));
end COMPONENT;

SUBTYPE t_mag_huff is integer range 0 to 37;
SUBTYPE t_mag_out is integer range 0 to 2**(input_exp-1) -1; -- for LPP could be 9 bit unsigned

signal sel_9_12:bool;
signal in1 : BIT_VECTOR(1 to 4);
signal in2 : BIT_VECTOR(1 to 5);
signal token : BIT_VECTOR(1 to 2);
signal token_length : BIT_VECTOR(1 to 5);
signal input_decode : BIT_VECTOR(1 to 16);
signal mag_out2:t_mag_huff;
signal mag_out_huff:t_mag_huff;
signal mag_out:t_mag_out;
signal sign:bit;
BEGIN

WITH input_decode(9 to 12) SELECT
    sel_9_12 <= t WHEN B"1111" ,
    f WHEN OTHERS;

WITH sel_9_12 SELECT
    in1 <= REV(4,input_decode(13 to 16)) WHEN t ,
    REV(4,input_decode(9 to 12)) WHEN f;

WITH sel_9_12 SELECT
    in2 <= B"10110" WHEN t ,
    --add 22 to give value--
    B"00111" WHEN f,
    --add 7 to give value--

    mag_out2 <= u_to_i(in1) + u_to_i(in2);

```

```

5
10
15
20
25
30
35
40
45
50
55

mag_out_huff <= 0 WHEN input_decode(1) = '0' ELSE
1  WHEN input_decode(3) = '1' ELSE
2  WHEN input_decode(4) = '1' ELSE
3  WHEN input_decode(5) = '1' ELSE
4  WHEN input_decode(6) = '1' ELSE
5  WHEN input_decode(7) = '1' ELSE
6  WHEN input_decode(8) = '1' ELSE
   mag_out2;

--on lpf_still bit 1 is the sign bit--

sign <= input_decode(1) WHEN mode = lpf_still ELSE
   '0' WHEN mag_out_huff = 0 ELSE
   input_decode(2);

--select huff value; 0 in lpf_send or real value; rearrange the bits for real data--
--on lpf_still bit 1 is sign bit so discard--

WITH mode SELECT
mag_out <= U_TO_1(REV(9,input_decode(2 to 10))) WHEN lpf_still, --- CHECK RBV PN HERE!!!
   mag_out_huff WHEN OTHERS;

token_length <= B"000" & token_length_in;

--decode token; valid only during a token cycle--
WITH token_length(4 to 5) SELECT
token <= input_decode(1 to 2) WHEN B"10",
   input_decode(1) & B"0" WHEN B"01",
   B"00" WHEN OTHERS;

shift:U_SHIFT32_16 PORT MAP(buffer_in,s,input_decode);

--architecture outputs--
out_1 <= sign;
out_2 <= mag_out;
out_3 <= token;
-- cant happen--

```

```

5
10
15
20
25
30
35
40
45
50
55

END;

CONFIGURATION U_HUFFMAN_DECODE_CON OF U_HUFFMAN_DECODE IS
FOR behave
  FOR ALL:U_SHIFT32_16      USE ENTITY WORK.U_SHIFT32_16(behave);
  END FOR;
END FOR;
END U_HUFFMAN_DECODE_CON;

--the huffman encoder--
use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utils.all;

entity U_HUFFMAN_ENCODE IS
PORT(
  value : in t_input ;
  sign : in bit ;
  token : in BIT_VECTOR(1 to 2) ;
  mode : in t_mode ;
  cycle : in t_cycle ;
  buffer_in : in BIT_VECTOR(1 to 16) ;
  s : in BIT_VECTOR(1 to 5) ;

  out_1 : out      BIT_VECTOR(1 to 32) ;
end U_HUFFMAN_ENCODE;

architecture behave OF U_HUFFMAN_ENCODE IS
COMPONENT U_SHIFT16X16_32
PORT(
  o,n : in BIT_VECTOR(1 to 16) ;
  sel : in BIT_VECTOR(1 to 4) ;

  out_1 : out BIT_VECTOR(1 to 32) );

```

```

5
10
15
20
25
30
35
40
45
50
55

end COMPONENT;

signal header:BIT_VECTOR(1 to 2);
signal sub_value:BIT_VECTOR(1 to 11);
signal sub_value_int:integer;
signal value_bit:BIT_VECTOR(1 to 16);
signal huff_encode:BIT_VECTOR(1 to 16);
signal enc_value:BIT_VECTOR(1 to 16);
signal shift:BIT_VECTOR(1 to 32);
signal sub_const:natural;
BEGIN

--encode value--
header <= '1' & sign;

I_TO_S(value,value_bit);

WITH value SELECT
sub_const <= 7 WHEN 7 to 21,
22 WHEN 22 to 37,
0 WHEN OTHERS;

sub_value_int <= value - sub_const;

I_TO_S(sub_value_int,sub_value);

--if value is to Huffman encoded so out 16 bit as this is the max the shift removes the extra bits--
WITH value SELECT
huff_encode <= B"0"& ZERO(15) WHEN 0,
header & B"1"& ZERO(13) WHEN 1,
header & B"01"& ZERO(12) WHEN 2,
header & B"001"& ZERO(11) WHEN 3,
header & B"0001"& ZERO(10) WHEN 4,
header & B"00001"& ZERO(9) WHEN 5,
header & B"000001"& ZERO(8) WHEN 6,
header & B"000000"& REV(4,sub_value(8 to 11)) & ZERO(4) WHEN 7 to 21,

--sub 7 to give value & reverse bits --

```



```

5
10
15
20
25
30
35
40
45
50
55

        header & B"0000001111" & REV(4,sub_value(8 to 11))      WHEN 22 to 37 ,
--sub 22 to give value & reverse bits--
        header & B"0000001111111"      WHEN OTHERS;

-- final huffman encode 16 bit output value

enc_value <= token & ZERO(14)      WHEN cycle = token_cycle ELSE  --token is msb, max 2 bits--

--on intra & LPP pass thro value as 16 bit word, and reverse bit order, place sign first next to lab--

        sign & REV(15,value_bit(2 to 16))      WHEN mode = lpf_still AND cycle = data_cycle ELSE
        huff_encode      WHEN cycle = data_cycle ELSE
        ZERO(16);      -- dummy value--

shift_block: U_SHIFT16X16_32 PORT MAP(buffer_in,enc_value,s(2 to 5),shift);
--max value is 37 so 8 bits enough--

--architecture outputs--
out_1 <= shift;

END behave;

CONFIGURATION U_HUFFMAN_ENCODE_CON OF U_HUFFMAN_ENCODE IS
FOR behave
    FOR ALL:U_SHIFT16X16_32 USE ENTITY WORK.U_SHIFT16X16_32(behave);
    END FOR;
END FOR;
END U_HUFFMAN_ENCODE_CON;
--generates the new_mode from the old, and outputs control signals to the tokeniser--

use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_MODE_CONTROL IS
PORT(
    ck : in bit ;

```

```

5
10
15
20
25
30
reset : in t_reset ;
intra_inter : in t_intra ;
lpf_done : in bit ;
flags : in BIT_VECTOR(1 to 7) ;
token_in : in BIT_VECTOR(1 to 2) ;
octave : in t_octave ;
state : in t_state ;
direction : in t_direction ;
load_mode_in : in t_load ;
cycle : in t_cycle ;

out_1:out t_mode ;
out_2:out t_mode ;
out_3:out BIT_VECTOR(1 to 2) ;
out_4:out t_diff ;
out_5:out BIT_VECTOR(1 to 2) ;
out_6:out t_mode ;

```

```

end U_MODE_CONTROL ;

```

```

architecture behave OF U_MODE_CONTROL IS

```

```

--new_mode,proposed_mode,current_token,difference,token_length, --
signal nzflag:bit ;
signal origin:bit ;
signal noflag:bit ;
signal ozflag:bit ;
signal motion:bit ;
signal pro_new_z:bit ;
signal pro_no_z:bit ;
signal lpf_done_del:bit ;
signal load_mode:t_load_vec(1 to 4) ;
signal load_next:t_load ;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal pre_mode_sig:t_mode;
signal pro_mode_sig:t_mode;
signal new_mode_sig:t_mode;
signal mode:t_mode;
signal diff_sig:t_diff;
signal diff_outft_diff;
signal mode_reg:t_mode_vec(1 to 4);
BEGIN

    nzflag <= flags(1);
    origin <= flags(2);
    noflag <= flags(3);
    osflag <= flags(4);
    motion <= flags(5);
    pro_new_z <= flags(6);
    pro_no_z <= flags(7);

    prf(ck,lpf_done,lpf_done_del);

    --synchronise mode change at end of LPP--

    --the proposed value for the mode at that octave, flags etc will change this value as necessary--
    --proposed, or inherited mode from previous tree--

MODE_CONTROL:PROCESS( nzflag,origin,noflag,osflag,motion,pro_new_z,pro_no_z,lpf_done_del,token_in,direction,
    mode_reg ,state,reset,intra_inter,octave)

    variable      pro_mode :t_mode;
    variable      new_mode :t_mode;
    variable      token_out :bit_vector(1 to 2);
    variable      difference :t_diff;
    variable      token_length :bit_vector(1 to 2);
    variable      pro_flag :bit;

    BEGIN

    --initialise variables

    CASE reset IS

```

```

WHEN rat => CASE intra_inter IS
--reset on frame start, so do lpf--
    WHEN intra => pro_mode:= lpf_still,
    OTHERS => pro_mode:=lpf_send,
END CASE,
WHEN OTHERS => CASE lpf_done_del IS
    WHEN '1' => CASE intra_inter IS
        WHEN intra => pro_mode:=still,
        WHEN OTHERS => pro_mode:= send,
        END CASE,
    WHEN OTHERS => CASE state IS
        WHEN downl => pro_mode:= mode_regs(3),
        WHEN up0 => pro_mode:= mode_regs(4),
        WHEN OTHERS => CASE octave IS
            WHEN 0 =>pro_mode:=
                WHEN 1 =>pro_mode:=
                WHEN 2 =>pro_mode:=
                WHEN 3 =>pro_mode:=
            END CASE,
        END CASE,
    END CASE,
END CASE,
new_mode := pro_mode;
token_out := B"00";
difference := nodiff;
token_length := B"00";
pro_flag := '0';
CASE direction IS

```

```

5
10
15
20
25
30
35
40
45
50
55

WHEN forward =>
CASE pro_mode IS
WHEN lpf_stop|stop => null;
WHEN void => CASE ozflag IS
WHEN '1' => new_mode := stop;
WHEN OTHERS => null;
END CASE;

WHEN void_still => null;
--intra so must zero out all of tree--

WHEN still_send => token_length := B'01';
IP nzflag='1' OR pro_new_z='1' THEN token_out := B'00';
CASE ozflag IS
WHEN '1' => new_mode := stop;
WHEN OTHERS => new_mode := void;
END CASE;
ELSE token_out := B'10';
new_mode := still_send;
END IF;

WHEN send => CASE ozflag IS
WHEN '1' => token_length := B'01';
IP nzflag = '1' OR pro_new_z='1' THEN token_out := B'00';
new_mode := stop;
ELSE token_out := B'10';
new_mode := still_send;
END IF;

WHEN OTHERS => token_length := B'10';
IP (NOT(nzflag) = '1' OR motion = '1') AND NOT(nzflag) = '1'
THEN
CASE origin IS

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN '1'=> pro_flag := pro_new_z;
    WHEN OTHERS => pro_flag := pro_no_z;
        difference:= diff;
    END CASE;

    CASE pro_flag IS
    WHEN '1' => token_out := B"10";
        new_mode:= void;
    WHEN OTHERS => CASE origin IS
        WHEN '1' => token_out :=
            B"01";
            still_send;
            token_out := B"11";
            new_mode:= send;
            new_mode:=
                WHEN OTHERS =>
                    END CASE;
            ELSE
                IF (motion = '1' OR origin = '1') AND nzflag = '1'
                THEN
                    token_out := B"10";
                    new_mode:= void;
                ELSE
                    token_out := B"00";
                    new_mode:= stop;
                END IF;
            END IF;
        END CASE;

    WHEN still => token_length := B"01";
        IF nzflag = '1' OR pro_new_z = '1'

```

```

5
10
15
20
25
30
35
40
45
50
55

--zero out tree
    THEN token_out := B"00";
    new_mode := void_still;

    ELSE token_out := B"10";
    new_mode := still;
    END IF;

    WHEN lpf_still => token_out := B"00";
    token_length := B"00";

    WHEN lpf_send => difference := diff;
    token_length := B"01";

    IF noflag = '1' OR pro_no_z = '1'
    THEN
        token_out := B"00";
        new_mode := lpf_stop;
    ELSE
        token_out := B"10";
        new_mode := lpf_send;
    END IF;

--as mode stop but for this block only--
    END CASE;

    WHEN inverse =>
        CASE pro_mode IS
        WHEN lpf_stop|stop => null;

        WHEN void => CASE ozflag IS
            WHEN '1' => new_mode := stop;
            WHEN OTHERS => null;
            END CASE;

        WHEN void_still => null;

        WHEN send => CASE ozflag IS
            WHEN '1' => token_length := B"01";
            CASE token_in(1) IS

```

```

5
10
15
20
25
30
35
40
45
50
55

    WHEN '1' => new_mode := still_send;
    WHEN '0' => new_mode := stop;
    END CASE;

    WHEN OTHERS => token_length := B"10";
    CASE token_in IS
        WHEN B"11" => difference := diff;
                        new_mode := send;
        WHEN B"01" => new_mode := still_send;
        WHEN B"10" => new_mode := void;
        WHEN B"00" => new_mode := stop;
    END CASE;

    END CASE;

    WHEN still_send => token_length := B"01";
    CASE token_in(1) IS
        WHEN '1' => new_mode := still_send;
        WHEN '0' =>
            WHEN '1' => new_mode := stop;
            WHEN '0' => new_mode := void;
            CASE oflag IS
                CASE OFlag IS
                    OTHERS => new_mode := void;
            END CASE;
        END CASE;

    WHEN still => token_length := B"01";
    CASE token_in(1) IS
        WHEN '1' => new_mode := still;
        WHEN '0' => new_mode := void_still;
    END CASE;

    WHEN lpf_send => difference := diff;
    token_length := B"01";
    CASE token_in(1) IS
        WHEN '0' => new_mode := lpf_atop;
        WHEN '1' => new_mode := lpf_send;
    END CASE;

```



```

5
10
15
20
25
30
35
40
45
50
55

    WHEN lpf_still => null;
    END CASE;

END CASE;

--relate variable to corresponding signals

out_2 <= pro_mode;
pro_mode_sig <= pro_mode;
out_3 <= token_out;
out_5 <= token_length;
out_6 <= new_mode;
new_mode_sig <= new_mode;
diff_sig <= difference;

END PROCESS MODE_CONTROL;

out_1 <= mode;
out_4 <= diff_out;

pre_mode_sig <= pro_mode_sig WHEN reset = rat OR lpf_done_del= '1' ELSE
mode;

--save the new modes difference during a token cycle, when the flags and tokens are valid--
-- on lpf_still & inverse no token cycles so load on skip cycle, just so next_mode is defined

load_next <= write WHEN cycle = token_cycle ELSE
    write WHEN cycle = skip_cycle AND pro_mode_sig=lpf_still AND direction = inverse ELSE
    read;

DPP_INIT(ck,no_rat,load_next,new_mode_sig,mode);
DPP_INIT(ck,no_rat,load_next,diff_sig,diff_out);

--now write the new mode value into the mode stack at end of cycle, for later use --
--dont update modes at tree base from lpf data, on reset next[] is undefined--

```

```

5
10
15
20
25
30
35
40
45
50
55

--store base mode in mode(3)& mode(4), base changes after lpf--
load_mode <=
  (read,read,write,write) WHEN reset_rst OR lpf_done_del= '1' ELSE
  (write,write,read,read) WHEN octave= 1 AND load_mode_in= write ELSE
  (read,write,write,read) WHEN octave = 2 AND load_mode_in=write ELSE
  (read,read,read,read) ;

DFF_INIT(ck,no_rst,load_mode(1),pre_mode_sig,mode_rege(1));
DFF_INIT(ck,no_rst,load_mode(2),pre_mode_sig,mode_rege(2));
DFF_INIT(ck,no_rst,load_mode(3),pre_mode_sig,mode_rege(3));
DFF_INIT(ck,no_rst,load_mode(4),pre_mode_sig,mode_rege(4));

END behave;

CONFIGURATION MODE_CONTROL_CON OF U_MODE_CONTROL is
FOR behave
END FOR;
END MODE_CONTROL_CON;

--the tree coder chip--
--threshold = 2*quant_norm--

use work.DWT_TYPES.all;
use work.dff_package.all;
use work.utile.all;

entity U_PALMAS is
PORT(
  ck : in bit ;
  reset : in t_reset ;
  direction : in t_direction ;
  intra_inter : in t_intra ;
  channel_factor : in t_channel_factor ;
  quant_norm_1 : in t_quant ;
  quant_norm_2 : in t_quant ;
  quant_norm_3 : in t_quant ;
  quant_norm_4 : in t_quant ;

```

```

45  buffer_in : in BIT_VECTOR(1 to 16) ;
      nw_old : in t_input ;
      threshold_1 : in t_result ;
      threshold_2 : in t_result ;
      threshold_3 : in t_result ;
      threshold_4 : in t_result ;
      fifo_full, fifo_empty : in t_fifo ;
      col_length : in BIT_VECTOR(1 to xsize) ;
      row_length : in BIT_VECTOR(1 to ysize) ;
      ximage_string : in BIT_VECTOR(1 to xsize) ;
      yimage_string : in BIT_VECTOR(1 to ysize) ;
      yimage_string_3 : in BIT_VECTOR(1 to 11) ;

```

```

--old, address, rv_new, cs_new, rv_old, cs_old, buffer_out, fifo_read, fifo_write, frame done, cycle--

```

```

      out_1 : out t_input;
      out_2 : out t_memory_addr;
      out_3 : out t_load:=read;
      out_4 : out t_cs;
      out_5 : out t_load:=read;
      out_6 : out t_cs;
      out_7 : out BIT_VECTOR(1 to 16);
      out_8 : out t_load:=read;
      out_9 : out t_load:=read;
      out_10 : out bit;
      out_11 : out t_cycle);

```

```

end U_PALMAS;

```

architecture behave OF U_PALMAS IS

COMPONENT U_DECIDE

PORT(

```

      ck : in bit ;
      reset : in t_reset ;
      q_int : in t_result ;

```

```

nw,old : in t_input ;
threshold,comparison : in t_result ;
octa : in t_octave ;
load_flags : in t_load ;

```

```

out_1 : out BIT_VECTOR(1 to 7) ;
end COMPONENT ;

```

COMPONENT U_MODE_CONTROL

```

PORT(
  ck : in bit ;
  reset : in t_reset ;
  intra_inter : in t_intra ;
  lpf_done : in bit ;
  flags : in BIT_VECTOR(1 to 7) ;
  token_in : in BIT_VECTOR(1 to 2) ;
  octave : in t_octave ;
  state : in t_state ;
  direction : in t_direction ;
  load_mode_in : in t_load ;
  cycle : in t_cycle ;

```

```

  out_1 : out t_mode ;
  out_2 : out t_mode ;
  out_3 : out BIT_VECTOR(1 to 2) ;
  out_4 : out t_diff ;
  out_5 : out BIT_VECTOR(1 to 2) ;
  out_6 : out t_mode ;
end COMPONENT ;

```

COMPONENT U_ADDR_GEN

```

port(
  ck : in bit ;
  reset : in t_reset ;
  new_channel , channel : in t_channel ;

```

```

5
10
15
20
25
30
35
40
45
50
55

load_channel : in t_load ;
sub_count : in BIT_VECTOR(1 to 2) ;
col_length : in BIT_VECTOR(1 to xsize) ;
row_length : in BIT_VECTOR(1 to ysize) ;
ximage_string : in BIT_VECTOR(1 to xsize) ;
yimage_string : in BIT_VECTOR(1 to ysize) ;
yimage_string_3 : in BIT_VECTOR(1 to 11) ;
read_enable,write_enable : in bit ;
new_mode : in t_mode ;

out_1 : out t_memory_addr ;
out_2 : out t_octave ;
out_3 : out bit ;
out_4 : out bit ;
out_5 : out bit ;
out_6 : out t_state ;

end COMPONENT ;

COMPONENT U_CONTROL_COUNTER
PORT(
  ck : in bit ;
  reset : in t_reset ;
  mode,new_mode : in t_mode ;
  direction : in t_direction ;

  out_0 : out t_load ;
  out_1 : out t_cycle ;
  out_2 : out t_reset ;
  out_3 : out bit ;
  out_4 : out bit ;
  out_5 : out t_load ;
  out_6 : out t_cs ;
  out_7 : out t_load ;
  out_8 : out t_cs ;
end COMPONENT ;
-- DECIDE OUTPUTS

```

```

5
10
15
20
25
30
35
40
45
50
55

signal decide: BIT_VECTOR(1 to 7);
signal nzflag :bit;
signal orig_q :bit;
signal noflag :bit;
signal ozflag :bit;
signal motion :bit;
signal pro_no_z :bit;
signal pro_new_z :bit;

COMPONENT U_QUANT
  GENERIC (input_exp:integer:=10);
  PORT( nw,old,lev_inv : in BIT_VECTOR(1 to input_exp) ;
        sign_lev_inv : in bit ;
        direction : in t_direction ;
        q : in t_quant ;
        difference : in t_diff ;
        mode : in t_mode ;

        --pro,lev& sign
        out_1 : out t_input;
        out_2 : out BIT_VECTOR(1 to input_exp);
        out_3 : out bit);

end COMPONENT;

COMPONENT FIFO_BUFFER
  PORT(
    ck : in bit ;
    reset : in t_reset ;
    direction : in t_direction ;
    cycle : in t_cycle ;
    mode : in t_mode ;
    value,mag_out_huff : in t_input ;
    fifo_in : in BIT_VECTOR(1 to 16) ;
    fifo_full,fifo_empty : in t_fifo ;
    shift : in BIT_VECTOR(1 to 32) ;
    token_length : in BIT_VECTOR(1 to 2) ;

```

```

5
10
15
20
25
30
35
40
45
50
55

flush_buffer : in bit ;
lpf_quant : in t_quant ;
-- fifo_out, s{fifo_read fifo_write
out_1 : out BIT_VECTOR(1 to 16);
out_2 : out BIT_VECTOR(1 to 16);
out_3 : out BIT_VECTOR(1 to 16);
out_4 : out BIT_VECTOR(1 to 5);
out_5 : out t_load;
out_6 : out t_load);
end COMPONENT;

COMPONENT U_HUFFMAN_ENCODE
PORT(
    value : in t_input ;
    sign : in bit ;
    token : in BIT_VECTOR(1 to 2) ;
    mode : in t_mode ;
    cycle : in t_cycle ;
    buffer_in : in BIT_VECTOR(1 to 16) ;
    s : in BIT_VECTOR(1 to 5) ;
    out_1 : out BIT_VECTOR(1 to 32) ;
end COMPONENT;

COMPONENT U_HUFFMAN_DECODE
PORT(
    mode : in t_mode ;
    token_length_in : in BIT_VECTOR(1 to 2) ;
    buffer_in : in BIT_VECTOR(1 to 32) ;
    s : in BIT_VECTOR(1 to 5) ;

    out_1 : out bit;
    out_2 : out t_input;
    out_3 : out BIT_VECTOR(1 to 2) ;
end COMPONENT;

COMPONENT BLK_SUB_COUNT
PORT(

```

```

5
10
15
20
25
30
35
40
45
50
55

ck:in bit ;reset:in t_reset,en,cin_en,cout_en:in bit;qout bit_vector(1 to 2);carry:out bit);
end COMPONENT;

```

```

-- MODE_CONTROL outputs
signal new_mode :t_mode ;
signal pro_mode :t_mode ;
signal mode_6 :t_mode ;
signal token_out :BIT_VECTOR(1 to 2) ;
signal difference :t_diff;
signal token_length : BIT_VECTOR(1 to 2);

-- QUANT
signal pro : t_input;
signal lev_out :natural ;
signal quant_2 :BIT_VECTOR(1 to 10);
signal sign :bit ;

-- ADDR_GEN
signal address:t_memory_addr;
signal octs : t_octave;
signal sub_en : bit;
signal tree_done :bit ;
signal lpf_done : bit;
signal state : t_state;

--FIFO_BUFFER
signal buffer_out : BIT_VECTOR(1 to 16);
signal fifo_buffer_2 : BIT_VECTOR(1 to 16);
signal fifo_buffer_3 : BIT_VECTOR(1 to 16);
signal a : BIT_VECTOR(1 to 5);
signal fifo_read : t_load;
signal fifo_write : t_load;

--HUFFMAN ENCODE
signal huffman_encode: BIT_VECTOR(1 to 32);

--HUFFMAN DECODE
signal sign_in : bit;
signal lev_in : t_input;
signal token_in : BIT_VECTOR(1 to 2);
--SUB_COUNT
signal sub_count_1:BIT_VECTOR(1 to 2);

```



```

5      signal sub_count_2:bit;
--CONTROL_COUNTER
10      signal load_mode : t_load;
      signal cycle : t_cycle;
      signal decide_del : t_reset;
      signal mode_reset : t_reset;
      signal channel_reset : t_reset;
      signal channel_reset_del : t_reset;
      signal read_addr_enable : bit;
      signal write_addr_enable : bit;
      signal load_flag : t_load;
      signal cs_new : t_cs;
      signal rv_new : t_load;
      signal rv_old : t_load;
      signal cs_old : t_cs;
--
      signal load_channel:t_load;
      signal load_channel_del:t_load;
      signal new_channel:t_channel;
      signal channel:t_channel;
--
      signal flush_next : bit;
      signal flush_buffer : bit;
      signal flush_buffer_1 : bit;
      signal frame_done : bit;
--
      signal del_new:t_input;
      signal new_1:t_input;
      signal new_2:t_input;
      signal new_3:t_input;
      signal new_4:t_input;
      signal del_new_str:BIT_VECTOR(1 to 10);
      signal del_old:t_input;
      signal old_1:t_input;
      signal old_2:t_input;
      signal old_3:t_input;
      signal old_4:t_input;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal old_5:t_input;
signal del_old_str:BIT_VECTOR(1 to 10);
signal dumb_vhd:BIT_VECTOR(1 to 32);
signal lev_in_str:BIT_VECTOR(1 to 10);
signal decide_reset:t_reset;
signal threshold_oct:t_result;
signal quant_oct:t_quant;
signal data_out:t_input;

signal one:bit:='1';

BEGIN

    nzflag <= decide(1);
    origin <= decide(2);
    noflag <= decide(3);
    ozflag <= decide(4);
    motion <= decide(5);
    pro_no_z <= decide(7); --pro_no_z or pro_new_z--
    pro_new_z <= decide(6);

    lev_out <= u_TO_I(quant_2); --corresponding level--

    rw_new <= read;

    --change channel--
    load_channel <= write WHEN sub_en='1' AND sub_count_2 = '1' ELSE read;
    channel_reset <= rst WHEN load_channel=write ELSE no_rst;

    DFL(ck,channel_reset,channel_reset_del);
    -- reset mode at start of new channel
    mode_reset <= rst WHEN reset = rst OR channel_reset_del = rst ELSE no_rst;

    new_channel <= y WHEN channel_factor= luminance ELSE
        u WHEN channel = y ELSE
        v WHEN channel = u ELSE
        y;

```

```

5
10
15
20
25
30
35
40
45
50
55

--flush the buffer in the huffman encoder--
flush_next <= '1' WHEN channel_factor = luminance AND load_channel = write ELSE
    '1' WHEN channel_factor = color AND channel = v AND load_channel = write ELSE
    '0';

DF1(ck, flush_next, flush_buffer);

DF1(ck, flush_buffer, flush_buffer_1);
DF1(ck, flush_buffer_1, frame_done);

DF1(ck, nw, new_1);
DF1(ck, new_1, new_2);
DF1(ck, new_2, new_3);
DF1(ck, new_3, new_4);
DF1(ck, new_4, del_new);

I_TO_S(del_new, del_new_str);

DF1(ck, old, old_1);
DF1(ck, old_1, old_2);
DF1(ck, old_2, old_3);
DF1(ck, old_3, old_4);
DF1(ck, old_4, old_5);
--old has variable delays for inverse--
del_old <= old_5 WHEN direction = forward ELSE
    old_5 WHEN direction = inverse AND
    (pro_mode = send OR pro_mode = still_send OR pro_mode = lpf_send OR pro_mode = void) ELSE
    old_2;

I_TO_S(del_old, del_old_str);

I_TO_S(lev_in, lev_in_str);

WITH reset SELECT
decide_reset <= rst WHEN
    decide_del WHEN OTHERS;

threshold_oct <= threshold_4 WHEN pro_mode = lpf_still OR pro_mode = lpf_send OR pro_mode = lpf_stop ELSE

```

```

threshold_1 WHEN octs=0 AND channel=y ELSE
threshold_2 WHEN (octs=1 AND channel=y) OR (octs=0 AND (channel=u OR channel=v)) ELSE
threshold_3;

quant_oct <= quant_norm_4 WHEN pro_mode = lpf_still OR pro_mode = lpf_send OR pro_mode = lpf_stop ELSE
quant_norm_1 WHEN octs=0 AND channel=y ELSE
quant_norm_2 WHEN (octs=1 AND channel=y) OR (octs=0 AND (channel=u OR channel=v)) ELSE
quant_norm_3;

dec_map:U_DECIDE PORT MAP(ck,decide_reset,threshold_oct,nw,old,threshold_oct,threshold_oct,
octs,load_flags,decide);

mode_map:U_MODE_CONTROL
PORT MAP(ck,mode_reset,intra_inter,lpf_done,decide,token_in,octs,state,direction,load_mode,cycle,
new_mode,pro_mode,token_out,difference,token_length,mode_6);

quant_map:U_QUANT_GENERIC MAP(input_exp) PORT MAP(del_new_str,del_old_str,lev_in,sign_in,
direction,quant_oct,difference,pro_mode, pro_quant_2,sign);

addr_map:U_ADDR_GEN PORT MAP(ck,reset,new_channel,load_channel,sub_count_1,col_length,row_length,
ximage_string,yimage_string,yimage_string_3,read_addr_enable,write_addr_enable,new_mode,
address,octs,sub_en,tree_done,lpf_done,state);

fifo_map:FIFO_BUFFER PORT MAP (ck,reset,direction,cycle,pro_mode,lev_out,lev_in,buffer_in,
fifo_full, fifo_empty,huffman_encode, token_length, flush_buffer,quant_norm_4,
buffer_out,fifo_buffer_2,fifo_buffer_3,s,fifo_read ,fifo_write );

huff_enc:U_HUFFMAN_ENCODE PORT MAP (lev_out ,sign,token_out,pro_mode,cycle,fifo_buffer_2,s,huffman_encode);

dumb_vhdl <= fifo_buffer_2 & fifo_buffer_3;

huff_dec:U_HUFFMAN_DECODE PORT MAP (pro_mode,token_length,dumb_vhdl,s,
sign_in,lev_in,token_in);

one <= '1';

sub_map:BLK_SUB_COUNT PORT MAP (ck,reset,sub_en,one,one,sub_count_1,sub_count_2);

```

```

5
10
15
20
25
30
35
40
45
50
55

control_map:U_CONTROL_COUNTER PORT MAP(ck,reset,pro_mode,new_mode,direction
,load_mode,cycle,decide_del,read_addr_enable,write_addr_enable,load_flags,cs_new,rw_old,cs_old);

DFF_INIT(ck,no_reset,load_channel,new_channel,channel);

--architecture outputs--

WITH new_mode SELECT
data_out <= 0 WHEN void|void_still,
           pro WHEN OTHERS;

out_1 <= data_out;
out_2 <= address;
out_3 <= rw_new;
out_4 <= cs_new;
out_5 <= rw_old;
out_6 <= cs_old;
out_7 <= buffer_out;
out_8 <= fifo_read;
out_9 <= fifo_write;
out_10 <= frame_done;
out_11 <= cycle;

END behave;

CONFIGURATION U_PALMAS_CON OF U_PALMAS IS
FOR behave
FOR ALL:U_DECIDE USE ENTITY WORK.U_DECIDE(behave);
END FOR;

FOR ALL:U_MODE_CONTROL USE ENTITY WORK.U_MODE_CONTROL(behave);
END FOR;

FOR ALL:U_ADDR_GEN USE ENTITY WORK.U_ADDR_GEN(behave);
END FOR;

FOR ALL:U_QUANT USE ENTITY WORK.U_QUANT(behave);

```

```

5
10
15
20
25
30
35
40
45
50
55

END FOR;

FOR ALL: FIFO_BUFFER USE ENTITY WORK.FIFO_BUFFER(bhaved);
END FOR;

FOR ALL: U_HUFFMAN_ENCODE USE ENTITY WORK.U_HUFFMAN_ENCODE(bhaved);
END FOR;

FOR ALL: U_HUFFMAN_DECODE USE ENTITY WORK.U_HUFFMAN_DECODE(bhaved);
END FOR;

FOR ALL: BLK_SUB_COUNT USE ENTITY WORK.BLK_SUB_COUNT(bhaved);
END FOR;

FOR ALL: U_CONTROL_COUNTER USE ENTITY WORK.U_CONTROL_COUNTER(bhaved);
END FOR;

END FOR;
END U_PALMAS_CON;
use work.DWT_TYPES.all;

--now the selector fn to shift the level depending on q--

entity U_BARREL_SHIFT_RIGHT IS
  GENERIC (n: Integer);
  PORT (
    q : in t_quant ;
    data: in BIT_VECTOR(1 to n) ;

    out_1 : out BIT_VECTOR(1 to n) );
  end U_BARREL_SHIFT_RIGHT;

architecture behave of U_BARREL_SHIFT_RIGHT IS
  --the bahift for the inverse, to generate the rounded level --
  BEGIN
  WITH q SELECT

```

```

5
10
15
20
25
30
35
40
45
50
55

out_1 <= data WHEN 0,
    b"0" & data(1 to n-1) WHEN 1,
    b"00" & data(1 to n-2) WHEN 2,
    b"000" & data(1 to n-3) WHEN 3,
    b"0000" & data(1 to n-4) WHEN 4,
    b"00000" & data(1 to n-5) WHEN 5,
    b"000000" & data(1 to n-6) WHEN 6,
    b"0000000" & data(1 to n-7) WHEN 7;
END behave;

configuration behift_right_con of U_BARREL_SHIFT_RIGHT is
    FOR behave
        END for;
    end behift_right_con;

use work.DWT_TYPES.all;

entity U_BARREL_SHIFT_LEFT is
    GENERIC (n:integer);
    PORT(
        q : in t_quant;
        data : in BIT_VECTOR(1 to n) );
    out_1 : out BIT_VECTOR(1 to n) );
    end U_BARREL_SHIFT_LEFT;

architecture behave of U_BARREL_SHIFT_LEFT is
    --the function to return the quantised level(UNSIGNED), and proposed value given,--
    -- the newold values,inverse direction --
BEGIN
    WITH q SELECT
        out_1 <= data WHEN 0,
            data(2 to n) & b"0" WHEN 1,
            data(3 to n) & b"01" WHEN 2,
            data(4 to n) & b"011" WHEN 3,

```

```

5
10
15
20
25
30
35
40
45
50
55

    data(5 to n) & b"0111" WHEN 4,
    data(6 to n) & b"0111" WHEN 5,
    data(7 to n) & b"01111" WHEN 6,
    data(8 to n) & b"011111" WHEN 7;
END behave;

configuration behift_left_con of U_BARREL_SHIFT_LEFT is
    FOR behave
        END for;
    end behift_left_con;

    use work.DWT_TYPES.all;
    use work.UTILS.all;

    entity U_QUANT IS
        GENERIC (input_exp:integer:=10);
        PORT( nw,old,lev_inv : in BIT_VECTOR(1 to input_exp) ;
            sign_lev_inv : in bit ;
            direction : in t_direction ;
            q : in t_quant ;
            difference : in t_diff ;
            mode : in t_mode ;

            --pro,lev& sign
            out_1 : out t_input;
            out_2 : out BIT_VECTOR(1 to input_exp);
            out_3 : out bit);
    end U_QUANT;

architecture behave OF U_QUANT IS

    COMPONENT U_BARREL_SHIFT_LEFT
        GENERIC (n:integer);
        PORT(

```



```

5
10
15
20
25
    q : in t_quant ;
    data : in BIT_VECTOR(1 to n) ;
    out_1 : out BIT_VECTOR(1 to n) ;
end COMPONENT;

COMPONENT U_BARREL_SHIFT_RIGHT
  GENERIC (n:integer);
  PORT(
    q : in t_quant ;
    data: in BIT_VECTOR(1 to n) ;

    out_1 : out BIT_VECTOR(1 to n) ;
  end COMPONENT;

```

```

    signal no_bits:BIT_VECTOR(1 to input_exp+1);
    signal abs_no_bits:BIT_VECTOR(1 to input_exp+1);
    signal lev_no:BIT_VECTOR(1 to input_exp+1);
    signal lev_data:BIT_VECTOR(1 to input_exp+1);
    signal lev:BIT_VECTOR(1 to input_exp+1);
    signal zero_factor:BIT_VECTOR(1 to input_exp+1);
    signal round_lev:BIT_VECTOR(1 to input_exp+1);
    signal r_lev:BIT_VECTOR(1 to input_exp+1);
    signal minus_round_lev:BIT_VECTOR(1 to input_exp+1);
    signal pro_no_bits:BIT_VECTOR(1 to input_exp+1);
    signal pro_new_bits:BIT_VECTOR(1 to input_exp+1);
    signal sgn_level:bit;
    signal lev_z :bit;
    signal inv_lev_z :bit;
    signal sub_in :BIT_VECTOR(1 to input_exp);
    -- actually input_exp+1/2 bits
    signal pro_no:integer;
    signal pro_new:integer;
    signal round_lev_int:natural;
    signal no:integer;

```

5
10
15
20
25
30
35
40
45
50
55

```

signal abs_no:natural;
BEGIN
    --decide which of new-old or new will be quantised, and the sign of the level--
    --level is stored in sign & magnitude form--
    --put old=0, when new is to be quantised

    WITH difference SELECT
    sub_in <= old WHEN diff,
    ZERO(input_exp) WHEN notdiff;

    --sign & magnitude of value to be quantised
    no <= S_TO_I(nw) - S_TO_I(sub_in);
    abs_no <= ABS(no);    --now input_exp+1 bits--

    -- convert to bits--
    I_TO_S(no,no_bits);
    I_TO_S(abs_no,abs_no_bits);

    WITH direction SELECT
    sgn_level <= no_bits(1) WHEN forward,
    sign_level_inv WHEN inverse;

    --find the quant. level by shifting by q, for the inverse it comes from the Huffman decoder--

    be_right: U_BARREL_SHIFT_RIGHT_GENERIC_MAP(input_exp+1) PORT MAP(q,abs_no_bits,lev_data);

    --saturate the lev at 37; for the Huffman table,except in lpf_still mode; send all the bits--
    lev_no <= lev_data WHEN mode = lpf_still ELSE
    B"00000100101" WHEN U_TO_I(lev_data)>37 ELSE
    lev_data;

    --pick level from lev_no or lev_inv depending on the direction, still input_exp+1 bits

```

```

5
10
15
20
25
30
35
40
45
50
55

WITH direction SELECT
  lev <= lev_no WHEN forward,
    8*0 & lev_inv WHEN inverse;

--the level <= 0 flag--
lev_z <= '1' WHEN U_TO_I(lev) = 0 ELSE
  '0';

WITH lev_z SELECT
  inv_lev_z <= '0' WHEN '1',
    '1' WHEN OTHERS;

--the level value shifted up, and rounded--
bs_left: U_BARREL_SHIFT_LEFT GENERIC MAP(input_exp+1) PORT MAP(q,lev,r_lev);

--if lev=0 out all 0's--
--clear out extra bit for lpf still case--

WITH mode SELECT
  zero_factor <= 8*00 & ALL_SAME(input_exp-1,'1') WHEN lpf_still,
    ALL_SAME(input_exp+1,inv_lev_z) WHEN OTHERS,
  round_lev <= r_lev AND zero_factor;

--calculate the proposed value
round_lev_int <= U_TO_I(round_lev);

WITH sgn_level SELECT
  pro_no <= S_TO_I(old) + round_lev_int WHEN '0',
    S_TO_I(old) - round_lev_int WHEN '1';

--now pro_new = +/- round_lev--

WITH sgn_level SELECT
  pro_new <= round_lev_int WHEN '0',
    - round_lev_int WHEN '1'; --NEG --

```

```

5
10
15
20
25
30
35
--outputs--
-- shift to input_exp bits for output
I TO S(pro_new,pro_new_bits);
I TO S(pro_no,pro_no_bits);

WITH difference SELECT
out_1 <= S_TO_I( pro_no_bits(3 to input_exp+2)) WHEN diff,
      S_TO_I(pro_new_bits(2 to input_exp+1)) WHEN modiff;

out_2 <= lev(2 to input_exp+1);

out_3 <= sgn_level;

END behave;

CONFIGURATION QUANT_CON OF U_QUANT IS
FOR behave

FOR ALL: U_BARREL_SHIFT_RIGHT USE CONFIGURATION WORK.bshift_right_con;
END FOR;

FOR ALL: U_BARREL_SHIFT_LEFT USE CONFIGURATION WORK.bshift_left_con;
END FOR;

END FOR;
END QUANT_CON;
use work.DWT_TYPES.all;

-- returns a signal with n copies of the zero
package util is

FUNCTION ZERO ( CONSTANT n:NATURAL) RETURN BIT_VECTOR;

-- returns a signal with n copies of the input bit
FUNCTION ALL_SAME (CONSTANT n:NATURAL; s:bit) RETURN BIT_VECTOR;

-- reverses the bit order

```

```

5
10
15
20
25
30
35
40
45
50
55
FUNCTION REV (CONSTANT n:natural; in_in:BIT_VECTOR) RETURN BIT_VECTOR;

```

```

end utils;

```

```

package body utils is

```

```

FUNCTION ALL_SAME (CONSTANT n:natural; s:bit) RETURN BIT_VECTOR IS

```

```

variable out_b:BIT_VECTOR(1 to n);

```

```

BEGIN

```

```

for i IN 1 to n LOOP

```

```

out_b(i):= s;

```

```

END LOOP;

```

```

RETURN out_b;

```

```

END ALL_SAME;

```

```

FUNCTION ZERO (CONSTANT n:natural) RETURN BIT_VECTOR IS

```

```

variable out_b:BIT_VECTOR(1 to n);

```

```

BEGIN

```

```

for i IN 1 to n LOOP

```

```

out_b(i):='0';

```

```

END LOOP;

```

```

RETURN out_b;

```

```

END ZERO;

```

```

FUNCTION REV (CONSTANT n:natural; in_in:BIT_VECTOR) RETURN BIT_VECTOR IS

```

```

variable temp:BIT_VECTOR(1 to n);

```

```

BEGIN

```

```

for i IN 1 to n LOOP

```

```

temp(i):=in_in(n-i+ in_in'left);

```

```

END LOOP;

```

```

RETURN temp;

```

```

END;

```

```

END utils;

```

```

--VHDL Description of CONTROL_COUNTER--
--a counter to control the sequencing ofw, token, huffman cycles--
--decide reset, enable 1 cycle early, and latched to avoid glitches--
--lpf_stop is a dummy mode to disable the block writes huffman data--
--cycles for that block--
use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_CONTROL_COUNTER is
  port(
    ck : in bit ;
    reset : in t_reset ;
    mode,new_mode : in t_mode ;
    direction : in t_direction ;

    out_0 : out t_load;
    out_1 : out t_cycle;
    out_2 : out t_reset;
    out_3 : out bit;
    out_4 : out bit;
    out_5 : out t_load;
    out_6 : out t_cs;
    out_7 : out t_load;
    out_8 : out t_cs) ;

    --mode load,cycle,decide reset,read_addr_enable,write_addr_enable,load flag--
    --decode write_addr_enable early and latch to avoid feedback loop with pro_mode--
    --in MODE_CONTROL--
    and U_CONTROL_COUNTER;

architecture behave of U_CONTROL_COUNTER is
  component COUNT_SYNC
    generic (n:integer);
    port(
      ck:in bit ;
      reset:in t_reset;

```

```

5
10
15
20
25
30
35
40
45
50
55

en:in bit;
q:out bit_vector(1 to n);
carry:out bit);
end COMPONENT;

signal write_del:bit;
signal write_sig:bit;
signal decide_del:bit;
signal decide_sig:bit;
signal count_len:bit;
signal count_1:bit_vector(1 to 4);
signal count_2:bit;
signal always_one:bit;

BEGIN

count_len <= u_to_1(count_1);

control:PROCESS(ck,count_reset,direction,mode,new_mode,count_len)

VARIABLE
cycle : t_cycle;
decide_reset : t_reset;
load_mode : t_load;
load_flags : t_load;
cs_new : t_cs;
cs_old : t_cs;
rw_old : t_load;
read_addr_enable : bit;
write_addr_enable : bit;

BEGIN

cycle := skip_cycle;
decide_reset := no_rst;
load_mode := read;
load_flags := read;
cs_new := no_sel;
cs_old := sel;

```

```

5
10
15
20
25
30
35
40
45
50
55

    rv_old := read;
    read_addr_enable := '0';
    write_addr_enable := '0';

CASE direction IS
WHEN forward =>
    CASE mode IS
    WHEN send|still_send|lpf_send =>
        CASE count_len IS
        WHEN 0 to 3 => read_addr_enable := '1';
            ce_new := sel;
        WHEN 4 => cycle := token_cycle;
            load_flags := write;
            write_addr_enable := '1';
        WHEN 5 to 7 => write_addr_enable := '1';
        CASE new_mode IS
        WHEN stop|lpf_stop => cycle := skip_cycle;
            rv_old := read;
            ce_old := no_sel;
        WHEN vld => cycle := skip_cycle;
            rv_old := write;
        WHEN OTHERS => cycle := data_cycle;
            rv_old := write;
        END CASE;
    WHEN 8 => decide_reset := rst;
        CASE new_mode IS
        WHEN stop|lpf_stop => cycle := skip_cycle;
            rv_old := read;
            ce_old := no_sel;
        WHEN vld => cycle := skip_cycle;
            load_mode := write;
            rv_old := write;
        WHEN OTHERS => cycle := data_cycle;
            load_mode := write;
            rv_old := write;
        END CASE;
    WHEN OTHERS => null;

```



```

5
10
15
20
25
30
35
40
45
50
55

END CASE;

WHEN still =>
    CASE count_len IS
        WHEN 0 to 3 => read_addr_enable := '1';
            cs_new := sel;
        WHEN 4 => cycle := token_cycle;
            write_addr_enable := '1';
            load_flags := write;
        WHEN 5 to 7 => rw_old := write;
            write_addr_enable := '1';
        CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
        END CASE;

        WHEN 8 => decide_reset := rst;
            rw_old := write;
            load_mode := write;
        CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
        END CASE;

        WHEN OTHERS => null;
    END CASE;

WHEN lpf_still =>
    CASE count_len IS
        WHEN 0 to 3 => read_addr_enable := '1';
            cs_new := sel;
        WHEN 4 => cycle := token_cycle;
            write_addr_enable := '1';
            load_flags := write;
        WHEN 5 to 7 => cycle := data_cycle;
            rw_old := write;
        WHEN 8 => cycle := data_cycle;
            rw_old := write;
    END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

    decide_reset:= rst;
    load_mode:= write;

    WHEN OTHERS => null;
    END CASE;

    CASE count_len IS
    WHEN 0 to 3 => read_addr_enable := '1';
        cs_new:= sel;
    WHEN 4 => load_flag:= write;
        cycle:= token_cycle;

    WHEN 5 to 7 => write_addr_enable := '1';
        write_addr_enable := '1';

    CASE new_mode IS
    WHEN stop => rw_old := read;
        cs_old:= no_sel;
    WHEN OTHERS => rw_old := write;
    END CASE;

    WHEN 8 => decide_reset := rst;
    CASE new_mode IS
    WHEN stop => rw_old := read;
        cs_old:= no_sel;
    WHEN OTHERS => load_mode := write;
        rw_old:= write;
    END CASE;

    WHEN OTHERS => null;
    END CASE;

    WHEN void_still => CASE count_len IS
    WHEN 0 => write_addr_enable := '1';

    WHEN 1 to 3 => write_addr_enable := '1';
        rw_old:= write;
    WHEN 4 => rw_old := write;
        load_mode:= write;

    --allow for delay--

```

```

5
10
15
20
25
30
35
40
45
50
55

        decide_reset := ret;
        WHEN OTHERS => null;
        END CASE;

        WHEN OTHERS => null;
        END CASE;

        WHEN inverse =>
            CASE mode IS
                WHEN send|still_send|lpf_send =>
                    CASE count_len IS
                        WHEN 0 to 3 => read_addr_enable := '1';
                        WHEN 4 => cycle := token_cycle;
                            write_addr_enable := '1';
                            load_flags := write;
                    WHEN 5 to 7 => write_addr_enable := '1';
                    CASE new_mode IS
                        WHEN stop|lpf_stop => cycle := skip_cycle;
                            rw_old := read;
                            cs_old := no_sel;
                        WHEN void => cycle := skip_cycle;
                            rw_old := write;
                        WHEN OTHERS => cycle := data_cycle;
                            rw_old := write;
                    END CASE;
                WHEN 8 => decide_reset := ret;
                CASE new_mode IS
                    WHEN stop|lpf_stop => cycle := skip_cycle;
                        rw_old := read;
                        cs_old := no_sel;
                    WHEN void => cycle := skip_cycle;
                        load_mode := write;
                        rw_old := write;
                    WHEN OTHERS => cycle := data_cycle;
                        load_mode := write;
                        rw_old := write;
                    END CASE;
                WHEN OTHERS => null;
            END CASE;
        END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

--skip to allow reset in huffman--
    WHEN still =>
        END CASE;
        CASE count_len IS
        WHEN 0 => null ;

        WHEN 1 => cycle := token_cycle;
            write_addr_enable := '1';
        WHEN 2 to 4 => rw_old := write;
            write_addr_enable := '1';
            CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
            END CASE;

        WHEN 5 => rw_old:=write;
            decide_reset:= rst;
            load_mode:= write;
            CASE new_mode IS
            WHEN void_still => cycle := skip_cycle;
            WHEN OTHERS => cycle := data_cycle;
            END CASE;

        WHEN OTHERS => null;
        END CASE;

    WHEN lpf_still =>
        CASE count_len IS
        WHEN 0 => null ;

        WHEN 1 => write_addr_enable := '1';

        WHEN 2 to 4 => cycle := data_cycle;
            rw_old:= write;
            write_addr_enable := '1';
        WHEN 5 => cycle := data_cycle;
            rw_old:= write;
            decide_reset:= rst;
            load_mode:= write;

        WHEN OTHERS => null;
        END CASE;

        CASE count_len IS
        WHEN void =>

```

```

--match with previous--

```

```

--skip for write enb delay--

```

```

5
10
15
20
25
30
35
40
45
50
55

--dummy token cycle for mode update--
WHEN 0 to 3 => read_addr_enable := '1';
WHEN 4 => load_flags := write;
cycle:= token_cycle;

WHEN 5 to 7 => write_addr_enable := '1';
CASE new_mode IS
WHEN stop => rv_old := read;
ce_old:= no_sel;
WHEN OTHERS => rv_old := write;
END CASE;

WHEN 8 => decide_reset := rst;
CASE new_mode IS
WHEN stop => rv_old := read;
ce_old:= no_sel;
WHEN OTHERS => load_mode := write;
rv_old:= write;
END CASE;

WHEN OTHERS => null;
END CASE;

WHEN void_still => CASE count_len IS
WHEN 0 => null;

WHEN 1 => write_addr_enable := '1';

WHEN 2 to 4 => write_addr_enable := '1';
rv_old:= write;

WHEN 5 => rv_old := write;
load_mode:= write;
decide_reset:= rst;

WHEN OTHERS => null;
END CASE;

WHEN OTHERS => null;
END CASE;

END CASE;

write_sig <= write_addr_enable;

```

```

5
10
15
20
25
30
35
40
45
50
55

decide_sig <= decide_reset;

DPR(ck,reset,write_sig,write_del);
out_0 <= load_mode;
out_1 <= cycle;
out_2 <= decide_sig;
out_3 <= read_addr_enable;
out_4 <= write_del;
out_5 <= load_flags;
out_6 <= cs_new;
out_7 <= rw_old;
out_8 <= cs_old;

END PROCESS;

WITH reset SELECT
count_reset <= rst WHEN rst,
               decide_sig WHEN OTHERS;

control_cnt: count_sync GENERIC MAP(4) PORT MAP(ck,count_reset,always_one,count_1,count_2);

END behave;

CONFIGURATION CONTROL_COUNTER_CON OF U_CONTROL_COUNTER IS
FOR behave
  FOR ALL:count_sync USE ENTITY WORK.count_sync(behave);
  END FOR;
END FOR;
END CONTROL_COUNTER_CON;

```

```

5
10
15
20
25
30
35
40
45
50
55

--VHDL Description of MODE_CONTROL--
--generates the new_mode from the old, and outputs control signals to the tokeniser--

use work.DWT_TYPES.all;
use work.dff_package.all;

entity U_MODE_CONTROL IS
PORT(
    ck : in bit ;
    reset : in t_reset ;
    intra_inter : in t_intra ;
    lpf_done : in bit ;
    flag : in BIT_VECTOR(1 to 7);
    token_in : in BIT_VECTOR(1 to 2) ;
    octave : in t_octave ;
    state : in t_state ;
    direction : in t_direction ;
    load_mode_in : in t_load ;
    cycle : in t_cycle ;

    out_1:out t_mode;
    out_2:out t_mode;
    out_3:out BIT_VECTOR(1 to 2) ;
    out_4:out t_diff;
    out_5:out BIT_VECTOR(1 to 2) ;
    out_6:out t_mode);

end U_MODE_CONTROL;

architecture behave of U_MODE_CONTROL IS

--new_mode,proposed mode,current token,difference,token_length, ---
signal nzflag:bit;

```

```

5
10
15
20
25
30
35
40
45
50
55

signal origin:bit;
signal noflag:bit;
signal ozflag:bit;
signal motion:bit;
signal pro_new_z:bit;
signal pro_no_z:bit;
signal lpf_done_del:bit;
signal load_mode:t_load_vec(1 to 4);
signal load_next:t_load;
signal pre_mode_sig:t_mode;
signal pro_mode_sig:t_mode;
signal new_mode_sig:t_mode;
signal mode:t_mode;
signal diff_sig:t_diff;
signal diff_out:t_diff;
signal mode_regs:t_mode_vec(1 to 4);
BEGIN

    nzflag <= flags(1);
    origin <= flags(2);
    noflag <= flags(3);
    ozflag <= flags(4);
    motion <= flags(5);
    pro_new_z <= flags(6);
    pro_no_z <= flags(7);

    DFL(ck,lpf_done,lpf_done_del);

    --synchronise mode change at end of LPF--

    --the proposed value for the mode at that octave, flags etc will change this value as necessary--
    --proposed, or inherited mode from previous tree--

    MODE_CONTROL:PROCESS( nzflag,origin,noflag,ozflag,motion,pro_new_z,pro_no_z,lpf_done_del,token_in,direction,
                           mode_regs ,state,reset,intra_inter,octave)

        variable      pro_mode :t_mode;
        variable      new_mode :t_mode;
        variable      token_out :bit_vector(1 to 2);

```



```

5
10
15
20
25
30
35
40
45
50
55

variable      difference :t_diff;
variable      token_length ibit_vector(1 to 2);
variable      xpro_flag ibit;

BEGIN
--initialise variables

CASE reset IS
WHEN rst => CASE intra_inter IS
--reset on frame start, so do lpf--
WHEN intra => pro_mode:= lpf_still;
WHEN OTHERS => pro_mode:=lpf_send;
END CASE;
WHEN OTHERS => CASE lpf_done_del IS
WHEN '1' => CASE intra_inter IS
WHEN intra => pro_mode:=still;
WHEN OTHERS => pro_mode:= send;
END CASE;
WHEN OTHERS => CASE state IS
WHEN down1 => pro_mode:= mode_regs(3);
WHEN up0 => pro_mode:= mode_regs(4);
WHEN OTHERS => CASE octave IS
WHEN 0 =>pro_mode:=
WHEN 1 =>pro_mode:=
WHEN 2 =>pro_mode:=
WHEN 3 =>pro_mode:=
END CASE;
END CASE;

END CASE;

END CASE;

mode_regs(1);
mode_regs(2);
mode_regs(3);
mode_regs(4);

--jump sideways in oct 1--

```

```

5
10
15
20
25
30
35
40
45
50
55

END CASE;

% new_mode := pro_mode;
  token_out := B"00";
  difference := nodiff;
  token_length := B"00";
  pro_flag := '0';

--inherit the previous mode--

CASE direction IS
  WHEN forward =>

    CASE pro_mode IS
      WHEN lpf_stop/stop => null;
      WHEN void => CASE ozflag IS
        WHEN '1' => new_mode := stop;
        WHEN OTHERS => null;
      END CASE;

      WHEN void_still => null;

      --intra so must zero out all of tree--
      WHEN still_send => token_length := B"01";
        IF nzflag='1' OR pro_new_z='1' THEN token_out := B"00";
        CASE ozflag IS
          WHEN '1' => new_mode := stop;
          WHEN OTHERS => new_mode := void;
        END CASE;
        ELSE token_out := B"10";
          new_mode := still_send;
        END IF;

      WHEN send => CASE ozflag IS
        WHEN '1' => token_length := B"01";
        IF nzflag = '1' OR pro_new_z='1' THEN token_out := B"00";

```

```

5
10
15
20
25
30
35
40
45
50
55

new_mode:= stop;
token_out := B"10";
new_mode:= still_send;

END IF;

WHEN OTHERS => token_length := B"10";

IF(NOT(noflag) = '1' OR motion = '1') AND NOT(nzflag) = '1'
THEN
CASE origin IS
WHEN '1'=> pro_flag := pro_new_z;
WHEN OTHERS => pro_flag := pro_no_z;
difference:= diff;
END CASE;

CASE pro_flag IS
WHEN '1' => token_out := B"10";
new_mode:= void;
WHEN OTHERS => CASE origin IS
WHEN '1' => token_out :=
new_mode:=
WHEN OTHERS =>

END CASE;

END CASE;

END CASE;

ELSE
IF (motion = '1' OR origin = '1')AND nzflag = '1'
THEN
token_out := B"10";
new_mode:= void;
ELSE
token_out := B"00";

```

new_mode:= stop;

END IF;

END IF;

END CASE;

WHEN still => token_length := B'01';
IF nflag = '1' OR pro_new_z = '1'
THEN token_out := B'00';
new_mode:= void_still;

--zero out tree--

ELSE token_out := B'10';
new_mode:= still;

END IF;

WHEN lpf_still => token_out := B'00';
token_length:= B'00';

WHEN lpf_send => difference := diff;
token_length:= B'01';

IF nflag = '1' OR pro_no_z = '1'
THEN token_out := B'00';
new_mode:= lpf_stop;
ELSE token_out := B'10';
new_mode:= lpf_send;

--as mode stop but for this block only--

END IF;

END CASE;

WHEN inverse =>

CASE pro_mode IS

WHEN lpf_stop/stop => null;

WHEN void => CASE ozflag IS
WHEN '1' => new_mode := stop;

```

5
10
15
20
25
30
35
40
45
50
55

      WHEN OTHERS => null;
      END CASE;

      WHEN void_still => null;

      WHEN send => CASE ozflag IS
--repeat of still-send code--
        WHEN '1' => token_length := B'01';

        CASE token_in(1) IS
          WHEN '1' => new_mode := still_send;
          WHEN '0' => new_mode := stop;
        END CASE;

        WHEN OTHERS => token_length := B'10';

        CASE token_in IS
          WHEN B'11' => difference := diff;
                        new_mode := send;
          WHEN B'01' => new_mode := still_send;
          WHEN B'10' => new_mode := void;
          WHEN B'00' => new_mode := stop;
        END CASE;

      END CASE;

      WHEN still_send => token_length := B'01';
      CASE token_in(1) IS
        WHEN '1' => new_mode := still_send;
        WHEN '0' =>
          CASE ozflag IS
            WHEN '1' => new_mode := stop;
            WHEN OTHERS => new_mode := void;
          END CASE;
        END CASE;

      WHEN still => token_length := B'01';
      CASE token_in(1) IS
        WHEN '1' => new_mode := still;
        WHEN '0' => new_mode := void_still;
      END CASE;

```

```

5
10
15
20
25
30
35
40
45
50
55

END CASE;

WHEN lpf_send => difference := diff;
token_length := B'01';
CASE token_in(1) IS
WHEN '0' => new_mode := lpf_stop;
WHEN '1' => new_mode := lpf_send;
END CASE;

WHEN lpf_still => null;
END CASE;

END CASE;

--relate variable to corresponding signals

out_2 <= pro_mode;
pro_mode_sig <= pro_mode;
out_3 <= token_out;
out_5 <= token_length;
out_6 <= new_mode;
new_mode_sig <= new_mode;
diff_sig <= difference;

END CASE;

END PROCESS MODE_CONTROL;

out_1 <= mode;
out_4 <= diff_out;

pre_mode_sig <= pro_mode_sig WHEN reset = rat OR lpf_done_del = '1' ELSE
mode;

--save the new mode's difference during a token cycle, when the flag and tokens are valid--
-- on lpf_still & inverse no token cycles so load on skip cycle, just so next_mode is defined

load_next <= write WHEN cycle = token_cycle ELSE

```

```

write WHEN cycle = skip_cycle AND pro_mode_sig=lpf_still AND direction = inverse ELSE
  read;

```

```

DFF_INIT(ck,no_rst,load_next,new_mode_sig,mode);
DFF_INIT(ck,no_rst,load_next,diff_sig,diff_out);

```

```

--now write the new mode value into the mode stack at end of cycle, for later use --
--dont update modes at tree base from lpf data, on reset next(1) is undefined--
--store base mode in mode(3)& mode(4), base changes after lpf--

```

```

load_mode <= (read,read,write,write) WHEN reset=rst OR lpf_done=del='1' ELSE
  (write,write,read,read) WHEN octave= 1 AND load_mode_in= write ELSE
  (read,write,write,read) WHEN octave = 2 AND load_mode_in=write ELSE
  (read,read,read,read) ;

```

```

DFF_INIT(ck,no_rst,load_mode(1),pre_mode_sig,mode_regs(1));
DFF_INIT(ck,no_rst,load_mode(2),pre_mode_sig,mode_regs(2));
DFF_INIT(ck,no_rst,load_mode(3),pre_mode_sig,mode_regs(3));
DFF_INIT(ck,no_rst,load_mode(4),pre_mode_sig,mode_regs(4));

```

```

END behave;

```

```

CONFIGURATION MODE_CONTROL_CON OF U_MODE_CONTROL Is
FOR behave
END FOR;
END MODE_CONTROL_CON;

```

Claims

1. A method comprising the step of:
 using a number of intercoupled accumulators, each accumulator comprising at least an adder circuit
 and a storage circuit, to filter a sequence of input image data values with both a first digital filter having
 X coefficients and a second digital filter having Y coefficients into a sequence of transformed image data
 values, Y being substantially equal to the number of intercoupled accumulators, wherein $X < T$.
2. A method as claimed in claim 1 wherein a first data value of the sequence of transformed image data
 values is an output of the first digital filter and wherein a second data value of the sequence of transformed
 image data values is an output of the second digital filter.
3. A method as claimed in claim 1 or 2 wherein the first digital filter is a boundary filter.
4. A method comprising the steps of:
 multiplying a first data value of a sequence of data values by a first plurality of predetermined values
 to generate a first plurality of products;
 inputting a selected one of the plurality of products into a first input of a first accumulator;
 multiplying a second data value of the sequence of data values by a second plurality of predeter-
 mined values to generate a second plurality of products;
 inputting a selected one of the second plurality of products into a first input of a second accumu-
 lator, an output of the first accumulator being supplied to a second input of the second accumulator;
 multiplying a third data value of the sequence of data values by a third plurality of predetermined
 values to generate a third plurality of products;
 inputting a selected one of the third plurality of products into a first input of a third accumulator, an
 output of the second accumulator being supplied to a second input of the third accumulator;
 multiplying a fourth data value of the sequence of data values by at least one predetermined value
 to generate at least one fourth product; and
 inputting the at least one fourth product into a first input of a fourth accumulator, an output of the
 third accumulator being supplied to a second input of the fourth accumulator, wherein each of the plurality
 of predetermined values is selected from coefficients of a quasi-perfect reconstruction filter and coeffi-
 cients of a boundary quasi-perfect reconstruction filter.
5. A method as claimed in claim 4 wherein each of the first, second, third and fourth accumulators comprises
 at least an adder circuit and a storage circuit, each of the storage circuits storing a number A of partially
 transformed data values during processing of one octave of a decomposition, and each of the storage
 circuits storing a number B of partially transformed data values during processing of another octave of a
 decomposition, wherein A is not equal to B.
6. A method as claimed in claim 4 or 5 wherein the sequence of data values comprises data values of rows
 and columns of a two-dimensional image, the method transforming both the rows and the columns of data
 values into a sub-band decomposition.
7. A method comprising the steps of:
 multiplying a first data value of a sequence of data values by a first plurality of predetermined values
 to generate a first plurality of products;
 inputting a selected one of the plurality of products into a first input of a first accumulator;
 multiplying a second data value of the sequence of data values by a second plurality of predeter-
 mined values to generate a second plurality of products;
 inputting a selected one of the second plurality of products into a first input of a second accumu-
 lator, an output of the first accumulator being supplied to a second input of the second accumulator;
 multiplying a third data value of the sequence of data values by a third plurality of predetermined
 values to generate a third plurality of products;
 inputting a selected one of the third plurality of products into a first input of a third accumulator, an
 output of the second input of the third accumulator;
 multiplying a fourth data value of the sequence of data values by at least one predetermined value
 to generate at least one fourth product; and
 inputting the at least one fourth products into a first input of a fourth accumulator, an output of the
 third accumulator being supplied to a second input of the fourth accumulator, wherein at least some of

the plurality of predetermined values are selected from coefficients of a boundary perfect reconstruction filter and coefficients of a boundary quasi-perfect reconstruction filter.

8. A method as claimed in claim 7 wherein the output of the second accumulator is supplied to a second input of the first accumulator, wherein the output of the third accumulator is supplied to a third input of the second accumulator, and wherein the output of the fourth accumulator is supplied to a third input of the third accumulator.
9. A method comprising the steps of:
 - multiplying a first data value of a sequence of data values by a first plurality of predetermined values to generate a first plurality of products;
 - inputting a selected one of the plurality of products into a first input of a first accumulator, the first accumulator having a first output terminal;
 - multiplying a second data value of the sequence of data values by a second plurality of predetermined values to generate a second plurality of products;
 - inputting a selected one of the second plurality of products into a first input of a second accumulator, an output of the first accumulator being supplied to a second input of the second accumulator, the output of the second accumulator being supplied to a second input of the first accumulator;
 - multiplying a third data value of the sequence of data values by a third plurality of predetermined values to generate a third plurality of products;
 - inputting a selected one of the third plurality of products into a first input of a third accumulator, an output of the second accumulator being supplied to a second input of the third accumulator, the output of the third accumulator being supplied to a third input of the second accumulator;
 - multiplying a fourth data value of the sequence of data values by at least one predetermined value to generate at least one fourth product; and
 - inputting the at least one fourth product into a first input of a fourth accumulator, an output of the third accumulator being supplied to a second input of said fourth accumulator, the output of the fourth accumulator being supplied to a third input of the third accumulator, the fourth accumulator having a second output terminal.
10. A method comprising the steps of:
 - reading a sequence of data values from a plurality of memory locations, each of the data values being stored in a separate one of the plurality of memory locations; and
 - overwriting some of the memory locations one at a time into a sequence of transformed data values of a sub-band decomposition.
11. A wavelet transform circuit arranged for transforming a sequence of a number C of input image data values adjacent an image boundary into a corresponding sequence of the number C of transformed image data values, the transformed image data values consisting of alternating low and high pass transformed image data values.
12. A wavelet transform circuit for transforming a sequence of data values adjacent a boundary of the data values, the sequence comprising a boundary subsequence adjacent the boundary and a non-boundary subsequence, the circuit comprising:
 - a multiplier circuit having an input for receiving an input data value and a plurality outputs;
 - a first multiplexer having a plurality of data inputs each coupled to selected ones of the plurality of outputs of the multiplier circuit, at least one control input, and a data output,
 - a first accumulator circuit having a first data input coupled to the data output of the first multiplexer, a second data input, a plurality of control inputs, and a data output,
 - a second multiplexer having a plurality of data inputs each coupled to selected ones of the plurality of outputs of the multiplier circuit, at least one control input, and a data output,
 - a second accumulator circuit having a first data input coupled to the data output of the second multiplexer, a second data input coupled to the data output of the first accumulator circuit, a third data input, a plurality of control inputs, and a data output,
 - a third multiplexer having a plurality of data inputs, at least one control input, and a data output;
 - a third accumulator circuit having a first data input coupled to the data output of the third multiplexer, a second data input coupled to the data output of the second accumulator circuit, a third data input, a plurality of control inputs, and a data output,

a fourth accumulator circuit having a first data input, a second data input coupled to one of the plurality of outputs of the multiplier circuit, a plurality of control inputs, and a data output coupled to the third data input of the third accumulator circuit, and

a control circuit coupled to the control inputs of the first, second and third multiplexers and the first, second, third and fourth accumulator circuits, the control circuit controlling the first, second, third multiplexers and the first, second, third and fourth accumulator circuits to transform the boundary subsequence using a first digital filter and to transform the non-boundary subsequence using a second digital filter.

- 10 13. A method of addressing selected ones of a plurality of memory locations, the plurality of memory locations storing a plurality of image data values, the method comprising the steps of:
addressing a sequence of a plurality of memory locations in a first period of time;
transforming at least some of the sequence into a first octave of sub-band decomposition during the first period of time;
15 addressing a subsequence of the sequence of the plurality of memory locations in a second period of time; and
transforming at least some of the subsequence from the first octave into a second octave during the second period of time, the first and second octaves comprising at least part of a sub-band decomposition.
- 20 14. A method as claimed in claim 13 wherein the step of addressing a sequence comprises adding one of a first value and a second value to a first address to generate a second address, the first and second addresses of the sequence being addresses of successive ones of the sequence of memory locations, and wherein the step of addressing a subsequence comprises the step of adding one of a third value and a fourth value to a first address of the subsequence of memory locations to generate a second address of the subsequence of memory locations, the first and second addresses of the subsequence being addresses of successive ones of the subsequence of memory locations.
- 25 15. A method as claimed in claim 13 wherein the step of addressing a sequence comprises subtracting one of a first value and a second value from a first address to generate a second address, the first and second addresses of the sequence being addresses of successive ones of the sequence of memory locations, and wherein the step of addressing a subsequence comprises the step of subtracting one of a third value and a fourth value from a first address of the subsequence of memory locations to generate a second address of the subsequence of memory locations, the first and second addresses of the subsequence being addresses of successive ones of the subsequence of memory locations.
- 30 16. A method of addressing selected ones of a plurality of memory locations storing a plurality of data values of a sub-band decomposition comprising a number of octaves, the method comprising the step of:
generating a second address by adding or by subtracting one of a first number and a second number to a first address, the first number having a value which is octave-dependent, the second number having another value which is octave-dependent, and the first and second addresses being successive addresses of respective ones of said plurality of memory locations.
- 35 17. A circuit for generating an address of a memory location of a plurality of memory locations storing a plurality of data values of a sub-band decomposition, the circuit comprising:
45 a first counter which increments by an octave-dependent first variable number of counts,
a multiplier circuit having a first data input coupled to an output of the first counter, a second data input coupled to receive a value indicative of a size of an image, and a data output,
a second counter which increments by an octave-dependent second variable number of counts,
an adder circuit having a first data input coupled to the output of the multiplier circuit, a second data input coupled to an output of the second counter, and a data output on which the address is generated.
- 50 18. A circuit as claimed in claim 17 wherein the first counter receives a clock input signal and increments by the first variable number of counts in one clock cycle of the clock input signal, and wherein the second counter receives a clock input signal and increments by the second variable number of counts in one clock cycle of the clock input signal.
- 55 19. A state machine for generating an address of a memory location of a plurality of memory locations storing a plurality of data values of a sub-band decomposition, the data values comprising luminance data values,

the state machine receiving:

a first signal upon which the state machine changes states,

a second signal indicative of an octave of the sub-band decomposition, and

a third signal indicative of a whether the address being generated is a luminance data value, the

state machine changing state based upon the first, second, and third signals.

20. A state machine as claimed in claim 19 wherein the state machine also receives a fourth signal indicative of a whether the address being generated is a chrominance data value, the state machine changing state based upon the fourth signal.

21. A circuit for generating an address and comprising:

an accumulator circuit having a first data input for receiving a first increment value, a second data input for receiving a second increment value, a third data input, and a data output, said accumulator circuit, and

a storage element having a data input coupled to the data output of the accumulator circuit, and a data output coupled to the third data input of the accumulator circuit, the address being output on the data output of the storage element.

22. a circuit as claimed in claim 22 wherein the accumulator circuit has a fourth data input receiving an address base-offset chrominance value, and a fifth data input receiving an address base-offset luminance value.

23. A circuit comprising:

a flag generating circuit which receives a plurality of data values of a new block and a plurality of data values of an old block and generates a plurality of different flags, the new block being a block of a sub-band composition of a new frame, the old block being a block of a sub-band decomposition of an old frame, the different flags being selected from flags consisting of ORIGIN, NEW_Z, NO_Z, MOTION, NO-FLAG; and

a stage machine which assumes a given new state based on a current stage of the state machine and also on the different flags, the state machine outputting a token indicating the new state of the stage machine, the new stage being selected from VOID, SEND, STILL, SEND-STILL, STOP.

24. A method comprising the steps of:

using a digital circuit to low pass filter a first sequence of a first number of image data values into a second sequence of a second number of image values, the second number being less than the first number, and

using the digital circuit to low pass filter the second sequence of image data values to generate a sub-band decomposition.

25. A method as claimed in claim 24 wherein the digital circuit is a low pass forward transform perfect reconstruction digital filter.

26. A method as claimed in claim 26 wherein the digital circuit comprises a convolver.

27. A method comprising the steps of:

reading a number of first data values from a first number of rows of memory locations of a memory, the number of first data values comprising a part of a sub-band decomposition;

transforming or inverse transforming the number of first data values to generate a number of second data values, the number of second data values comprising high pass component data values and low pass component data values; and

after the transforming step, writing the number of second data values to a second number of rows of memory locations of the memory, the first number of rows being equal to the second number of rows.

28. A method as claimed in claim 28 wherein the rows of memory locations read in the reading step are different rows from the rows of memory locations written in the writing step or are the same rows of memory locations that are written in the writing step.

29. A method as claimed in claim 27 or 28 wherein the number of first data values equal the number of second data values.

30. A method as claimed in claim 27, 28 or 28 wherein the memory is a dynamic memory.

31. A method comprising the steps of:

reading a number of first image data values from a first number of twos of memory locations of a memory;

transforming the number of first data values to generate a number of second data values of a sub-band decomposition, the number of second data values comprising high pass component data values and low pass component data values; and

after the transforming step, writing the number of second data values to a second number of rows of memory locations of the memory, the first number of rows being equal to the second number of rows.

32. A method comprising the steps of:

reading a first block of low pass component data values and associated first high pass component trees of blocks of data values from a first plurality of dynamic memory locations, the first plurality of dynamic memory locations containing a plurality of blocks of low pass component data values, and each of the blocks of low pass component data values having three associated high pass component trees of blocks of data values,

writing the first block of low pass component data values and associated first high pass component trees of blocks of data values read from the first plurality of dynamic memory locations into a second memory;

reading a second block of low pass component data values and associated second high pass component trees of blocks of data values from a second plurality of dynamic memory locations, the second plurality of dynamic memory locations containing a plurality of blocks of low pass component data values, and each of the blocks of low pass component data values having three associated high pass component trees of blocks of data values,

writing the second block of low pass component data values and associated second high pass component trees of blocks of data values read from the second plurality of dynamic memory locations into a fourth memory, and

processing the data values written into the second and fourth memories to generate a compressed image data stream.

33. A method as claimed in claim 32 wherein the second memory is a plurality of static random access memory locations, and wherein the fourth memory is another plurality of static random access memory locations.

34. A method as claimed in claim 32 or 33 wherein the first block of low pass component data values and associated first high pass component trees of blocks of data values consist of 256 memory locations, and wherein the first block of low pass component data values and associated first high pass component trees of blocks of data values comprise a part of a three octave sub-band decomposition.

35. A method as claimed in claim 32, 33 or 34 wherein the processing step generates a third block of low pass component data values and associated third high pass component trees of blocks of data values, the method comprising the steps of:

writing the third block of low pass component data values and associated third high pass component trees of blocks of data values into the fourth memory;

reading a fourth block of low pass component data values and associated fourth high pass component trees of blocks of data values from the second plurality of dynamic memory locations, the dynamic memory locations from which the fourth block is read being different dynamic memory locations from the dynamic memory locations from which said second block was read; and

after the step of reading the fourth block, writing the third block of low pass component data values and associated third high pass component trees of blocks of data values into the second plurality of dynamic memory locations.

36. A method as claimed in claim 32, 33 or 34 wherein the processing step generates a third block of low pass component data values and associated third high pass component trees of blocks of data values, the method comprising the steps of:

writing the third block of low pass component data values and associated third high pass component trees of blocks of data values into the fourth memory, and

reading the third block of low pass component data values and associated third high pass compo-

nent trees of blocks of data values from the fourth memory and writing the third block of low pass component data values and associated third high pass component trees of blocks of data values into the second plurality of dynamic memory locations.

37. A method comprising the steps of:

reading a first block of low pass component data values and associated first high pass component trees of blocks of data values from a first plurality of dynamic memory locations, the first plurality of dynamic memory locations containing a plurality of blocks of low pass component data values, each of the blocks of low pass component data values having three associated high pass component trees of blocks of data values,

writing the first block of low pass component data values and associated first high pass component trees of blocks of data values into a second memory,

receiving an encoded data stream and using the encoded data stream to generate a second block of low pass component data values and associated second high pass component trees of blocks of data values;

writing the second block of low pass component data values and associated second high pass component trees of blocks of data values into the second memory; and

38. A method comprising the steps of:

during a first period of time, using a plurality of memory locations of static random access memory as a line delay in a convolver, and

during a second period of time, using the plurality of memory locations to store a block of low pass component data values and associated high pass component trees of blocks of data values.

39. A system comprising:

a first circuit which receives a input data stream and which generates quantized data stream and an inverse quantized data stream, the input data stream comprising tokens and transformed image data values, and

a second circuit which receives the quantized data stream and which generates an inverse quantized data stream and which generates an inverse quantized data stream, the inverse quantized data stream generated by the second circuit being identical to the inverse quantized data stream generated by the first circuit.

40. A system as claimed in claim 30 wherein the first circuit comprises:

means for generating a mode control signal from the input data stream,

quantizer means for receiving the transformed image data values of the input data stream and the mode control signal and for quantizing said transformed image data values into the quantized data stream using one of a first quantization method and a second quantization method depending on the mode control signal, the quantizer means also being for generating the inverse quantized data stream from the quantized data stream, and

means for generating a compressed data stream from the quantized data stream,

wherein the second circuit comprises:

means for generating a mode control signal from the compressed data stream,

means for generating quantized data values from the compressed data stream, and

quantizer means for receiving said quantized data values of the compressed data stream and the mode control signal and for inverse quantizing the quantized data values into a data stream.

41. A circuit for generating an address for traversing a sub-band decomposition, the circuit comprising:

a first counter receiving a first count enable signal, and generating a first plurality of count output signals and a first carry out signal,

a second counter receiving a second count enable signal, and generating a second plurality of count output signals and a second carry out signal,

a third counter receiving a third count enable signal, and generating a third plurality of count output signals and a third carry out signal,

a fourth counter receiving a fourth count enable signal and generating a fourth plurality of counter output signals and a fourth carry out signal, and

means for receiving the first, second, third, and fourth carry out signals and for generating the first, second, third and fourth count enable signals, the first, second, third and fourth plurality of count output

signals comprising the address.

42. A circuit as claimed in claim 41 wherein the sub-band decomposition comprises three high pass sub-bands, and wherein the first plurality of count output signals is indicative of an X address of a root of a tree, the second plurality of count output signals is indicative of a Y address of the root of the tree, the tree comprises a plurality of blocks of data values, the third plurality of count output signals is indicative of an address of data values in a block of data values of the tree, and the fourth plurality of counter output signals is indicative of the sub-band of the tree.

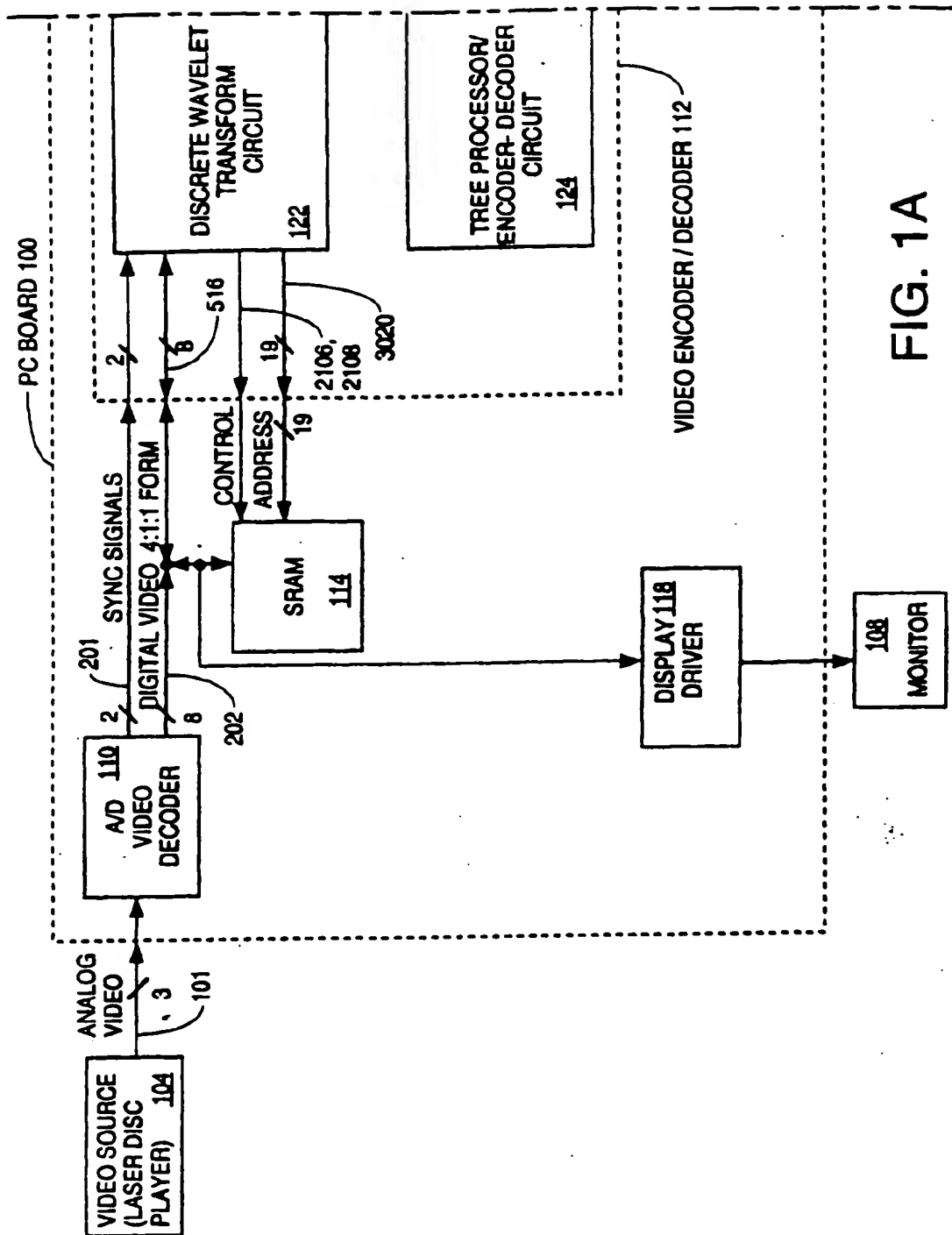
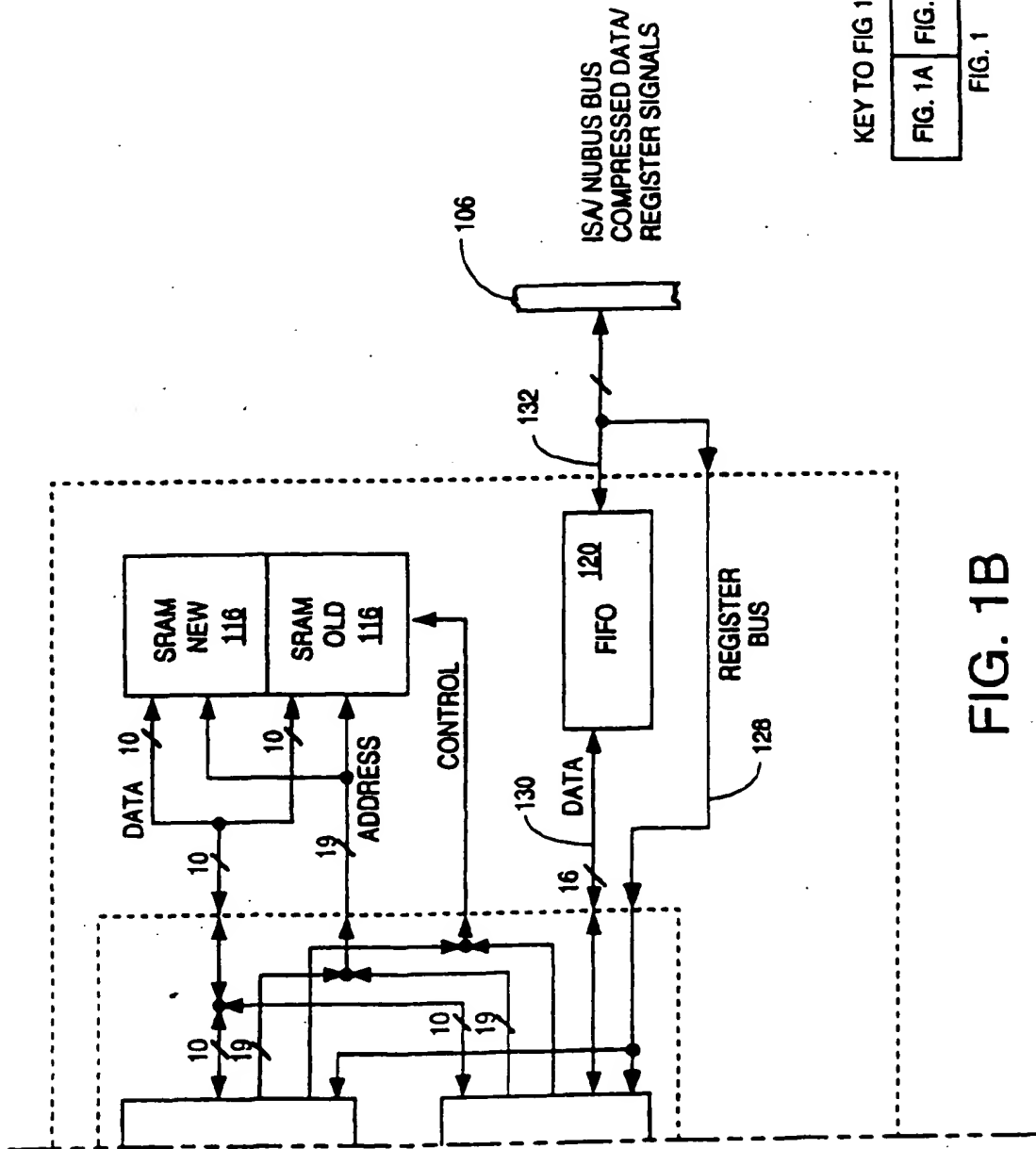


FIG. 1A



KEY TO FIG 1

FIG. 1A | FIG. 1B

FIG. 1A | FIG. 1B

FIG. 1

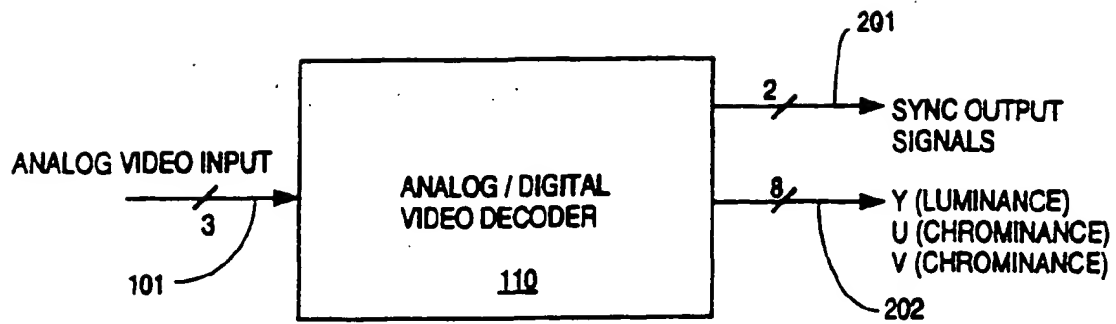


FIG. 2

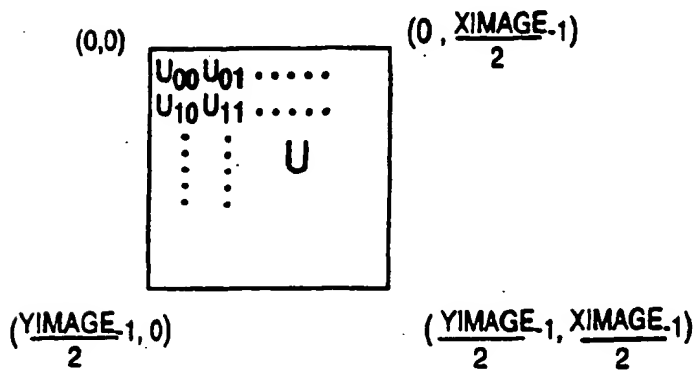
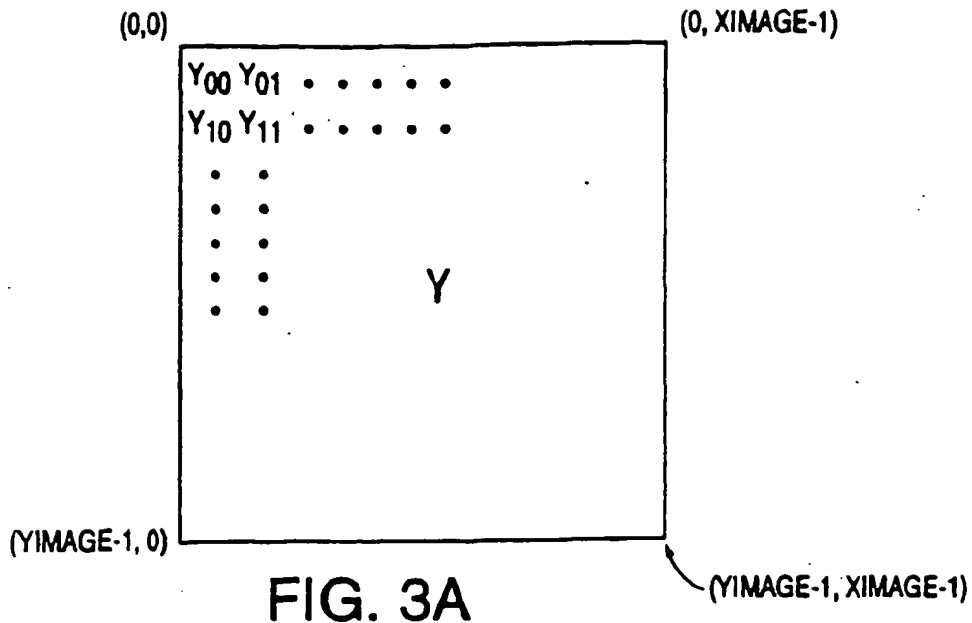


FIG. 3B

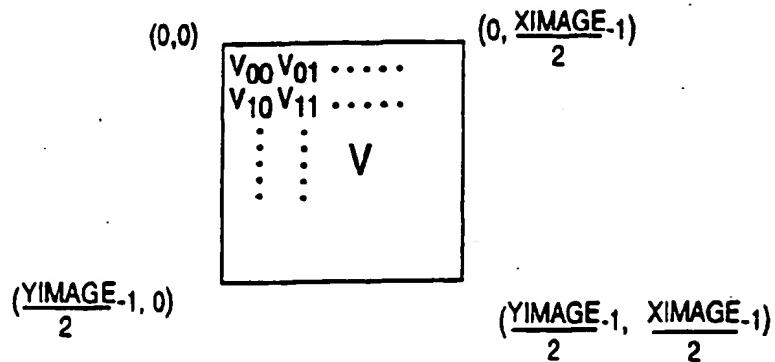


FIG. 3C

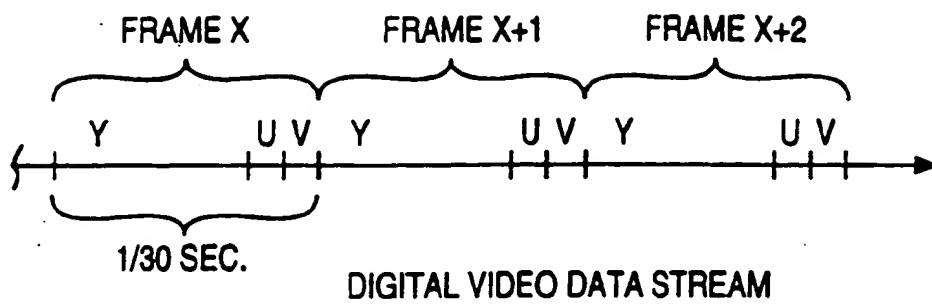


FIG. 4

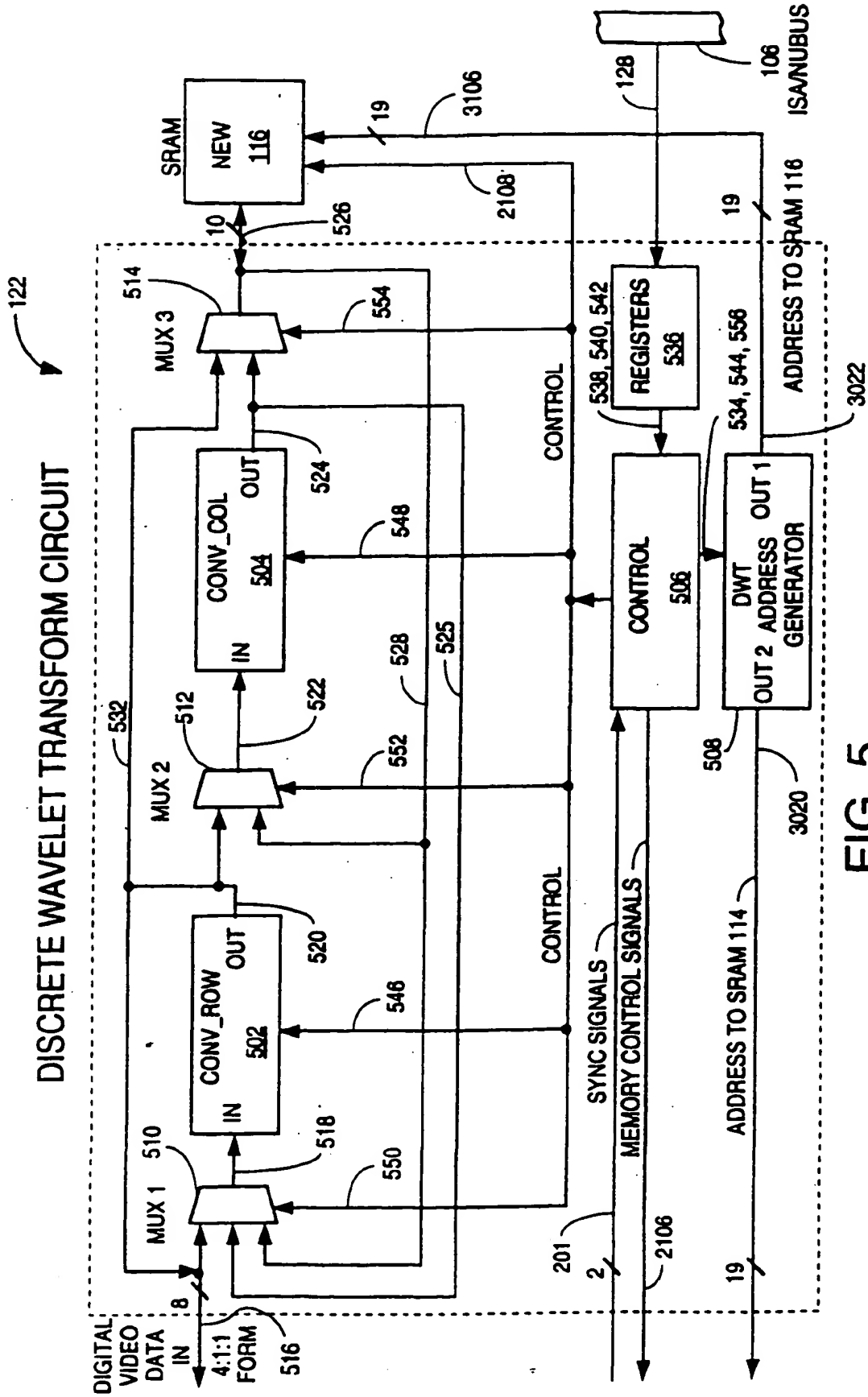


FIG. 5

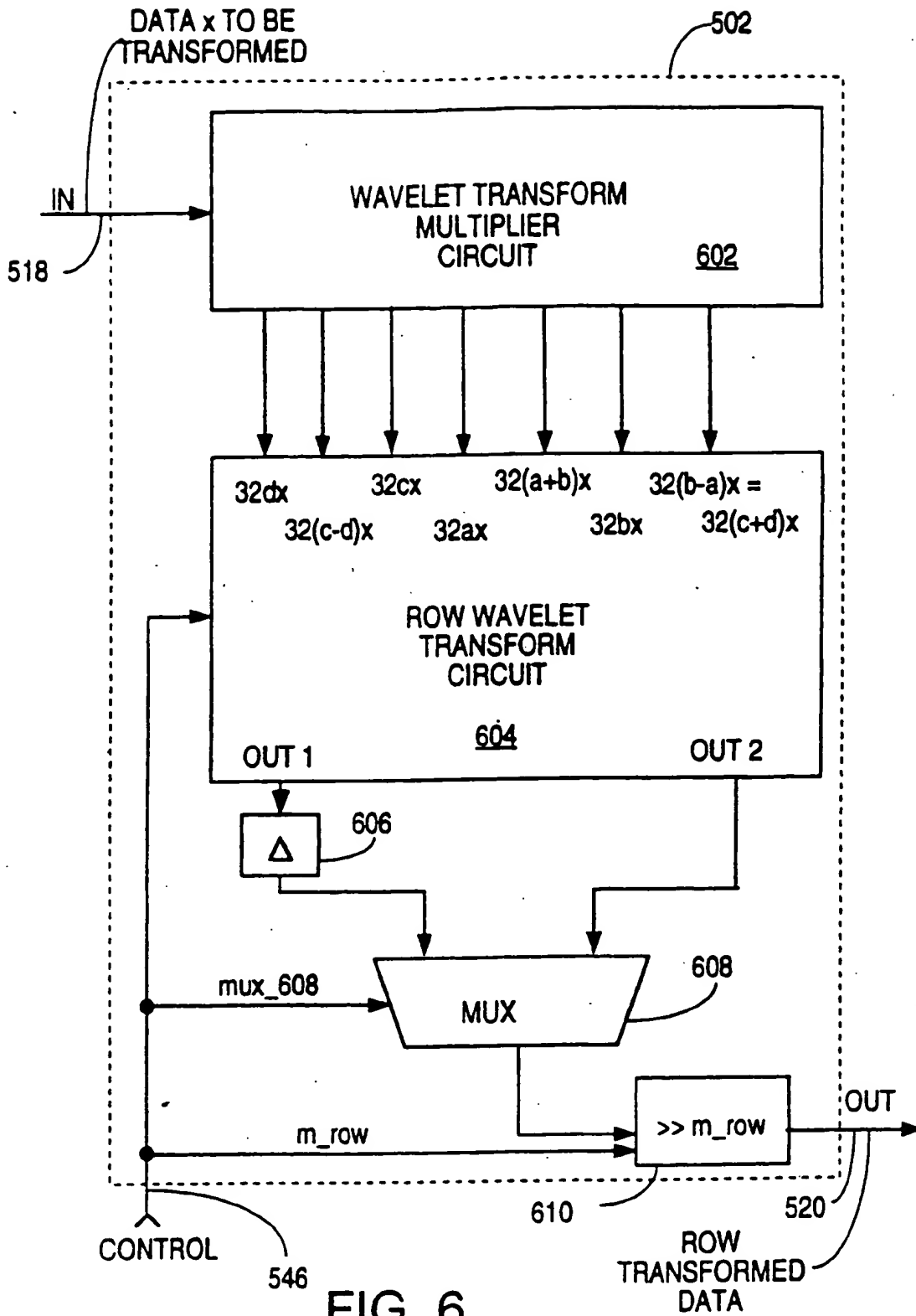


FIG. 6
CONV_ROW
BLOCK DIAGRAM

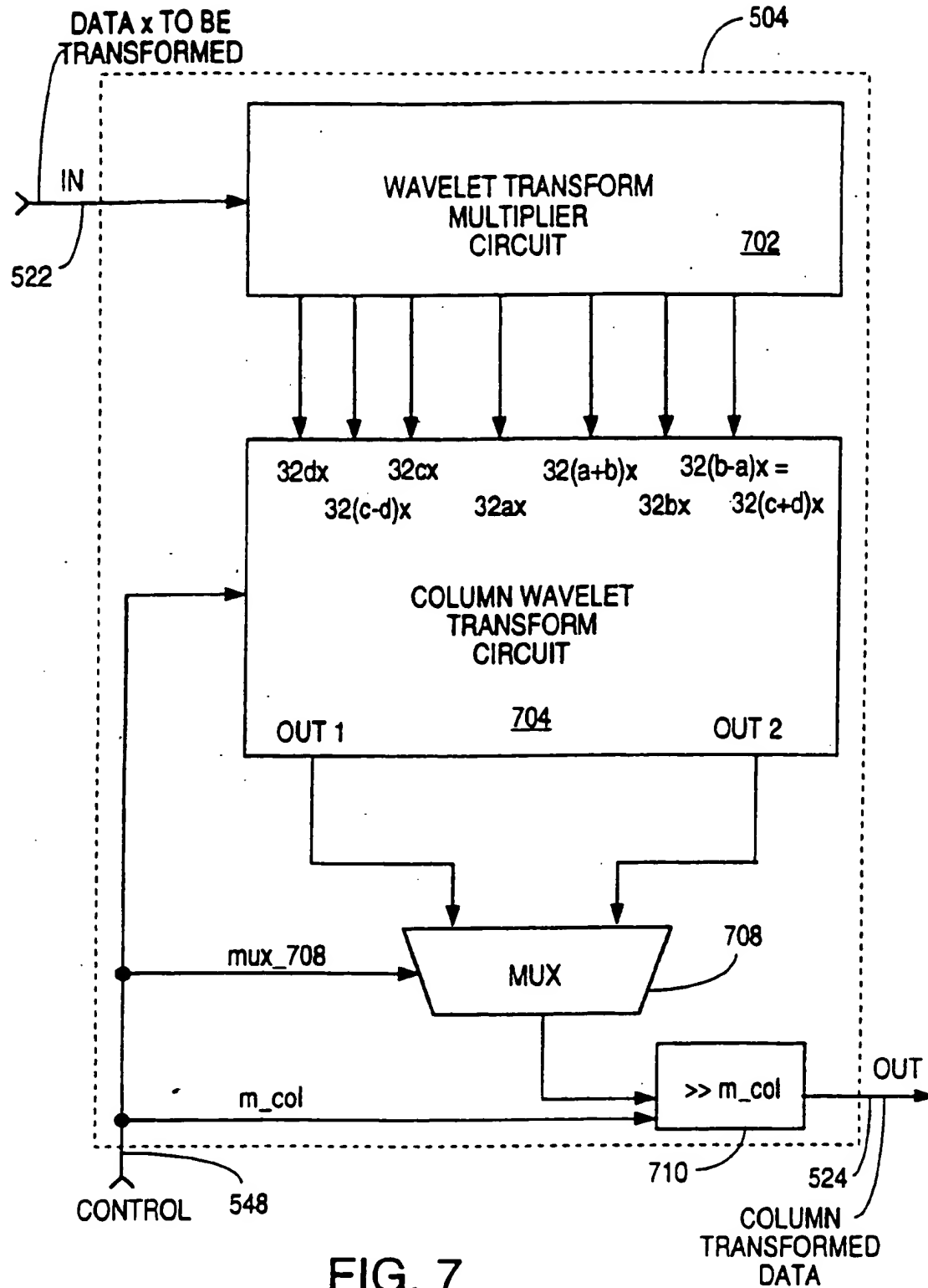
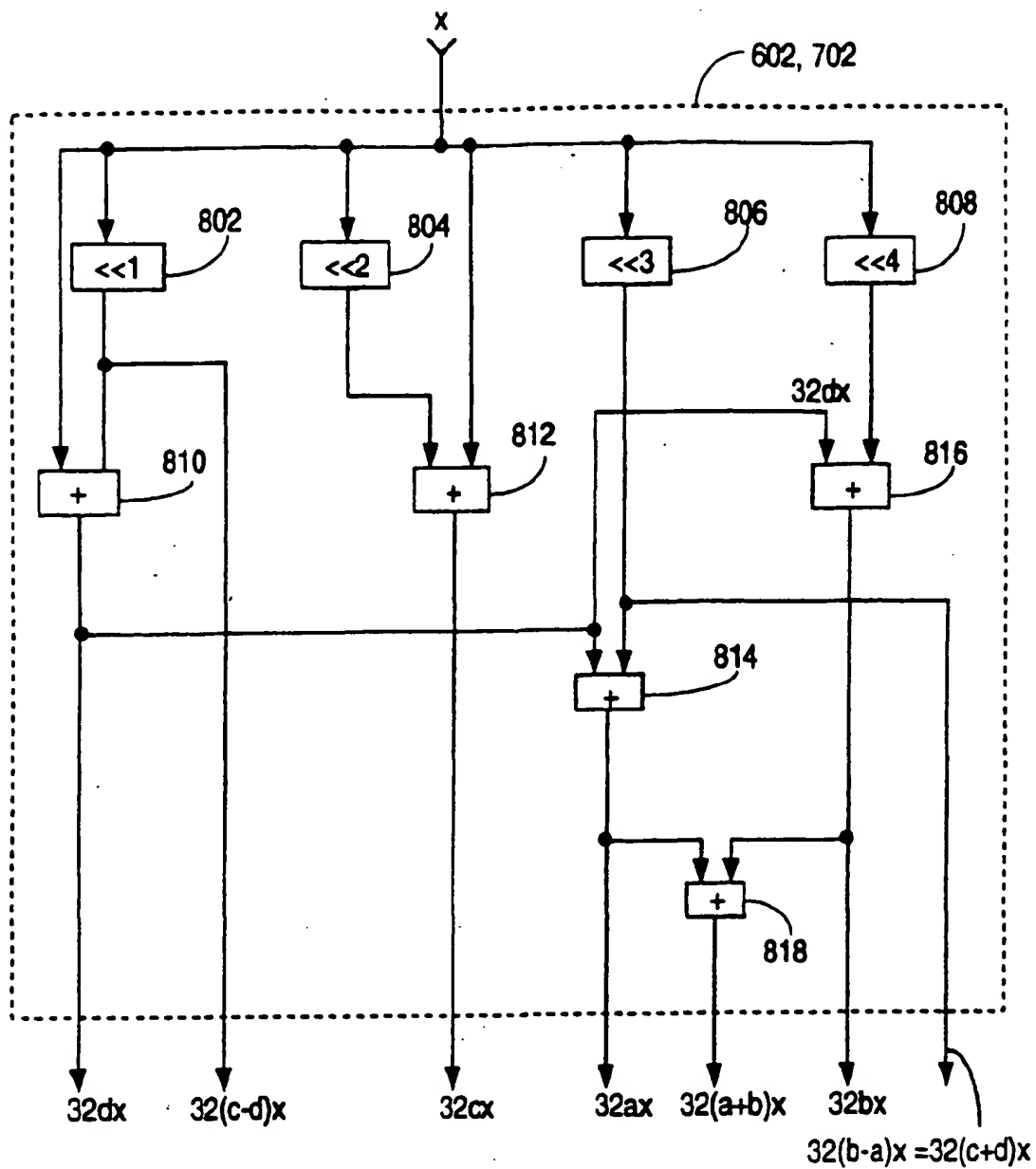


FIG. 7
CONV_COL
BLOCK DIAGRAM



$$32a = 11$$

$$32b = 19$$

$$32c = 5$$

$$32d = 3$$

$\ll n$ = SHIFT LEFT n BITS

FIG. 8
WAVELET TRANSFORM MULTIPLIER CIRCUIT

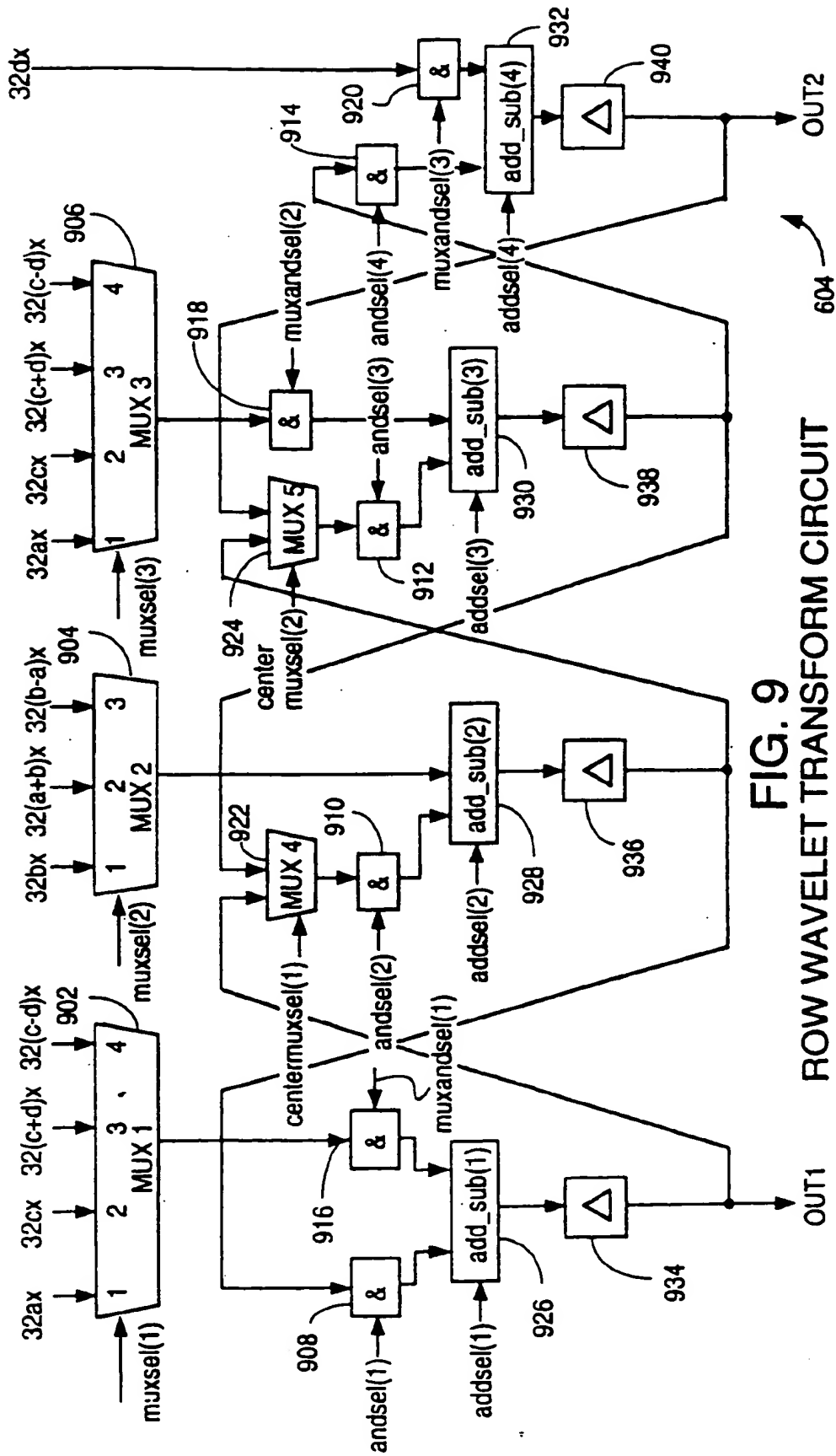


FIG. 9
ROW WAVELET TRANSFORM CIRCUIT

time t	0	1	2	3	4	5
Input Data Value x	D ₀₀	D ₀₁	D ₀₂	D ₀₃	D ₀₄	D ₀₅
muxsel(1)	1	1	1	1	1	1
muxsel(2)	2	1	1	1	1	1
muxsel(3)	3	2	2	2	2	2
andsel(1),(4)	pass	zero	pass	pass	pass	pass
andsel(2),(3)	zero	pass	pass	pass	pass	pass
addsel(1)	add	add	add	add	add	add
addsel(2)	add	sub	add	sub	add	sub
addsel(3)	add	add	add	add	add	add
addsel(4)	sub	add	sub	add	sub	add
centermuxsel (1)	l	r	l	r	l	r
centermuxsel (2)	r	l	r	l	r	l
muxandsel (1)	zero	pass	pass	pass	pass	pass
muxandsel (2)	pass	pass	pass	pass	pass	pass
muxandsel (3)	zero	pass	pass	pass	pass	pass
OUT 2				32H ₀₀		32H ₀₁
OUT 1				32G ₀₀		32G ₀₁
OUTPUT LEADS 520				H ₀₀	G ₀₀	H ₀₁

FIG. 10A

CONV_ROW CONTROL SIGNALS AND
OUTPUTS DURING FORWARD OCTAVE 0 TRANSFORM

6	7	8	9	10	11
D ₀₆	D ₀₇	D ₁₀	D ₁₁	D ₁₂	D ₁₃
1	1	1	1	1	1
1	3	2	1	1	1
2	4	3	2	2	2
pass	zero	pass	zero	pass	pass
pass	pass	zero	pass	pass	pass
add	add	add	add	add	add
add	sub	add	sub	add	sub
add	add	add	add	add	add
sub	add	sub	add	sub	add
l	r	l	r	l	r
r	l	r	l	r	l
pass	pass	zero	pass	pass	pass
pass	pass	pass	pass	pass	pass
pass	pass	zero	pass	pass	pass
	32H ₀₂		32H ₀₃		32H ₁₀
	32G ₀₂		32G ₀₃		32G ₁₀
G ₀₁	H ₀₂	G ₀₂	H ₀₃	G ₀₃	H ₁₀
					G ₁₀

KEY TO FIG. 10

FIG. 10A	FIG. 10B
----------	----------

FIG. 10

FIG. 10B

CONV_ROW CONTROL SIGNALS AND
OUTPUTS DURING FORWARD OCTAVE0 TRANSFORM

Time	Input Data Value	Output of block 926	Output of block 928
0	D ₀₀	0	32(a+b) D ₀₀
1	D ₀₁	32aD ₀₁	32((c+d)D ₀₀ - bD ₀₁)
2	D ₀₂	32G ₀₀ = 32((c+d)D ₀₀ - bD ₀₁ + aD ₀₂)	32(aD ₀₁ + bD ₀₂)
3	D ₀₃	32aD ₀₃	32(dD ₀₁ + cD ₀₂ - bD ₀₃)
4	D ₀₄	32G ₀₁ = 32(dD ₀₁ + cD ₀₂ - bD ₀₃ + aD ₀₄)	32(aD ₀₃ + bD ₀₄)
5	D ₀₅	32aD ₀₅	32(dD ₀₃ + cD ₀₄ - bD ₀₅)
6	D ₀₆	32G ₀₂ = 32(dD ₀₃ + cD ₀₄ - bD ₀₅ + aD ₀₆)	32(aD ₀₅ + bD ₀₆)
7	D ₀₇	0	32(dD ₀₅ + cD ₀₆ - (b-a)D ₀₇)
8	D ₁₀	32G ₀₃ =32(dD ₀₅ + cD ₀₆ - (b-a)D ₀₇)	32((a+b)D ₁₀)
9	D ₁₁	32aD ₁₁	32((c+d)D ₁₀ - bD ₁₁)
10	D ₁₂	32G ₁₀ =32((c+d)D ₁₀ - bD ₁₁ + aD ₁₂)	32(aD ₁₁ + bD ₁₂)
	.	.	.
	.	.	.
	.	.	.
	.	.	.
	.	.	.

FIG. 11A CONV_ROW DATA FLOW DURING THE FORWARD OCTAVE 0 TRANSFORM

Output of block 930	Output of block 932
$32(c+d)D_{00}$	0
$32((a+b)D_{00} + cD_{01})$	$32dD_{01}$
$32(dD_{01} + cD_{02})$	$32((a+b)D_{00} + cD_{01} - dD_{02}) = 32H_{00}$
$32(aD_{01} + bD_{02} + cD_{03})$	$32dD_{03}$
$32(dD_{03} + cD_{04})$	$32(aD_{01} + bD_{02} + cD_{03} - dD_{04}) = 32H_{01}$
$32(aD_{03} + bD_{04} + cD_{05})$	$32dD_{05}$
$32(dD_{05} + cD_{06})$	$32(aD_{03} + bD_{04} + cD_{05} - dD_{06}) = 32H_{02}$
$32(aD_{05} + bD_{06} + (c-d)D_{07})$	0
$32(c+d)D_{10}$	$32(aD_{05} + bD_{06} + (c-d)D_{07}) = 32H_{03}$
$32((a+b)D_{10} + cD_{11})$	$32dD_{11}$
$32(dD_{11} + cD_{12})$	$32((a+b)D_{10} + cD_{11} - dD_{12}) = 32H_{10}$
.	.
.	.
.	.
.	.
.	.

KEY TO FIG. 11

FIG. 11A	FIG. 11B
----------	----------

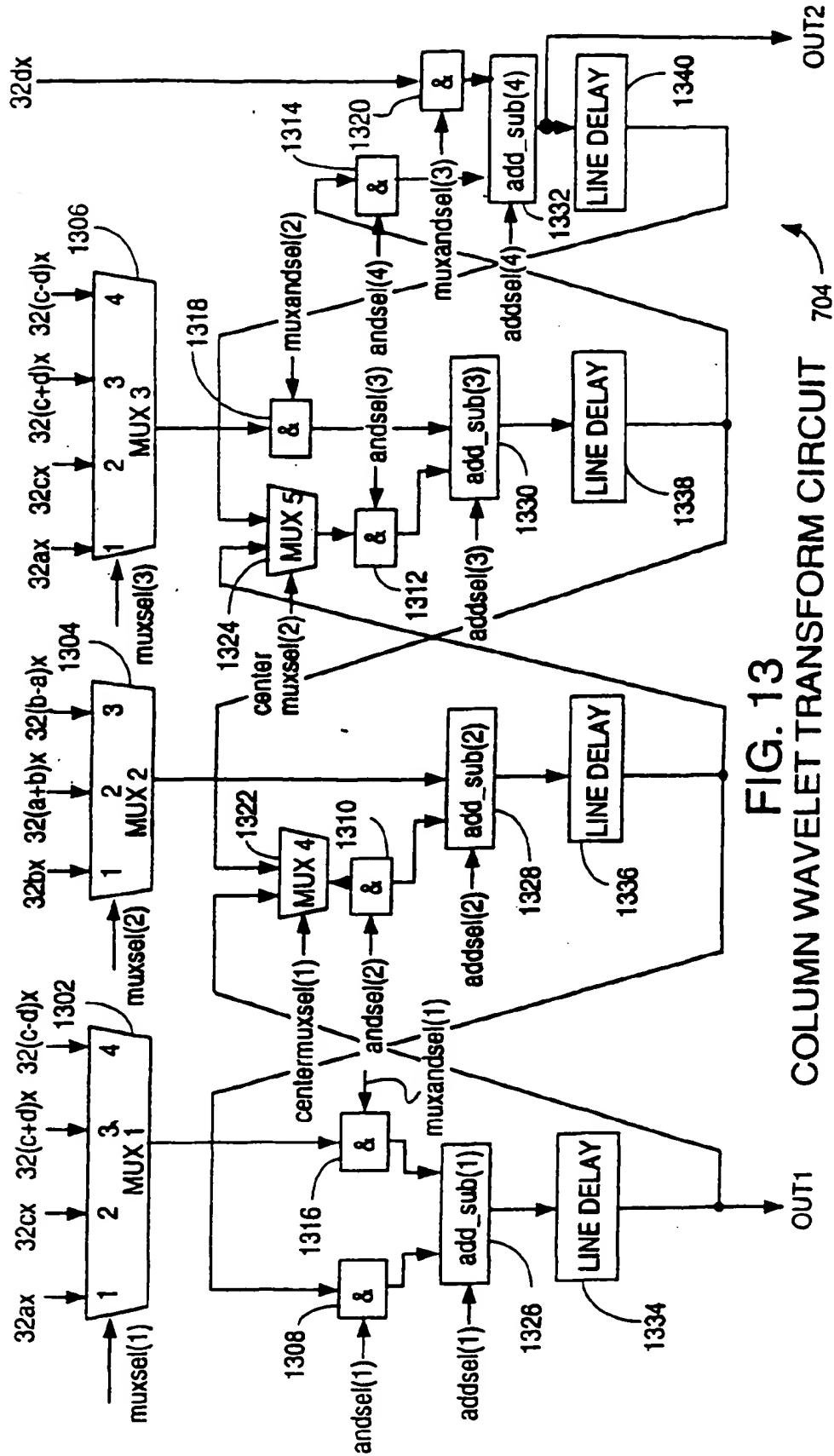
FIG. 11

FIG. 11B CONV_ROW DATA FLOW DURING THE FORWARD OCTAVE 0 TRANSFORM

H00	G00	H01	G01	H02	G02	H03	G03
H10	G10	H11	G11	H12	G12	H13	G13
H20	G20	H21	G21	H22	G22	H23	G23
H30	G30	H31	G31	H32	G32	H33	G33
H40	G40	H41	G41	H42	G42	H43	G43
H50	G50	H51	G51	H52	G52	H53	G53
H60	G60	H61	G61	H62	G62	H63	G63
H70	G70	H71	G71	H72	G72	H73	G73

FIG. 12

OUTPUT OF CONV_ROW AFTER FIRST
FORWARD TRANSFORM PASS



Time t	0 7	8 15	16 23
Input Data Value	H ₀₀ G ₀₃	H ₁₀ G ₁₃	H ₂₀ G ₂₃
muxsel(1)	1 1	1 1	1 1
muxsel(2)	2 2	1 1	1 1
Muxsel(3)	3 3	2 2	2 2
andsel(1),(4)	pass pass	zero zero	pass pass
andsel(2),(3)	zero zero	pass pass	pass pass
addsel(1)	add add	add add	add add
addsel(2)	add add	sub sub	add add
addsel(3)	add add	add add	add add
addsel(4)	sub sub	add add	sub sub
centermuxel(1)	l l	r r	l l
centermuxel(2)	r r	l l	r r
muxandsel(1)	zero zero	pass pass	pass pass
muxandsel(2)	pass pass	pass pass	pass pass
muxandsel(3)	zero zero	pass pass	pass pass
OUT2			32HH ₀₀ .. 32GH ₀₃
OUT1			
OUTPUT LEADS 524			HH ₀₀ .. GH ₀₃

FIG. 14A CONV_COL CONTROL SIGNALS
AND OUTPUTS DURING
FORWARD OCTAVE 0 TRANSFORM

24 31	32 39	56 63
H ₃₀ G ₃₃	H ₄₀ G ₄₃	H ₇₀ G ₇₃
1 1	1 1	1 1
1 1	1 1	3 3
2 2	2 2	4 4
pass pass	pass pass	zero zero
pass pass	pass pass	pass pass
add add	add add	add add
sub sub	add add	sub sub
add add	add add	add add
add add	sub sub	add add
r r	r r
....	r r
pass pass	pass pass	pass pass
pass pass	pass pass	pass pass
pass pass	pass pass	pass pass
32HG ₀₀ .. 32GG ₀₃	32HH ₁₀ .. 32GH ₁₃		32HG ₂₀ .. 32GG ₂₃
HG ₀₀ GG ₀₃	HH ₁₀ GH ₁₃		HG ₂₀ GG ₂₃

FIG. 14B CONV_COL CONTROL SIGNALS
AND OUTPUTS DURING
FORWARD OCTAVE O TRANSFORM

64	71	72	79
1	1	1	1
2	2	1	1
3	3	2	2
pass	pass	zero	zero
zero	zero	pass	pass
add	add	add	add
add	add	sub	sub
add	add	add	add
sub	sub	add	add
		r	r
r	r		
zero	zero	pass	pass
pass	pass	pass	pass
zero	zero	pass	pass
32HH ₃₀ .. 32GH ₃₃			32HG ₃₀ .. 32GG ₃₃		
HH ₃₀ GH ₃₃			HG ₃₀ GG ₃₃		

KEY TO FIG 14

FIG. 14A	FIG. 14B	FIG. 14C
-------------	-------------	-------------

FIG 14

FIG. 14C CONV_COL CONTROL SIGNALS
AND OUTPUTS DURING
FORWARD OCTAVE 0 TRANSFORM

Time	Input Data Value	Output of block 1326	Output of block 1328
0	H ₀₀	.	$32(a+b)H_{00}$
1	G ₀₀	.	$32(a+b)G_{00}$
.	.	.	.
7	G ₀₃	.	$32(a+b)G_{03}$
8	H ₁₀	$32aH_{10}$	$32((c+d)H_{00} - bH_{10})$
.	.	.	.
15	G ₁₃	$32aG_{13}$	$32((c+d)G_{03} - bG_{13})$
16	H ₂₀	$32HG_{00} = 32((c+d)H_{00} - bH_{10} + aH_{20})$	$32(aH_{10} + bH_{20})$
.	.	.	.
23	G ₂₃	$32GG_{03} = 32((c+d)G_{03} - bG_{13} + aG_{23})$	$32(aG_{13} + bG_{23})$
24	H ₃₀	$32aH_{30}$	$32(dH_{10} + cH_{20} - bH_{30})$
.	.	.	.
31	G ₃₃	$32aG_{33}$	$32(dG_{13} + cG_{23} - bG_{33})$
32	H ₄₀	$32HG_{10} = 32(dH_{10} + cH_{20} - bH_{30} + aH_{40})$	$32(aH_{30} + bH_{40})$
.	.	.	.

FIG. 15A
CONV_COL DATA FLOW DURING FORWARD OCTAVE 0 TRANSFORM

Output of block 1330	Output of block 1332
$32(c+d)H_{00}$.
$32(c+d)G_{00}$.
.	.
$32(c+d)G_{03}$.
$32((a+b)H_{00} + cH_{10})$	$32dH_{10}$
.	.
$32((a+b)G_{03} + cG_{13})$	$32dG_{13}$
$32(dH_{10} + cH_{20})$	$32((a+b)H_{00} + cH_{10} \cdot dH_{20}) = 32HH_{00}$
.	.
$32(dG_{13} + cG_{23})$	$32((a+b)G_{03} + cG_{13} \cdot dG_{23}) = 32GH_{03}$
$32(aH_{10} + bH_{20} + cH_{30})$	$32dH_{30}$
.	.
$a32(G_{13} + bG_{23} + cG_{33})$	$32dG_{33}$
$32(dH_{30} + cH_{40})$	$32(aH_{10} + bH_{20} + cH_{30} \cdot dH_{40}) = 32HH_{10}$
.	.

FIG. 15B

CONV_COL DATA FLOW DURING FORWARD OCTAVE 0 TRANSFORM

39	G ₄₃	$32GG_{13} = 32(dG_{13} + cG_{23} - bG_{33} + aG_{43})$	$32(aG_{33} + bG_{43})$		

56	H ₇₀		$32(dH_{50} + cH_{60} - (b-a)H_{70})$		

63	G ₇₃		$32(dG_{53} + cG_{63} - (b-a)G_{73})$		
64		$32HG_{30} = 32(dH_{50} + cH_{60} - (b-a)H_{70})$			

71		$32GG_{33} = 32(dG_{53} + cG_{63} - (b-a)G_{73})$			

FIG. 15C
CONV_COL DATA FLOW DURING FORWARD OCTAVE 0 TRANSFORM

$32\{dG_{33} + cG_{43}\}$	$32\{aG_{13} + bG_{23} + cG_{33} + dG_{43}\} = 32GH_{13}$
$32\{aH_{50} + bH_{60} + (c-d)H_{70}\}$	
$32\{aG_{53} + bG_{63} + (c-d)G_{73}\}$	
	$32\{aH_{50} + bH_{60} + (c-d)H_{70}\} = 32HH_{30}$
	$32\{aG_{53} + bG_{63} + (c-d)G_{73}\} = 32GH_{33}$

KEY TO FIG 15			
FIG. 15A		FIG. 15B	
FIG. 15C		FIG. 15D	

FIG 15

FIG. 15D
CONV_COL DATA FLOW DURING FORWARD OCTAVE 0 TRANSFORM

HH ₀₀	GH ₀₀	HH ₀₁	GH ₀₁	HH ₀₂	GH ₀₂	HH ₀₃	GH ₀₃
HG ₀₀	GG ₀₀	HG ₀₁	GG ₀₁	HG ₀₂	GG ₀₂	HG ₀₃	GG ₀₃
HH ₁₀	GH ₁₀	HH ₁₁	GH ₁₁	HH ₁₂	GH ₁₂	HH ₁₃	GH ₁₃
HG ₁₀	GG ₁₀	HG ₁₁	GG ₁₁	HG ₁₂	GG ₁₂	HG ₁₃	GG ₁₃
HH ₂₀	GH ₂₀	HH ₂₁	GH ₂₁	HH ₂₂	GH ₂₂	HH ₂₃	GH ₂₃
HG ₂₀	GG ₂₀	HG ₂₁	GG ₂₁	HG ₂₂	GG ₂₂	HG ₂₃	GG ₂₃
HH ₃₀	GH ₃₀	HH ₃₁	GH ₃₁	HH ₃₂	GH ₃₂	HH ₃₃	GH ₃₃
HG ₃₀	GG ₃₀	HG ₃₁	GG ₃₁	HG ₂₂	GG ₃₂	HG ₃₃	GG ₃₃

FIG. 16
OCTAVE 0 DECOMPOSITION
AFTER FIRST CONV_ROW AND
CONV_COL PASS

Time t	0	1	2	3	4	5
Input Data Value	HH ₀₀	HH ₀₁	HH ₀₂	HH ₀₃	HH ₁₀	HH ₁₁
muxsel(1)	1	1	1	1	1	1
muxsel(2)	2	1	1	3	2	1
muxsel(3)	3	2	2	4	3	2
andsel(1),(4)	pass	zero	pass	zero	pass	zero
andsel(2),(3)	zero	pass	pass	pass	zero	pass
addsel(1)	add	add	add	add	add	add
addsel(2)	add	sub	add	sub	add	sub
addsel(3)	add	add	add	add	add	add
addsel(4)	sub	add	sub	add	sub	add
center muxsel (1)	l	r	l	r	l	r
center muxsel (2)	r	l	r	l	r	l
muxandsel (1)	zero	pass	pass	pass	zero	pass
muxandsel (2)	pass	pass	pass	pass	pass	pass
muxandsel (3)	zero	pass	pass	pass	zero	pass
OUT 2				32HHH ₀₀		32HHH ₀₁
OUT 1				32HHG ₀₀		32HHG ₀₁
OUTPUT LEADS				HHH ₀₀	HHG ₀₀	HHH ₀₁
520						

FIG. 17A

CONV_ROW CONTROL SIGNALS AND OUTPUTS
DURING THE FORWARD OCTAVE 1 TRANSFORM

6	7	8	9	10	11
HH ₁₂	HH ₁₃	HH ₂₀	HH ₂₁	HH ₂₂	HH ₂₃
1	1	1	1	1	1
1	3	2	1	1	1
2	4	3	2	2	2
pass	zero	pass	zero	pass	zero
pass	pass	zero	pass	pass	pass
add	add	add	add	add	add
add	sub	add	sub	add	sub
add	add	add	add	add	add
sub	add	sub	add	sub	add
l	r	l	r	l	r
r	l	r	l	r	l
pass	pass	zero	pass	pass	pass
pass	pass	pass	pass	pass	pass
pass	pass	zero	pass	pass	pass
	32HHH ₁₀		32HHH ₁₁		32HHH ₂₀
	32HHG ₁₀		32HHG ₁₁		32HHG ₂₀
HHG ₀₁	HHH ₁₀	HHG ₁₀	HHH ₁₁	HHG ₁₁	HHH ₂₀
					HHG ₂₀

KEY TO FIG. 17

FIG. 17A	FIG. 17B
-------------	-------------

FIG. 17

FIG. 17B

CONV_ROW CONTROL SIGNALS AND OUTPUTS
DURING THE FORWARD OCTAVE 1 TRANSFORM

Time	Input Data Value	Output of Block 926	Output of Block 928
0	HH ₀₀	0	$32((a+b)HH_{00})$
1	HH ₀₁	$32aHH_{01}$	$32((c+d)HH_{00} - bHH_{01})$
2	HH ₀₂	$32HHG_{00} = 32((c+d)HH_{00} - bHH_{01} + aHH_{02})$	$32(aHH_{01} + bHH_{02})$
3	HH ₀₃	0	$32(dHH_{01} + cHH_{02} - (b-a)HH_{03})$
4	HH ₁₀	$32HHG_{01} = 32(dHH_{01} + cHH_{02} - (b-a)HH_{03})$	$32((a+b)HH_{10})$
5	HH ₁₁	$32aHH_{11}$	$32((c+d)HH_{10} - bHH_{11})$
6	HH ₁₂	$32HHG_{10} = 32((c+d)HH_{10} - bHH_{11} + aHH_{12})$	$32(aHH_{11} + bHH_{12})$
	.	.	.
	.	.	.
	.	.	.
	.	.	.
	.	.	.

FIG. 18A

CONV_ROW DATA FLOW DURING THE
FORWARD OCTAVE 1 TRANSFORM

Output of Block 930	Output of Block 932
32(c+d) HH ₀₀	0
32((a+b)HH ₀₀ + cHH ₀₁)	32dHH ₀₁
32(dHH ₀₁ + cHH ₀₂)	32((a+b)HH ₀₀ + cHH ₀₁ - dHH ₀₂)=32HH ₀₀
32(aHH ₀₁ + bHH ₀₂ + (c-d)HH ₀₃)	0
32(c+d)HH ₁₀	32(aHH ₀₁ + bHH ₀₂ + (c-d)HH ₀₃) =32HH ₀₁
32((a+b)HH ₁₀ +cHH ₁₁)	32dHH ₁₁
32(dHH ₁₁ + cHH ₁₂)	32((a+b)HH ₁₀ + cHH ₁₁ - dHH ₁₂)=32HH ₁₀
...	...
...	...
...	...
...	...
...	...

KEY TO FIG. 18

FIG. 18A	FIG. 18B
----------	----------

FIG. 18

FIG. 18B

CONV_ROW DATA FLOW DURING THE
FORWARD OCTAVE 1 TRANSFORM

Time t	0 3	4 7	8 11
Input Data Values	HHH ₀₀ ... HHG ₀₁	HHH ₁₀ ... HHG ₁₁	HHH ₂₀ ... HHG ₂₁
muxsel(1)	1 1	1 1	1 1
muxsel(2)	2 2	1 1	1 1
muxsel(3)	3 3	2 2	2 2
andsel(1),(4)	pass pass	zero zero	pass pass
andsel(2),(3)	zero zero	pass pass	pass pass
addsel(1)	add add	add add	add add
addsel(2)	add add	sub sub	add add
addsel(3)	add add	add add	add add
addsel(4)	sub sub	add add	sub sub
centermuxel(1)	l l	r r	l l
centermuxel(2)	r r	l l	r r
muxandsel(1)	zero zero	pass pass	pass pass
muxandsel(2)	pass pass	pass pass	pass pass
muxandsel(3)	zero zero	pass pass	pass pass
OUT2			32HHH ₀₀ ... 32HHG ₀₁
OUT1			
OUTPUT LEADS 524			HHH ₀₀ ... HHG ₀₁

FIG. 19A

CONV_COL CONTROL SIGNALS AND OUTPUTS
DURING THE FORWARD OCTAVE 1 TRANSFORM

12 15	16 19	20 23
HHH ₃₀ .. HHG ₃₁		
1 1	1 1	1 1
3 3	2 2	1 1
4 4	3 3	2 2
zero zero	pass pass	zero zero
pass pass	zero zero	pass pass
add add	add add	add add
sub sub	add add	sub sub
add add	add add	add add
add add	sub sub	add add
r r	r r
....	r r
pass pass	zero zero	pass pass
pass pass	pass pass	pass pass
pass pass	zero zero	pass pass
32HHHG ₀₀ .. 32HHGG ₀₁	32HHHH ₁₀ .. 32HHGH ₁₁	32HHHG ₁₀ .. 32HHGG ₁₁
HHHG ₀₀ .. HHGG ₀₁	HHHH ₁₀ .. HHGH ₁₁	HHHG ₁₀ .. HHGG ₁₁

FIG. 19B

KEY TO FIG 19

FIG. 19A	FIG. 19B
-------------	-------------

FIG 19

CONV_COL CONTROL SIGNALS AND OUTPUTS
DURING THE FORWARD OCTAVE 1 TRANSFORM

t	Input Data Value	Output of Block 1326	
0	HHH ₀₀	.	
.		.	
.		.	
3	HHG ₀₁	.	
4	HHH ₁₀	32aHHH ₁₀	
.		.	
.		.	
7	HHG ₁₁	32aHHG ₁₁	
8	HHH ₂₀	$32HHHG_{00} = 32\{(c+d)HHH_{00} - bHHH_{10} + aHHH_{20}\}$	
.		.	
.		.	
11	HHG ₂₁	$32HHGG_{01} = 32\{(c+d)HHG_{01} - bHHG_{11} + aHHG_{21}\}$	
12	HHH ₃₀	.	
.		.	
.		.	
15	HHG ₃₁	.	
16		$32HHHG_{10} = 32\{dHHH_{10} + cHHH_{20} - (b-a)HHH_{30}\}$	
.		.	
.		.	
19		$32HHGG_{11} = 32\{dHHG_{11} + HHG_{21} - (b-a)HHG_{31}\}$	

FIG. 20A

CONV_COL DATA FLOW FOR FORWARD
OCTAVE 1 TRANSFORM

Output of Block 1328	Output of Block 1330
$32(a+b)HHH_{00}$	$32(c+d)HHH_{00}$
.	.
$32(a+b)HHG_{01}$	$32(c+d)HHG_{01}$
$32((c+d)HHH_{00} - bHHH_{10})$	$32((a+b)HHH_{00} + cHHH_{10})$
.	.
$32((c+d)HHG_{01} - bHHG_{11})$	$32((a+b)HHG_{01} + cHHG_{11})$
$32(aHHH_{10} + bHHH_{20})$	$32(dHHH_{10} + cHHH_{20})$
.	.
$32(aHHG_{11} + bHHG_{21})$	$32(dHHG_{11} + cHHG_{21})$
$32(dHHH_{10} + cHHH_{20} - (b-a)HHH_{30})$	$32(aHHH_{10} + bHHH_{20} + (c-d)HHH_{30})$
.	.
$32(dHHG_{11} + cHHG_{21} - (b-a)HHG_{31})$	$32(aHHG_{11} + bHHG_{20} + (c-d)HHG_{31})$

FIG. 20B

CONV_COL DATA FLOW FOR FORWARD
OCTAVE 1 TRANSFORM

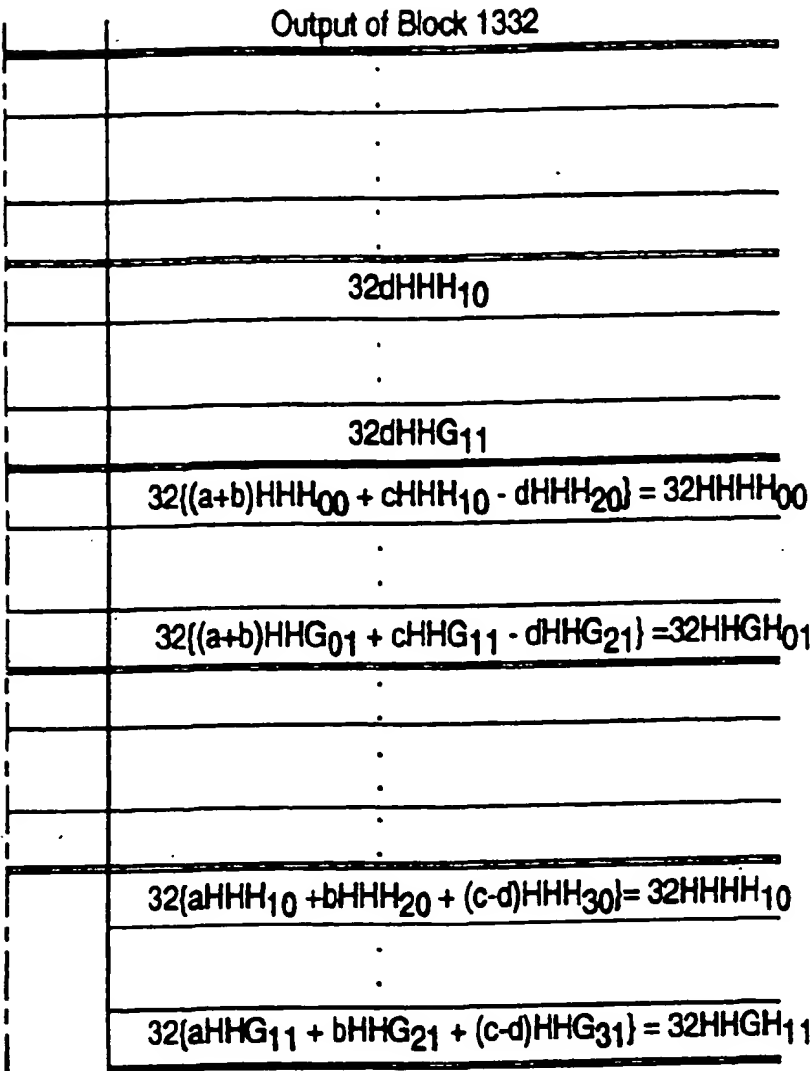


FIG. 20C
CONV_COL DATA FLOW FOR FORWARD
OCTAVE 1 TRANSFORM

KEY TO FIG. 20

FIG. 20A	FIG. 20B	FIG. 20C
-------------	-------------	-------------

FIG. 20

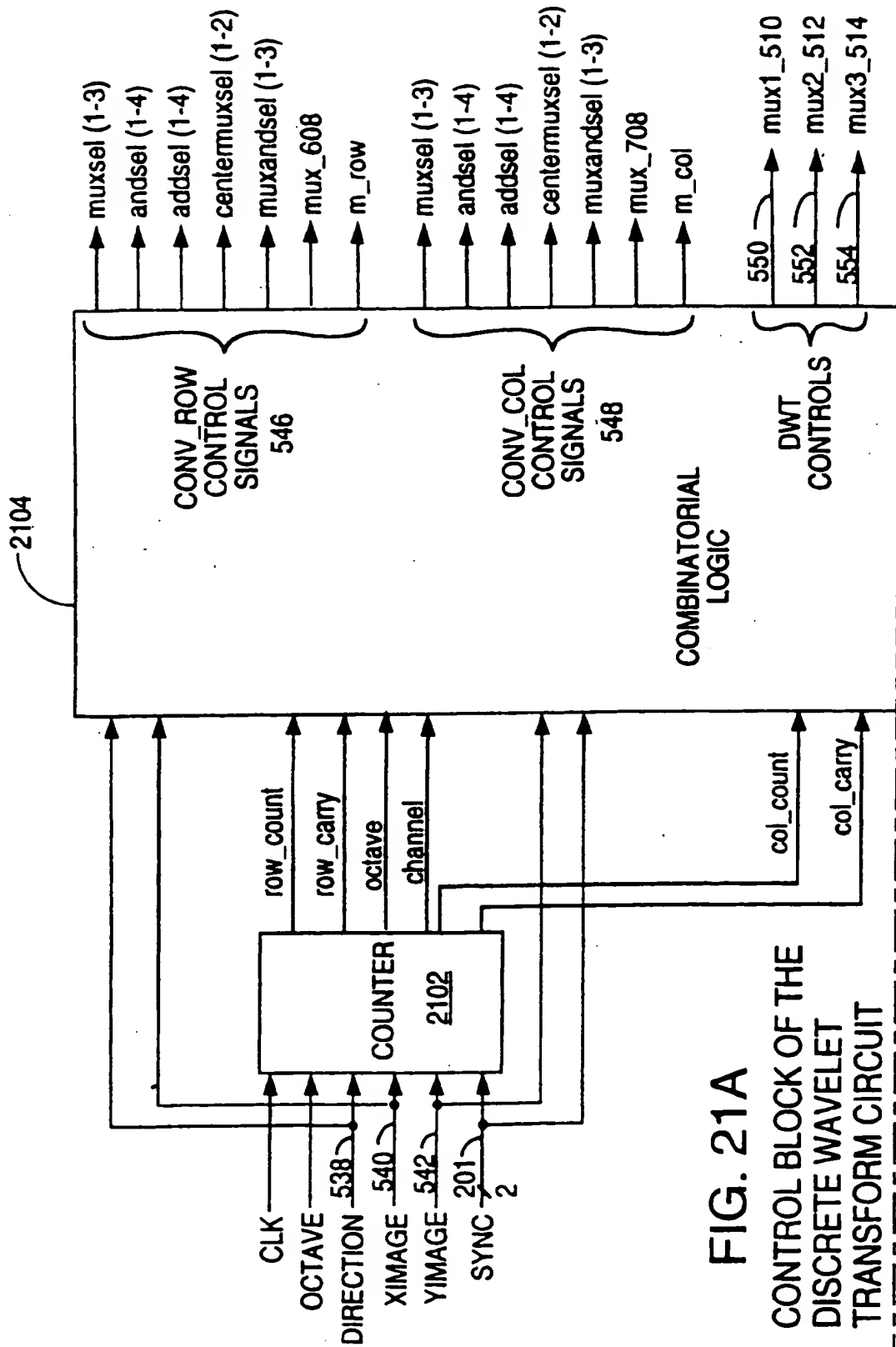
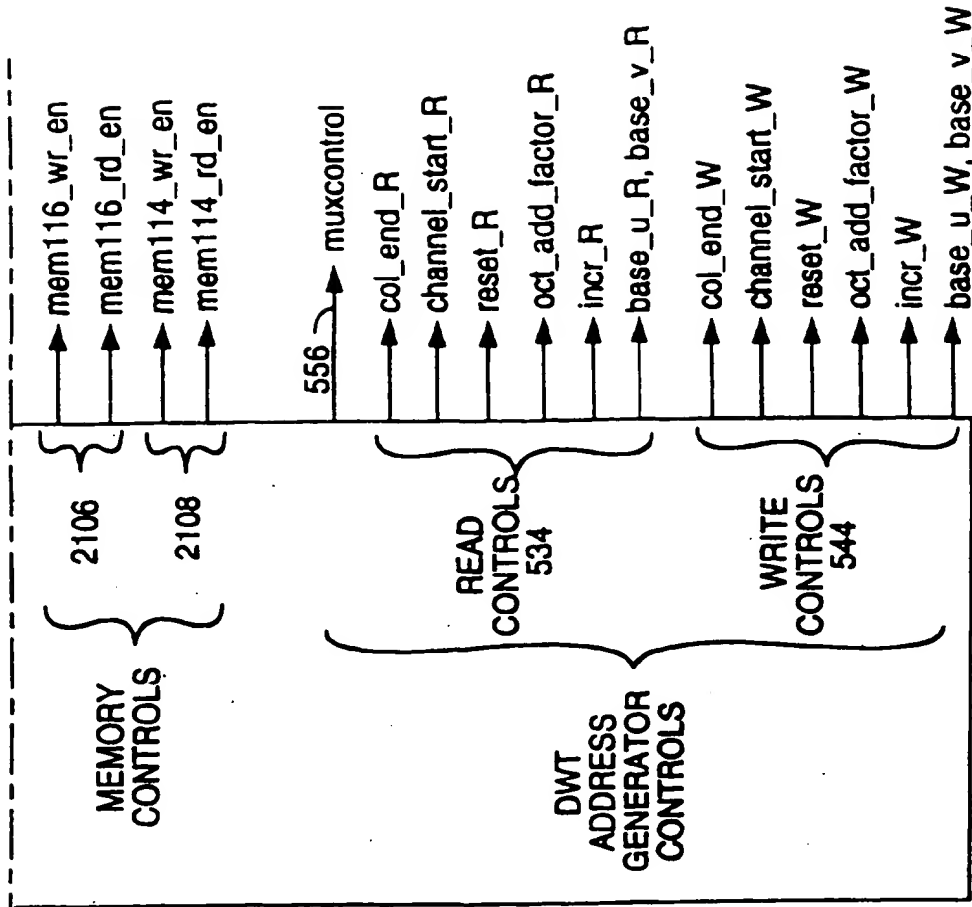


FIG. 21B
CONTROL BLOCK OF THE
DISCRETE WAVELET
TRANSFORM CIRCUIT



KEY TO FIG. 21

FIG. 21A
FIG. 21B

FIG. 21

Time t	0 3	4 7	8 11
muxsel(1)	2 2	4 4	3 3
muxsel(2)	3 3	1 1	1 1
muxsel(3)	1 1	1 1	1 1
andsel(1),(4)	zero zero	pass pass	zero zero
andsel(2)	zero zero	pass pass	pass pass
andsel(3)	pass pass	pass pass	pass pass
addsel(1)	add add	add add	add add
addsel(2)	add add	sub sub	add add
addsel(3)	add add	add add	add add
addsel(4)	sub sub	add add	sub sub
centermuxel(1)	r r	l l	r r
centermuxel(2)	l l	r r	l l
muxandsel(1)	pass pass	pass pass	pass pass
muxandsel(2)	zero zero	pass pass	pass pass
muxandsel(3)	pass pass	zero zero	pass pass
OUT2			
OUT1			8HHH ₀₀ .. 8HHG ₀₁
OUTPUT LEADS 524			HHH ₀₀ .. HHG ₀₁

FIG. 22A

CONV_COL CONTROL SIGNALS AND OUTPUTS
FOR INVERSE OCTAVE 1 TRANSFORM

12 15	16 19	20 23
2 2	2 2	4 4
2 2	3 3	1 1
1 1	1 1	1 1
pass pass	zero zero	pass pass
pass pass	zero zero	pass pass
pass pass	pass pass	pass pass
add add	add add	add add
sub sub	add add	sub sub
add add	add add	add add
add add	sub sub	add add
....	r r
r r	r r
pass pass	pass pass	pass pass
pass pass	zero zero	pass pass
pass pass	pass pass	zero zero
16HHH ₁₀ ... 16HHG ₁₁	16HHH ₂₀ ... 16HHG ₂₁	8HHH ₃₀ ... 8HHG ₃₁
HHH ₁₀ ... HHG ₁₁	HHH ₂₀ ... HHG ₂₁	HHH ₃₀ ... HHG ₃₁

FIG. 22B
 CONV_COL CONTROL SIGNALS AND OUTPUTS
 FOR INVERSE OCTAVE 1 TRANSFORM

KEY TO FIG 22

FIG. 22A	FIG. 22B
-------------	-------------

FIG 22

Time t	Input Data Value	OUTPUT OF BLOCK 1326
0	HHHH ₀₀	32cHHHH ₀₀
.	.	.
3	HHGH ₀₁	32cHHGH ₀₁
4	HHHG ₀₀	$8HHH_{00} = 32\{(b-a)HHHH_{00} + (c-d)HHHG_{00}\}$
.	.	.
7	HHGG ₀₁	$8HHG_{01} = 32\{(b-a)HHGH_{01} + (c-d)HHGG_{01}\}$
8	HHHH ₁₀	32(c+d)HHHH ₁₀
.	.	.
11	HHGH ₁₁	32(c+d)HHGH ₁₁
12	HHHG ₁₀	$16HHH_{20} = 32\{-dHHHH_{00} + aHHHG_{00} + bHHHH_{10} + cHHGG_{10}\}$
.	.	.
15	HHGG ₁₁	$16HHG_{21} = 32\{-dHHGH_{01} + aHHGG_{01} + bHHGH_{11} + cHHHG_{11}\}$
16		
.		
19		
20		
.		
23		

FIG. 23A

CONV_COL DATA FLOW FOR INVERSE
OCTAVE 1 TRANSFORM

OUTPUT OF BLOCK 1328	OUTPUT OF BLOCK 1330
$32(b-a)HHHH_{00}$	\vdots
\vdots	\vdots
$32(b-a)HHGH_{01}$	\vdots
$32(cHHHH_{00} - bHHHG_{00})$	$32(-dHHHH_{00} + aHHHG_{00})$
\vdots	\vdots
$32(cHHGH_{01} - bHHGG_{01})$	$32(-dHHGH_{01} + aHHGG_{01})$
$32(-dHHHH_{00} + aHHHG_{00} + bHHHH_{10})$	$32(cHHHH_{00} - bHHHG_{00} + aHHHH_{10})$
\vdots	\vdots
$32(-dHHGH_{01} + aHHGG_{01} + bHHGH_{11})$	$32(cHHGH_{01} - bHHGG_{01} + aHHGH_{11})$
$32((c+d)HHHH_{10} - (a+b)HHHG_{10})$	\vdots
\vdots	\vdots
$32((c+d)HHGH_{11} - (a+b)HHGG_{11})$	\vdots
	$32((c+d)HHHH_{10} - (a+b)HHHG_{10})$
	\vdots
	$32((c+d)HHGH_{11} - (a+b)HHGG_{11})$

FIG. 23B

CONV_COL DATA FLOW FOR INVERSE
OCTAVE 1 TRANSFORM

OUTPUT OF BLOCK 1332	
	$32\{-dHHHH_{00}\}$
	.
	$32\{-dHHGH_{01}\}$
	.
	.
	.
	.
	.
	.
	$32\{cHHHH_{00} - bHHHG_{00} + aHHHH_{10} + dHHHG_{10}\} = 16HHH_{10}$
	.
	$32\{cHHGH_{01} - bHHGG_{01} + aHHGH_{01} + dHHGG_{11}\} = 16HHG_{11}$
	$32\{(c+d)HHHH_{10} - (a+b)HHHG_{10}\} = 8 HHH_{30}$
	.
	.
	$32\{(c+d)HHGH_{11} - (a+b)HHGG_{11}\} = 8 HHG_{31}$

KEY TO FIG. 23

FIG. 23A	FIG. 23B	FIG. 23C
----------	----------	----------

FIG. 23C

CONV_COL DATA FLOW FOR INVERSE
OCTAVE 1 TRANSFORM

Time t	0	1	2	3	4	
INPUT DATA VALUE	HHH ₀₀	HHG ₀₀	HHH ₀₁	HHG ₀₁	HHH ₁₀	
muxsel(1)	2	4	3	2	2	
muxsel(2)	3	1	1	2	3	
muxsel(3)	1	1	1	1	1	
andsel(1),(4)	zero	pass	zero	pass	zero	
andsel(2)	zero	pass	pass	pass	zero	
andsel(3)	pass	pass	pass	pass	pass	
addsel(1)	add	add	add	add	add	
addsel(2)	add	sub	add	sub	add	
addsel(3)	add	add	add	add	add	
addsel(4)	sub	add	sub	add	sub	
centermuxsel (1)	r	l	r	l	r	
centermuxsel (2)	l	r	l	r	l	
muxandsel (1)	pass	pass	pass	pass	pass	
muxandsel (2)	zero	pass	pass	pass	zero	
muxandsel (3)	pass	zero	pass	pass	pass	
OUT 2					16HH ₀₁	
OUT 1			8HH ₀₀		16HH ₀₂	
OUTPUT LEADS 520				HH ₀₀	HH ₀₁	HH ₀₂

FIG. 24
CONV_ROW CONTROL SIGNALS AND OUTPUTS
FOR INVERSE OCTAVE 1 TRANSFORM

Time	Input Data Value	Output of Block 926	Output of Block 928
0	HHH ₀₀	32cHHH ₀₀	32(b-a)HHH ₀₀
1	HGG ₀₀	8HH ₀₀ = 32((b-a)HHH ₀₀ + (c-d)HHG ₀₀)	32(cHHH ₀₀ - bHHG ₀₀)
2	HHH ₀₁	32(c+d)HHH ₀₁	32(-dHHH ₀₀ + aHHG ₀₀ + bHHH ₀₁)
3	HGG ₀₁	16HH ₀₂ = 32(-dHHH ₀₀ + aHHG ₀₀ + bHHH ₀₁ + cHHG ₀₁)	32((c+d)HHH ₀₃ - (a+b)HHG ₀₁)
4	HHH ₁₀	32cHHH ₁₀	32(b-a)HHH ₁₀
5	HGG ₁₀	8HH ₁₀ = 32((b-a)HHH ₁₀ + (c-d)HHG ₁₀)	32(cHHH ₁₀ - bHHG ₁₀)
6			

FIG. 25A
CONV_ROW DATA FLOW FOR INVERSE OCTAVE 1 TRANSFORM

Output of Block 930	Output of Block 932
	$-32dHHH_{00}$
$32\{-dHHH_{00} + aHHG_{00}\}$	
$32\{cHHH_{00} - bHHG_{00} + aHHH_{01}\}$	
	$16HH_{01} = 32\{cHHH_{00} - bHHG_{00} + aHHH_{01} + dHHG_{01}\}$
$32\{(c+d)HHH_{01} - (a+b)HHG_{01}\}$	$32\{-dHHH_{10}\}$
$32\{-dHHH_{10} + aHHG_{10}\}$	$8HH_{03} = 32\{(c+d)HHH_{01} - (a+b)HHG_{01}\}$

FIG. 25B
CONV_ROW DATA FLOW FOR INVERSE OCTAVE 1 TRANSFORM

KEY TO FIG. 25

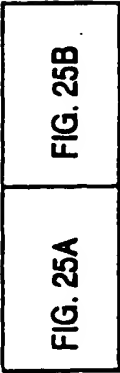


FIG. 25

Time t	0 7	8 15	16 23
muxsel(1)	2 2	4 4	2 2
muxsel(2)	3 3	1 1	1 1
muxsel(3)	1 1	1 1	1 1
andsel(1),(4)	zero zero	pass pass	zero zero
andsel(2)	zero zero	pass pass	pass pass
andsel(3)	pass pass	pass pass	pass pass
addsel(1)	add add	add add	add add
addsel(2)	add add	sub sub	add add
addsel(3)	add add	add add	add add
addsel(4)	sub sub	add add	sub sub
centermuxel(1)	r r	r r
centermuxel(2)	r r
muxandsel(1)	pass pass	pass pass	pass pass
muxandsel(2)	zero zero	pass pass	pass pass
muxandsel(3)	pass pass	zero zero	pass pass
out2			
out1			8H00 8G03
OUTPUT LEADS 524			H00 G03

FIG. 26A
 CONV_COL CONTROL SIGNALS AND OUTPUTS
 DURING INVERSE OCTAVE 0 TRANSFORM

24 31	32 39	48 55
2 2	2 2	3 3
1 1	1 1	1 1
1 1	1 1	1 1
pass pass	zero zero	zero zero
pass pass	pass pass	pass pass
pass pass	pass pass	pass pass
add add	add add	add add
sub sub	add add	add add
add add	add add	add add
add add	sub sub	sub sub
l l	r l	r r
r r	r l	l l
pass pass	pass pass	pass pass
pass pass	pass pass	pass pass
pass pass	pass pass	pass pass
16H ₁₀ ... 16G ₁₃	16H ₂₀ 16G ₂₃		16H ₄₀ 16G ₄₃
H ₁₀ G ₁₃	H ₂₀ G ₂₃		H ₄₀ G ₄₃

FIG. 26B

CONV_COL CONTROL SIGNALS AND OUTPUTS
DURING INVERSE OCTAVE 0 TRANSFORM

56 63	64 71	72 79
2 2	2 2	4 4
2 2	3 3	1 1
1 1	1 1	1 1
pass pass	zero zero	pass pass
pass pass	zero zero	pass pass
pass pass	pass pass	pass pass
add add	add add	add add
sub sub	add add	sub sub
add add	add add	add add
add add	sub sub	add add
....	r r
r r	r r
pass pass	pass pass	pass pass
pass pass	zero zero	pass pass
pass pass	pass pass	zero zero
16H ₅₀ 16G ₅₃	16H ₆₀ 16G ₆₃	8H ₇₀ 8G ₇₃
H ₅₀ G ₅₃	H ₆₀ G ₆₃	H ₇₀ G ₇₃

KEY TO FIG 26

FIG. 26A	FIG. 26B	FIG. 26C
-------------	-------------	-------------

FIG 26

FIG. 26C

CONV_COL CONTROL SIGNALS AND OUTPUTS
DURING INVERSE OCTAVE 0 TRANSFORM

Time t	Input Data Value	Output of Block 1326	
0	HH ₀₀	32cHH ₀₀	
.	.	.	
7	GH ₀₃	32cGH ₀₃	
8	HG ₀₀	$8H_{00} = 32\{(b-a)HH_{00} + (c-d)HG_{00}\}$	
.	.	.	
15	GG ₀₃	$8G_{03} = 32\{(b-a)GH_{03} + (c-d)GG_{03}\}$	
16	HH ₁₀	32cHH ₁₀	
.	.	.	
23	GH ₁₃	32cGH ₁₃	
24	HG ₁₀	$16H_{20} = 32\{-dHH_{00} + aHG_{00} + bHH_{10} + cHG_{10}\}$	
.	.	.	
31	GG ₁₃	$16G_{23} = 32\{-dGH_{03} + aGG_{03} + bGH_{13} + cGG_{13}\}$	
32	HH ₂₀	32cHH ₂₀	
.	.	.	
39	GH ₂₃	32cGH ₂₃	

FIG. 27A

CONV_COL DATA FLOW FOR THE
INVERSE OCTAVE 0 TRANSFORM

Output of Block 1328	Output of Block 1330
$32(b-a)HH_{00}$.
.	.
$32(b-a)GH_{03}$.
$32cHH_{00} - 32bHG_{00}$	$-32dHH_{00} + 32aHG_{00}$
.	.
$32cGH_{03} - 32bGG_{03}$	$-32dGH_{03} + 32aGG_{03}$
$32(-dHH_{00} + aHG_{00} + bHH_{10})$	$32(cHH_{00} - bHG_{00} + aHH_{10})$
.	.
$32(-dGH_{03} + aGG_{03} + bGH_{13})$	$32(cGH_{03} - bGG_{03} + aGH_{13})$
$32cHH_{10} - 32bHG_{10}$	$-32dHH_{10} + 32aHG_{10}$
.	.
$32cGH_{13} - 32bGG_{13}$	$-32dGH_{13} + 32aGG_{13}$
$32(-dHH_{10} + aHG_{10} + bHH_{20})$	$32(cHH_{10} - bHG_{10} + aHH_{20})$
.	.
$32(-dGH_{13} + aGG_{13} + bGH_{23})$	$32(cGH_{13} - bGG_{13} + aGH_{23})$
.	.

FIG. 27B

CONV_COL DATA FLOW FOR THE
INVERSE OCTAVE 0 TRANSFORM

Output of Block 1330	
	$-32dHH_{00}$
	.
	.
	$-32dGH_{03}$
	.
	.
	.
	$-32dHH_{10}$
	.
	.
	$-32dGH_{13}$
	$32(cHH_{00} - bHG_{00} + aHH_{10} + dHG_{10}) = 16H_{10}$
	.
	$32(cGH_{03} - bGG_{03} + aGH_{13} + dGG_{13}) = 16G_{13}$
	$-32dHH_{20}$
	.
	.
	$-32dGH_{23}$

FIG. 27C
 CONV_COL DATA FLOW FOR THE
 INVERSE OCTAVE 0 TRANSFORM

48	HH ₃₀	$32(c+d)HH_{30}$	
.		.	
55	GH ₃₃	$32(c+d)GH_{30}$	
56	HG ₃₀	$16H_{60} = 32\{-dHH_{20} + aHG_{20} + bHH_{30} + cHG_{30}\}$	
.		.	
63	GG ₃₃	$16G_{63} = 32\{-dGH_{23} + aGG_{23} + bGH_{33} + cGG_{33}\}$	
64		.	
.		.	
71		.	
72		.	
.		.	
79		.	

FIG. 27D

CONV_COL DATA FLOW FOR THE
INVERSE OCTAVE 0 TRANSFORM

$32\{-dHH_{20} + aHG_{20} + bHH_{30}\}$	$32\{cHH_{20} - bHG_{20} + aHH_{30}\}$
$32\{-dGH_{21} + aGG_{21} + bGH_{31}\}$	$32\{cGH_{21} - bGG_{21} + aGH_{31}\}$
$32(c+d)HH_{30} - 32(a+b)HG_{30}$	
$32(c+d)GH_{33} - 32(a+b)GG_{33}$	
	$32(c+d)HH_{30} - 32(a+b)HG_{30}$
	$32(c+d)GH_{33} - 32(a+b)GG_{33}$

FIG. 27E

CONV_COL DATA FLOW FOR THE INVERSE OCTAVE 0 TRANSFORM

$$32\{cHH_{20} - bHG_{20} + aHH_{30} + dHG_{30}\} = 16H_{50}$$
$$32\{cGH_{23} - bGG_{23} + aGH_{33} + dGG_{33}\} = 16G_{53}$$
$$32\{(c+d)HH_{30} - (a+b)HG_{30}\} = 8H_{70}$$
$$32\{(c+d)GH_{33} - (a+b)GG_{33}\} = 8G_{73}$$

KEY TO FIG. 27

FIG. 27A	FIG. 27B	FIG. 27C
FIG. 27D	FIG. 27E	FIG. 27F

FIG. 27

FIG. 27F

CONV_COL DATA FLOW FOR THE INVERSE OCTAVE 0 TRANSFORM

Time t	0	1	2	3	4	5
Input Data Value	H ₀₀	G ₀₀	H ₀₁	G ₀₁	H ₀₂	G ₀₂
muxsel(1)	2	4	2	2	2	2
muxsel(2)	3	1	1	1	1	1
muxsel(3)	1	1	1	1	1	1
andsel(1),(4)	zero	pass	zero	pass	zero	pass
andsel(2)	zero	pass	pass	pass	pass	pass
andsel(3)	pass	pass	pass	pass	pass	pass
addsel(1)	add	add	add	add	add	add
addsel(2)	add	sub	add	sub	add	sub
addsel(3)	add	add	add	add	add	add
addsel(4)	sub	add	sub	add	sub	add
centermuxsel (1)	r	l	r	l	r	l
centermuxsel (2)	l	r	l	r	l	r
muxandsel (1)	pass	pass	pass	pass	pass	pass
muxandsel (2)	zero	pass	pass	pass	pass	pass
muxandsel (3)	pass	zero	pass	pass	pass	pass
OUT 2					16D ₀₁	
OUT 1			8D ₀₀		16D ₀₂	
OUTPUT LEADS 520				D ₀₀	D ₀₁	D ₀₂

FIG. 28A
 CONV_ROW CONTROL SIGNALS AND
 OUTPUTS DURING INVERSE OCTAVE 0 TRANSFORM

6	7	8	9	10
H ₀₃	G ₀₃	H ₁₀	G ₁₀	H ₁₁
3	2	2	4	2
1	2	3	1	1
1	1	1	1	1
zero	pass	zero	pass	zero
pass	pass	zero	pass	pass
pass	pass	pass	pass	pass
add	add	add	add	add
add	sub	add	sub	add
add	add	add	add	add
sub	add	sub	add	sub
r	l	r	l	r
l	r	l	r	l
pass	pass	pass	pass	pass
pass	pass	zero	pass	pass
pass	pass	pass	zero	pass
16D ₀₃		16D ₀₅		8D ₀₇
16D ₀₄		16D ₀₆		8D ₁₀
D ₀₃	D ₀₄	D ₀₅	D ₀₆	D ₀₇

KEY TO Fig. 28

Fig. 28A	Fig. 28B
----------	----------

Fig. 28

FIG. 28B

CONV_ROW CONTROL SIGNALS AND
 OUTPUTS DURING INVERSE OCTAVE 0 TRANSFORM

Time	Data Values	Output of Block 926	Output of Block 928
0	H ₀₀	32cH ₀₀	32(b-a)H ₀₀
1	G ₀₀	8D ₀₀ = 32((b-a)H ₀₀ + (c-d)G ₀₀)	32(cH ₀₀ - bG ₀₀)
2	H ₀₁	32cH ₀₁	32(-dH ₀₀ + aG ₀₀ + bH ₀₁)
3	G ₀₁	16D ₀₂ = 32(-dH ₀₀ + aG ₀₀ + bH ₀₁ + cG ₀₁)	32(cH ₀₁ - bG ₀₁)
4	H ₀₂	32cH ₀₂	32(-dH ₀₁ + aG ₀₁ + bH ₀₂)
5	G ₀₂	16D ₀₄ = 32(-dH ₀₁ + aG ₀₁ + bH ₀₂ + cG ₀₂)	32(cH ₀₂ - bG ₀₂)
6	H ₀₃	32(c+d)H ₀₃	32(-dH ₀₂ + aG ₀₂ + bH ₀₃)
7	G ₀₃	16D ₀₆ = 32(-dH ₀₂ + aG ₀₂ + bH ₀₃ + cG ₀₃)	32((c+d)H ₀₃ - (a+b)G ₀₃)
8	H ₁₀	32cH ₁₀	32(b-a)H ₁₀
9	G ₁₀	8D ₁₀ = 32((b-a)H ₁₀ + (c-d)G ₁₀)	32(cH ₁₀ - bG ₁₀)
10			

FIG. 29A

CONV_ROW DATA FLOW FOR THE
INVERSE OCTAVE 0 TRANSFORM

Output of Block 930	Output of Block 932
	-32dH ₀₀
32{-dH ₀₀ + aG ₀₀ }	
32{cH ₀₀ - bG ₀₀ + aH ₀₁ }	32{-dH ₀₁ }
32{-dH ₀₁ + aG ₀₁ }	16D ₀₁ = 32{cH ₀₀ - bG ₀₀ + aH ₀₁ + dG ₀₁ }
32{cH ₀₁ - bG ₀₁ + aH ₀₂ }	32{-dH ₀₂ }
32{-dH ₀₂ + aG ₀₂ }	16D ₀₃ = 32{cH ₀₁ - bG ₀₁ + aH ₀₂ + dG ₀₂ }
32{cH ₀₂ - bG ₀₂ + aH ₀₃ }	
	16D ₀₅ = 32{cH ₀₂ - bG ₀₂ + aH ₀₃ + dG ₀₃ }
32{(c+d)H ₀₃ - (a+b)G ₀₃ }	32{-dH ₁₀ }
32{-dH ₁₀ + aG ₁₀ }	8D ₀₇ = 32{(c+d)H ₀₃ - (a+b)G ₀₃ }

KEY TO FIG. 29

FIG. 29A	FIG. 29B
----------	----------

FIG. 29

CONV_ROW DATA FLOW FOR THE
INVERSE OCTAVE 0 TRANSFORM

FIG. 29B

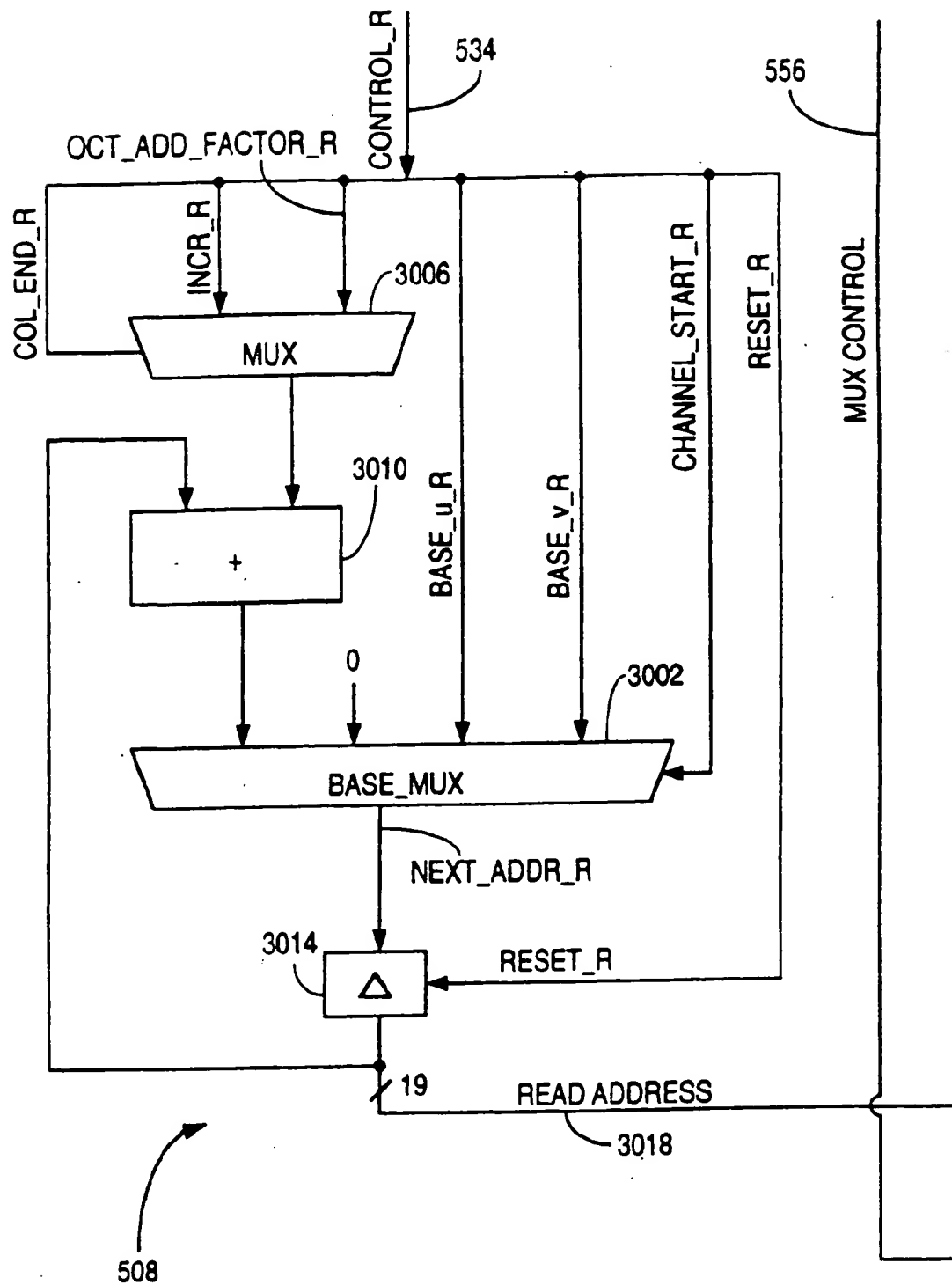
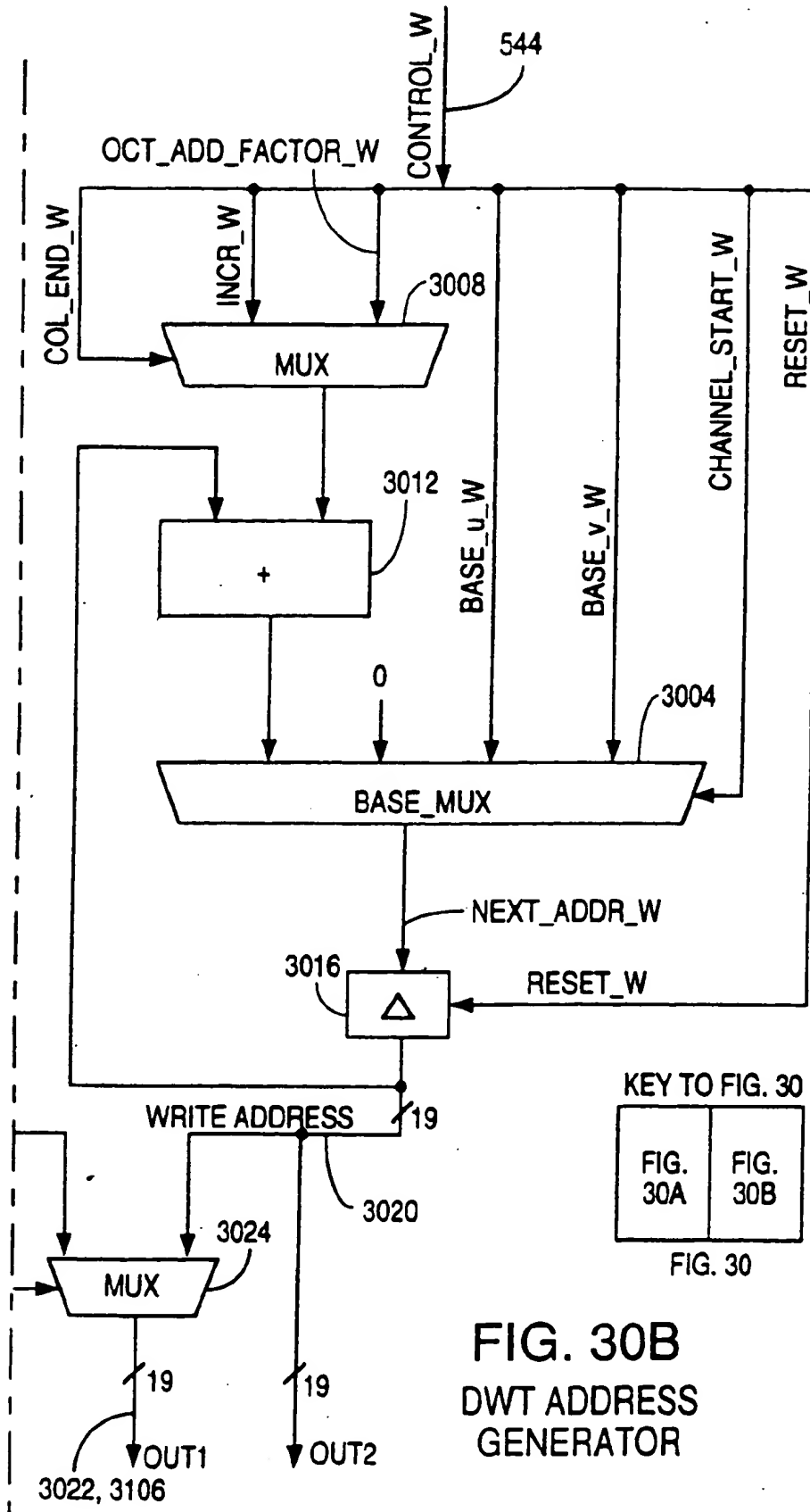
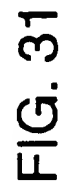
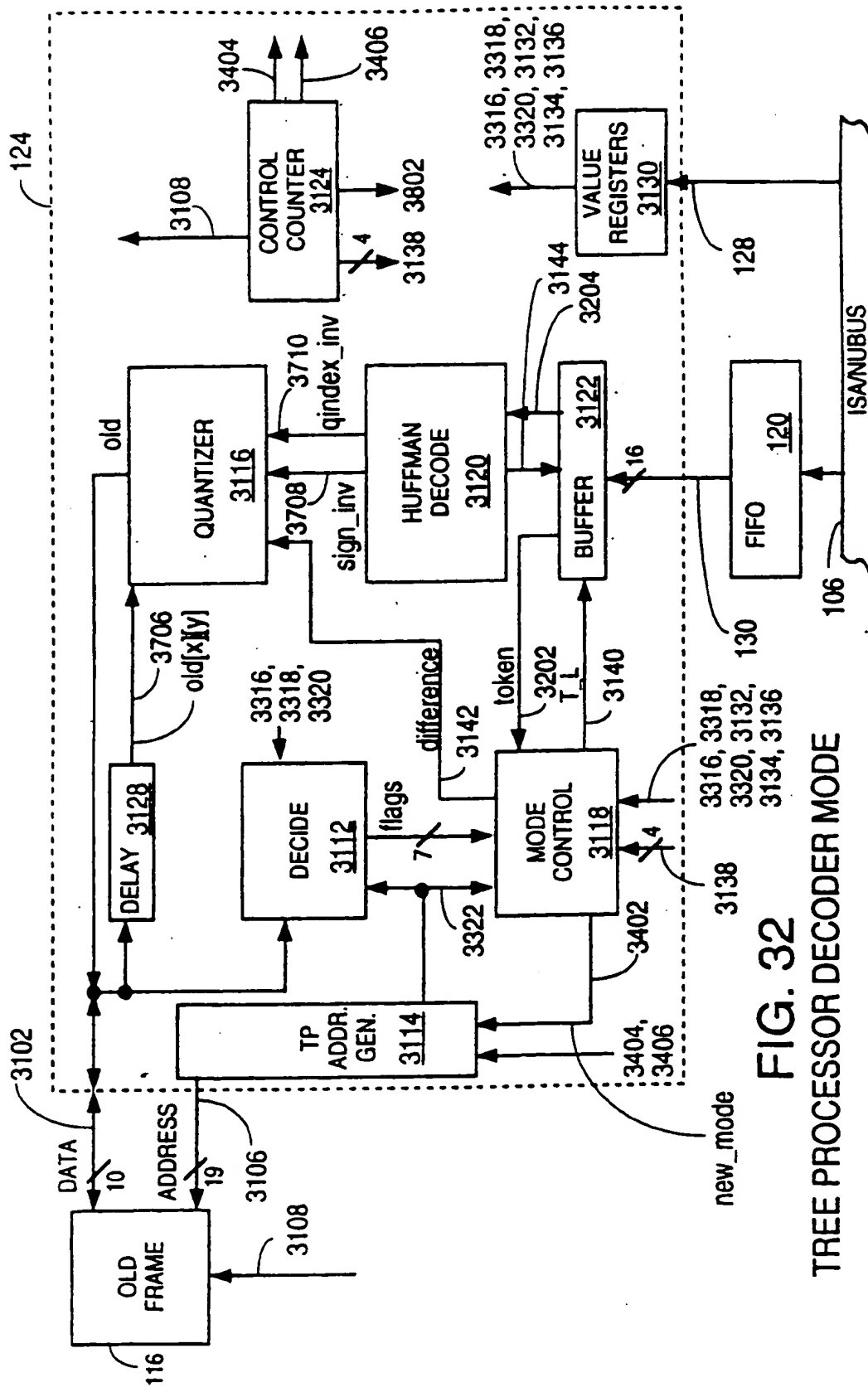


FIG. 30A
DWT ADDRESS
GENERATOR





TREE PROCESSOR ENCODER MODE



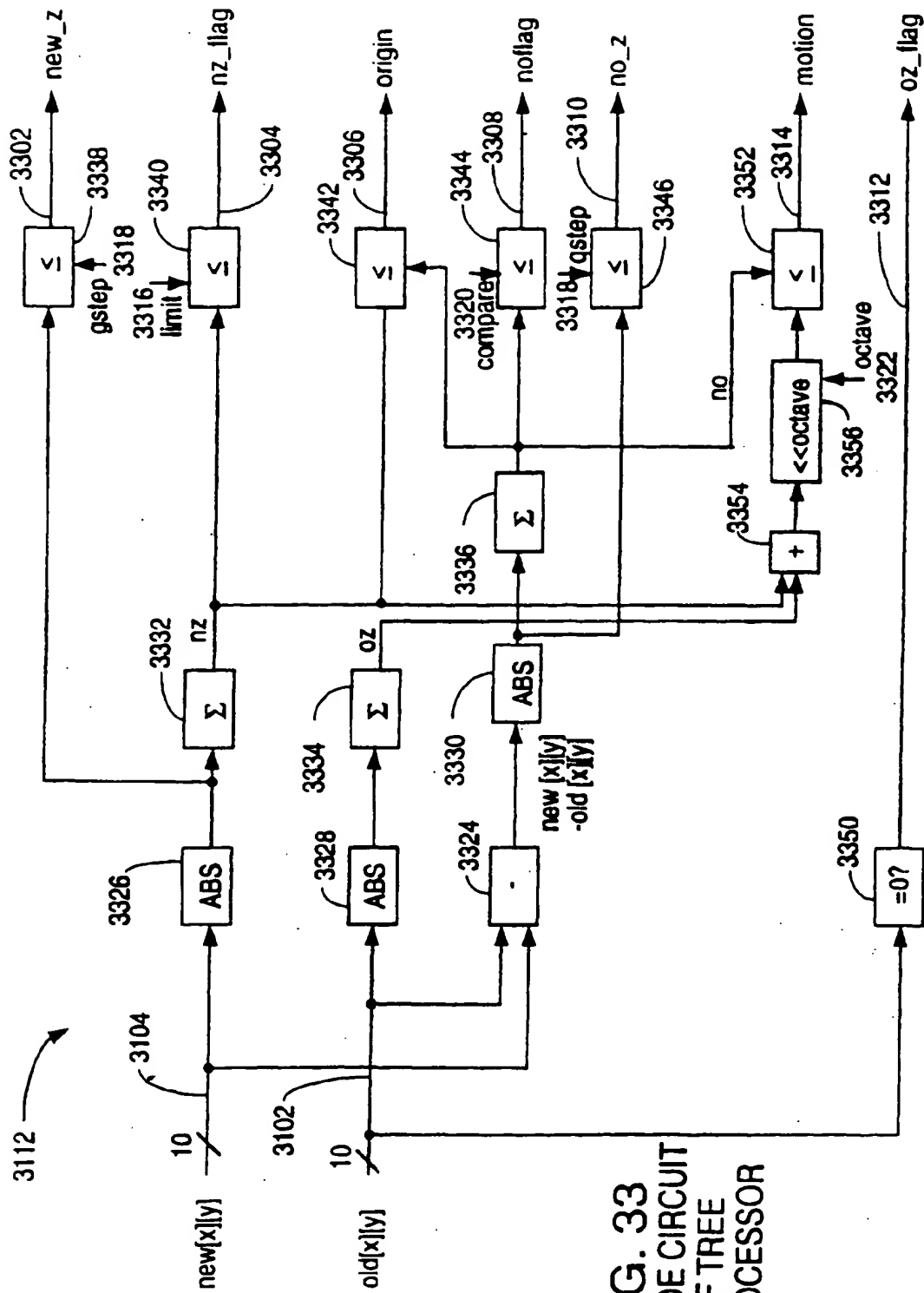


FIG. 33
DECIDE CIRCUIT
OF TREE
PROCESSOR

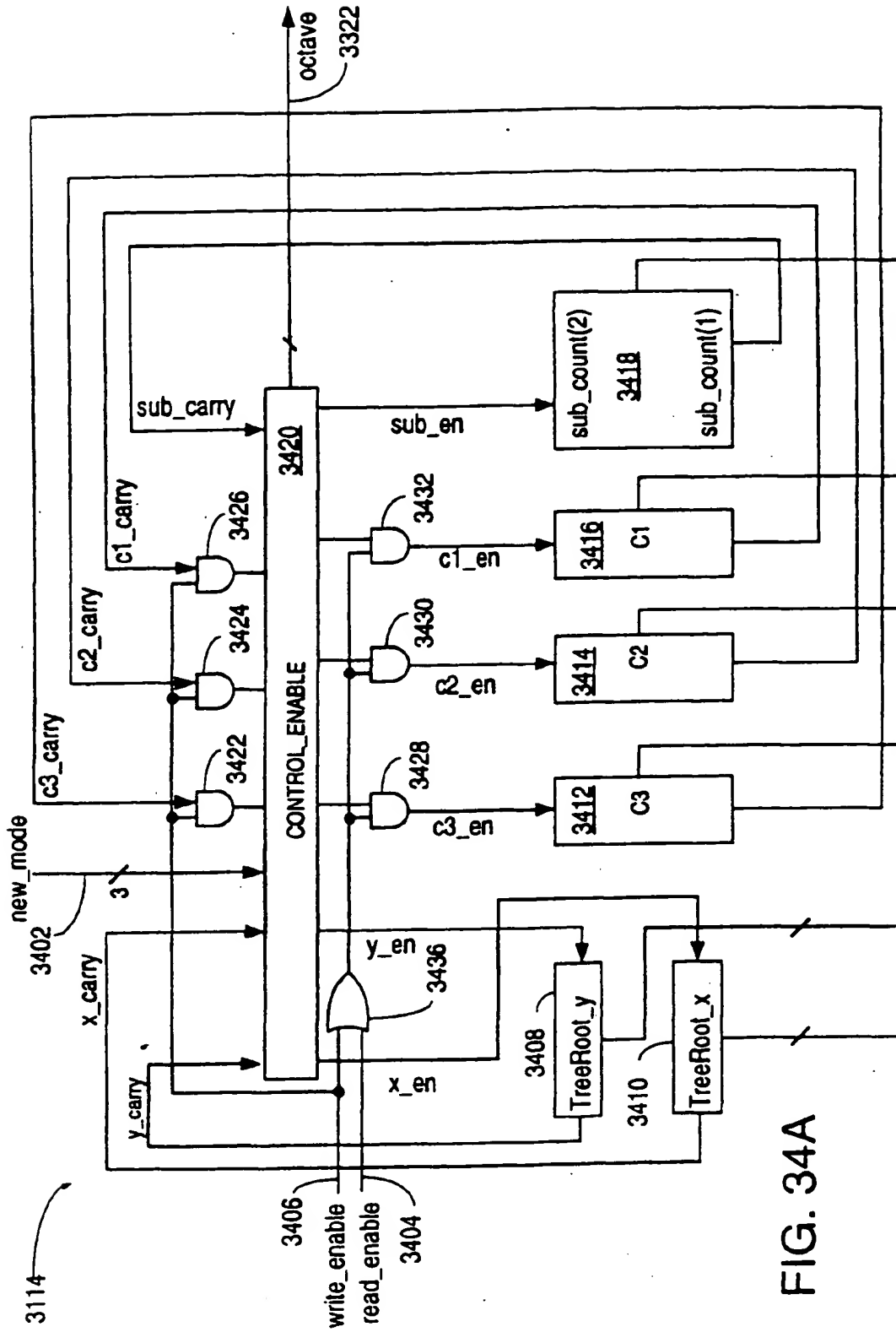


FIG. 34A

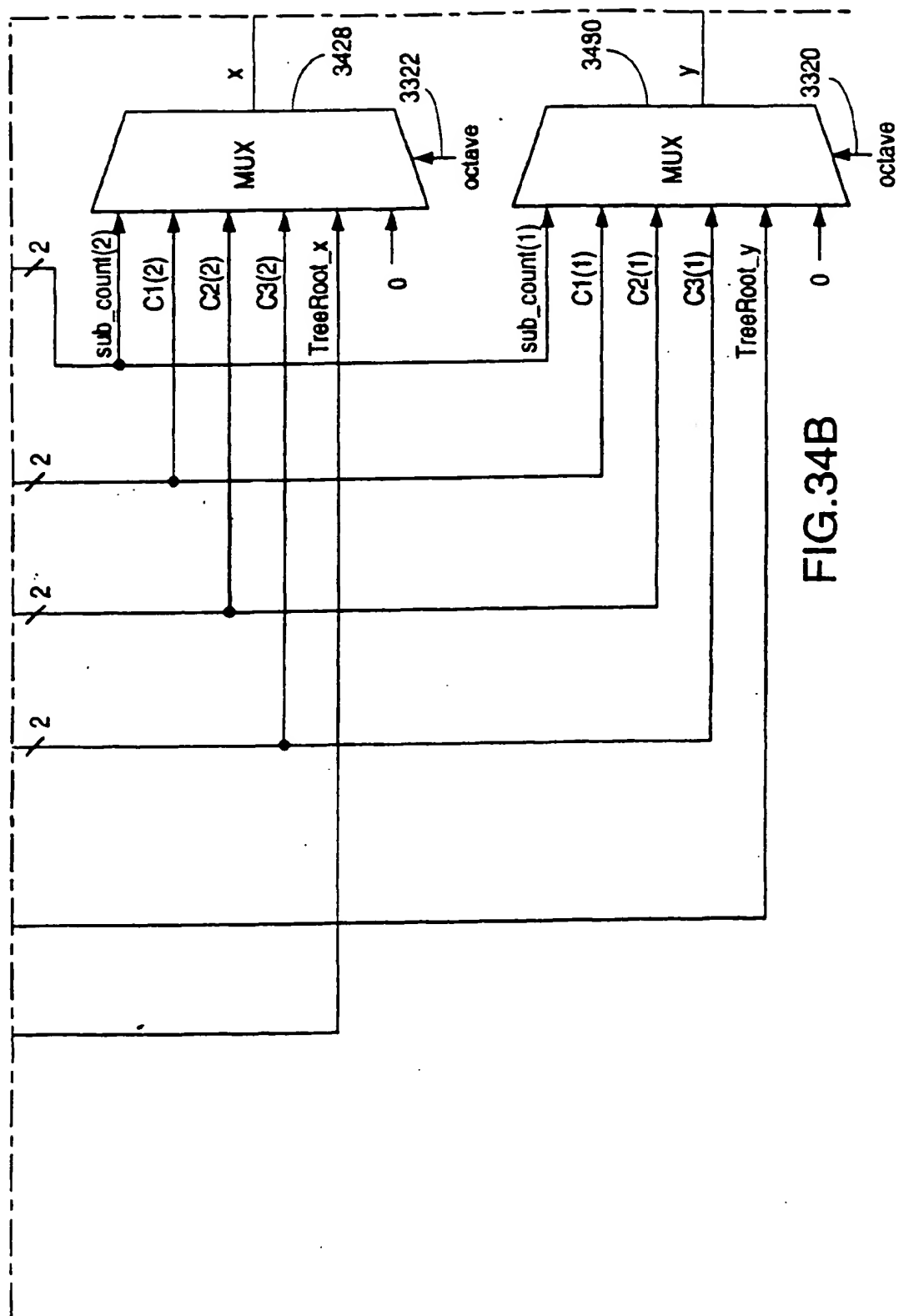
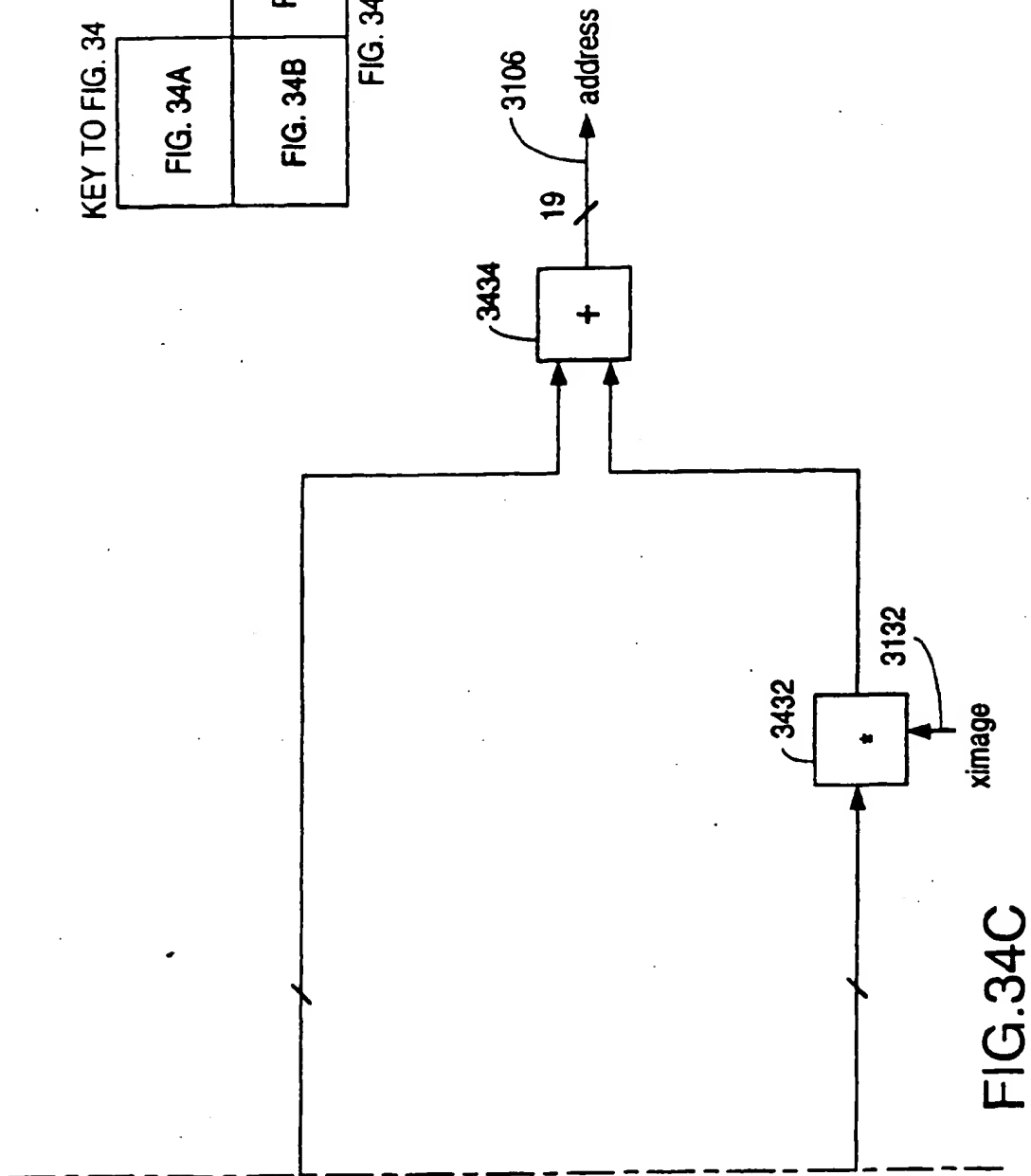
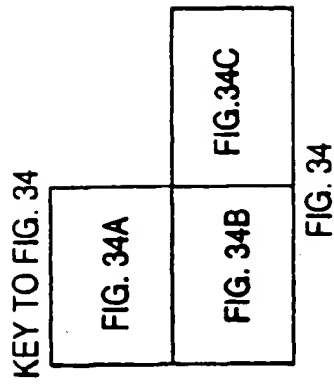


FIG. 34B



octave	STATE OF COUNTER C3	STATE OF COUNTER C2	STATE OF COUNTER C1	RELATIVE MOVEMENT BETWEEN TREE NODES
2	(0-3)	0	0	up0
1	0	(0-3)	0	up1
0	0	0	(0-3)	zz0
0	0	1	(0-3)	zz1
0	0	2	(0-3)	zz2
0	0	3	(0-3)	zz3
1	1	(0-3)	0	down1
0	1	0	(0-3)	zz0
0	1	1	(0-3)	zz1
0	1	2	(0-3)	zz2
0	1	3	(0-3)	zz3
1	2	(0-3)	0	down2
0	2	0	(0-3)	zz0
0	2	1	(0-3)	zz1
0	2	2	(0-3)	zz2
0	2	3	(0-3)	zz3
1	3	(0-3)	0	down3
0	3	0	(0-3)	zz0
0	3	1	(0-3)	zz1
0	3	2	(0-3)	zz2
0	3	3	(0-3)	zz3
2	(0-3)	0	0	up0
1	0	(0-3)	0	up1
.
.
.
.

FIG. 35
STATE TRANSITION DIAGRAM OF CONTROL
ENABLE BLOCK OF TREE PROCESSOR ADDRESS
GENERATOR

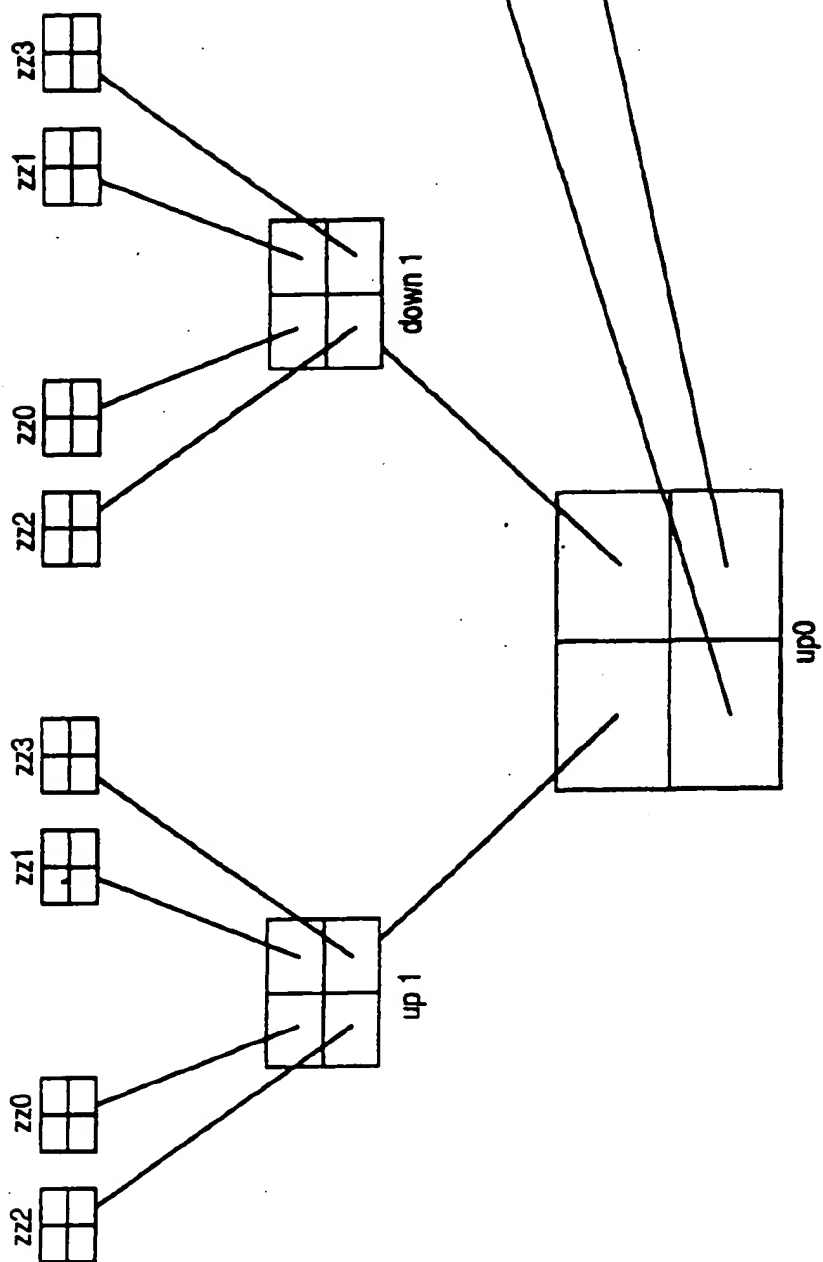
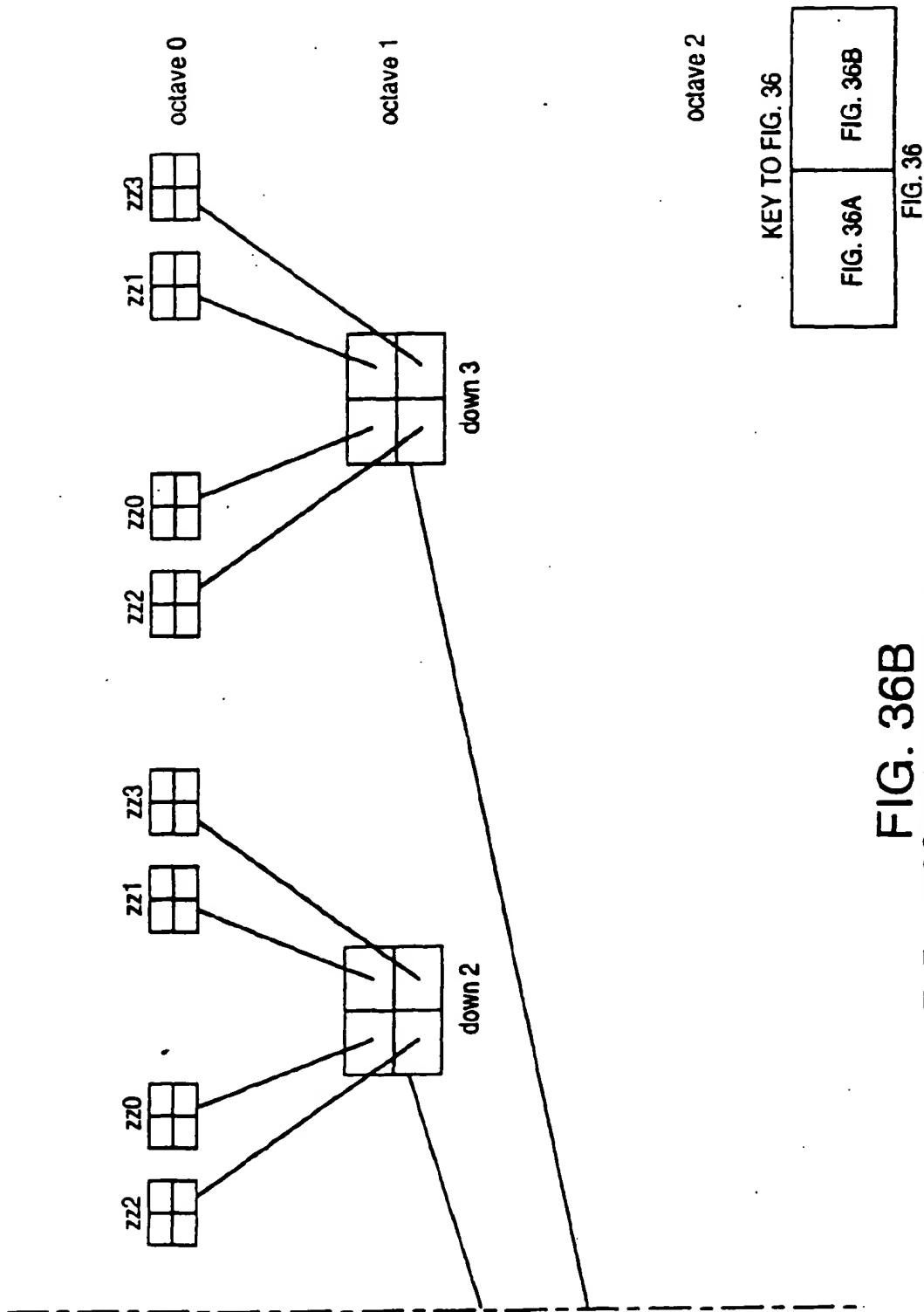
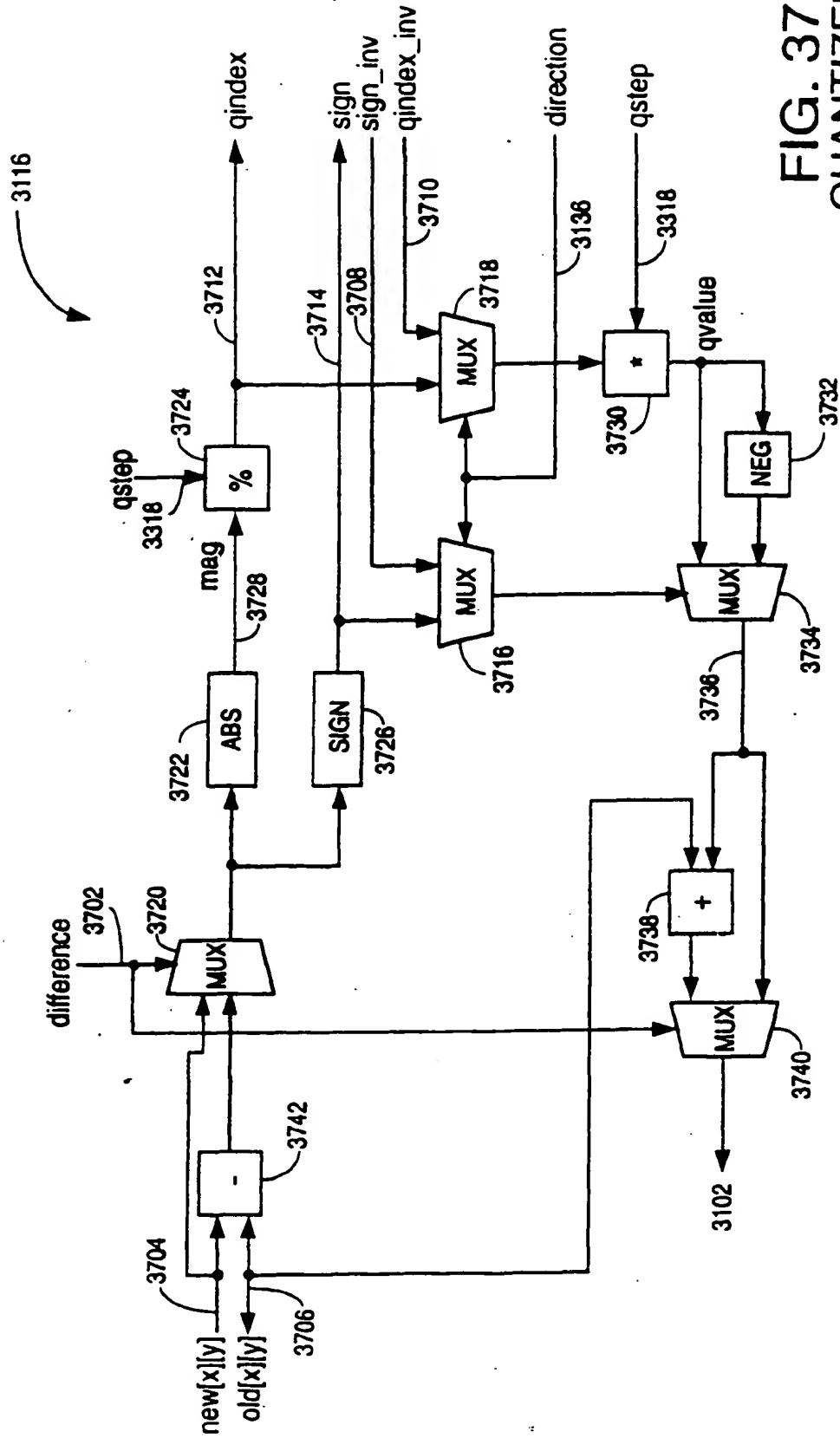


FIG. 36A
TREE DECOMPOSITION (3 OCTAVES)



FIG. 37
QUANTIZER

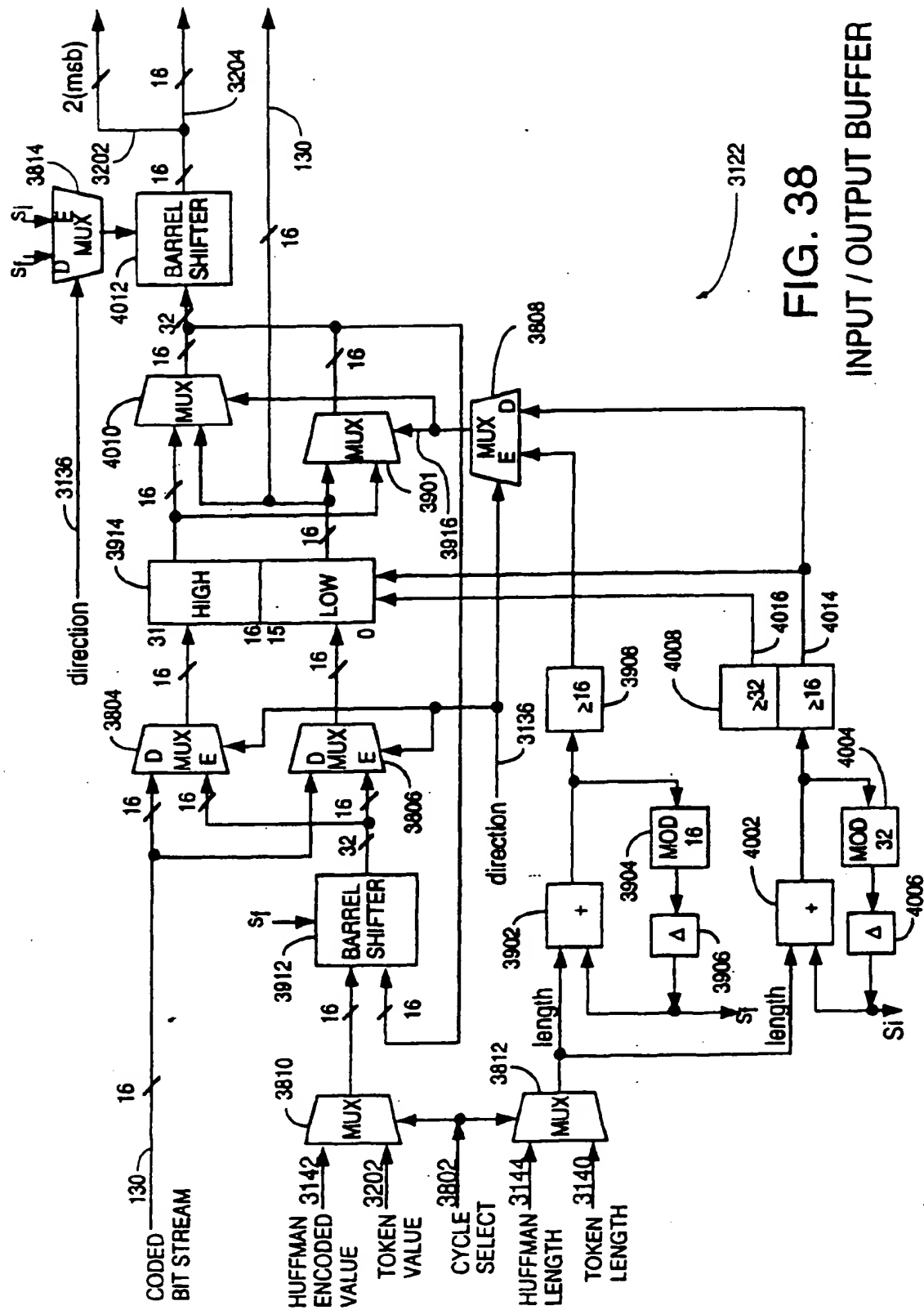


FIG. 38
INPUT / OUTPUT

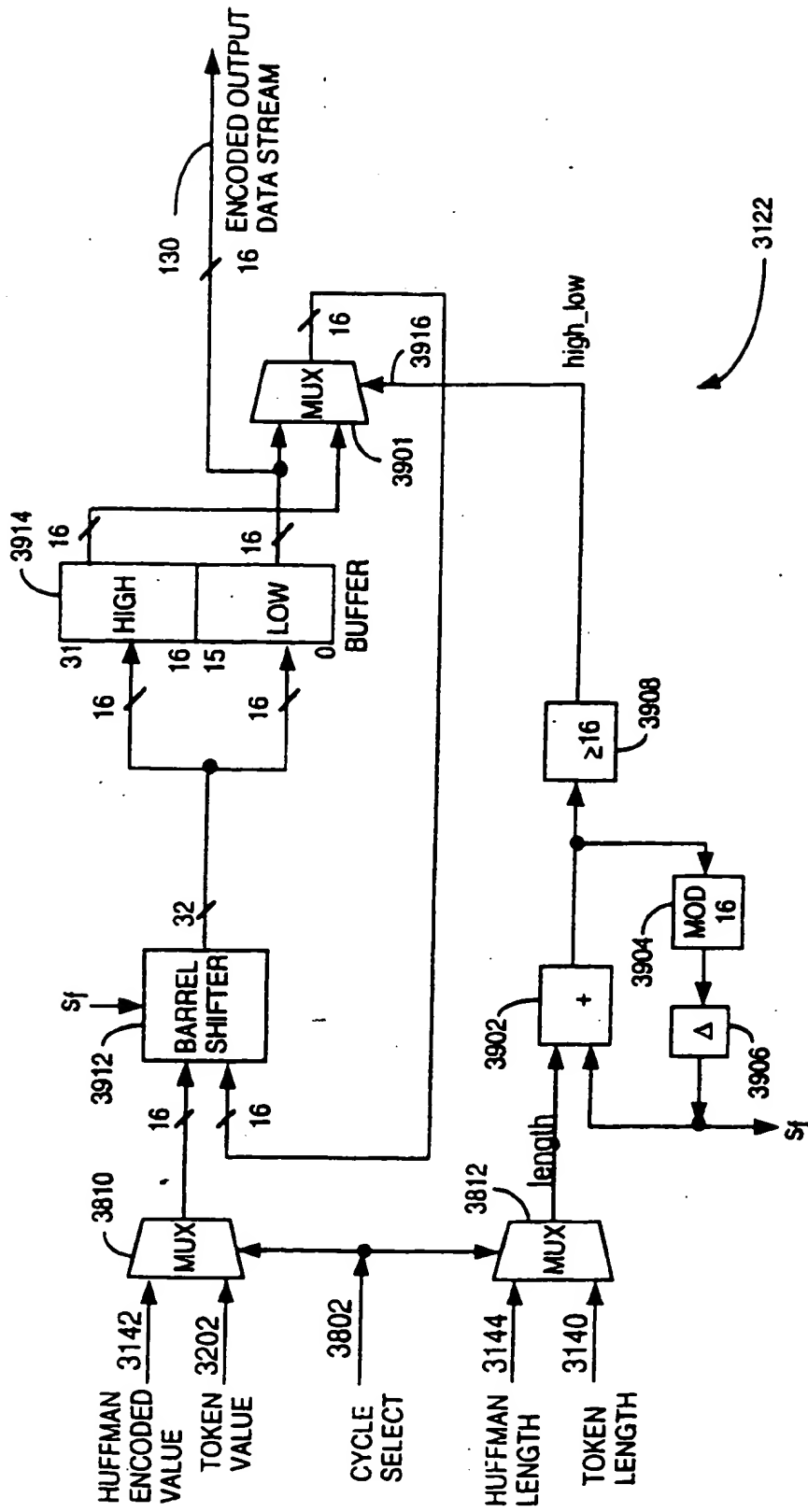


FIG. 39

BUFFER IN ENCODER MODE

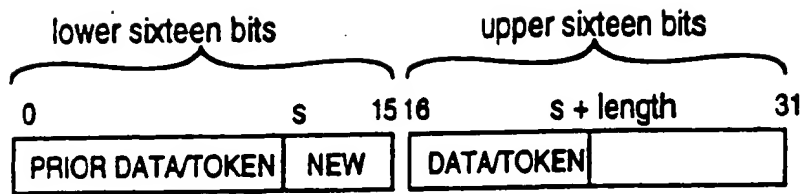
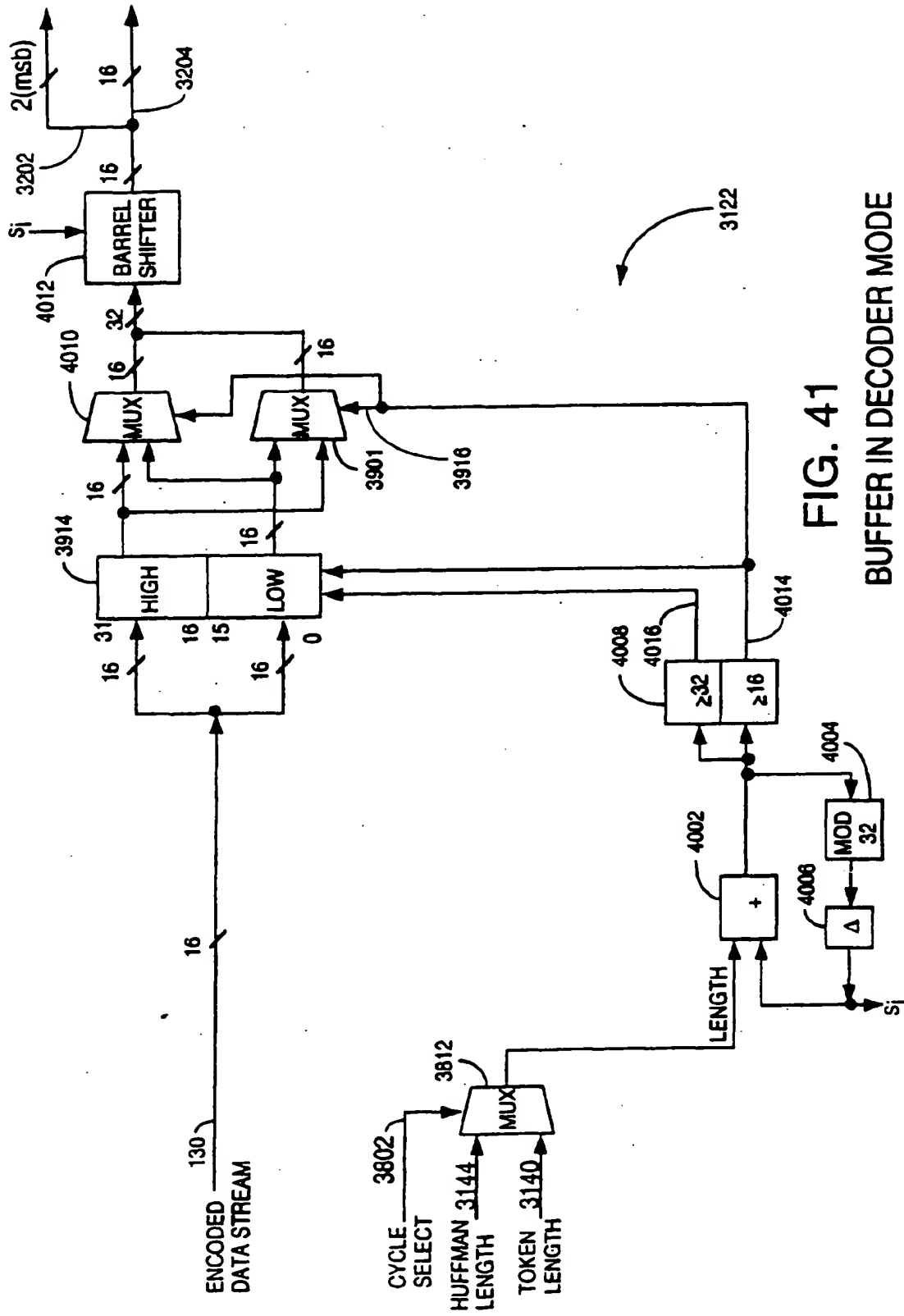
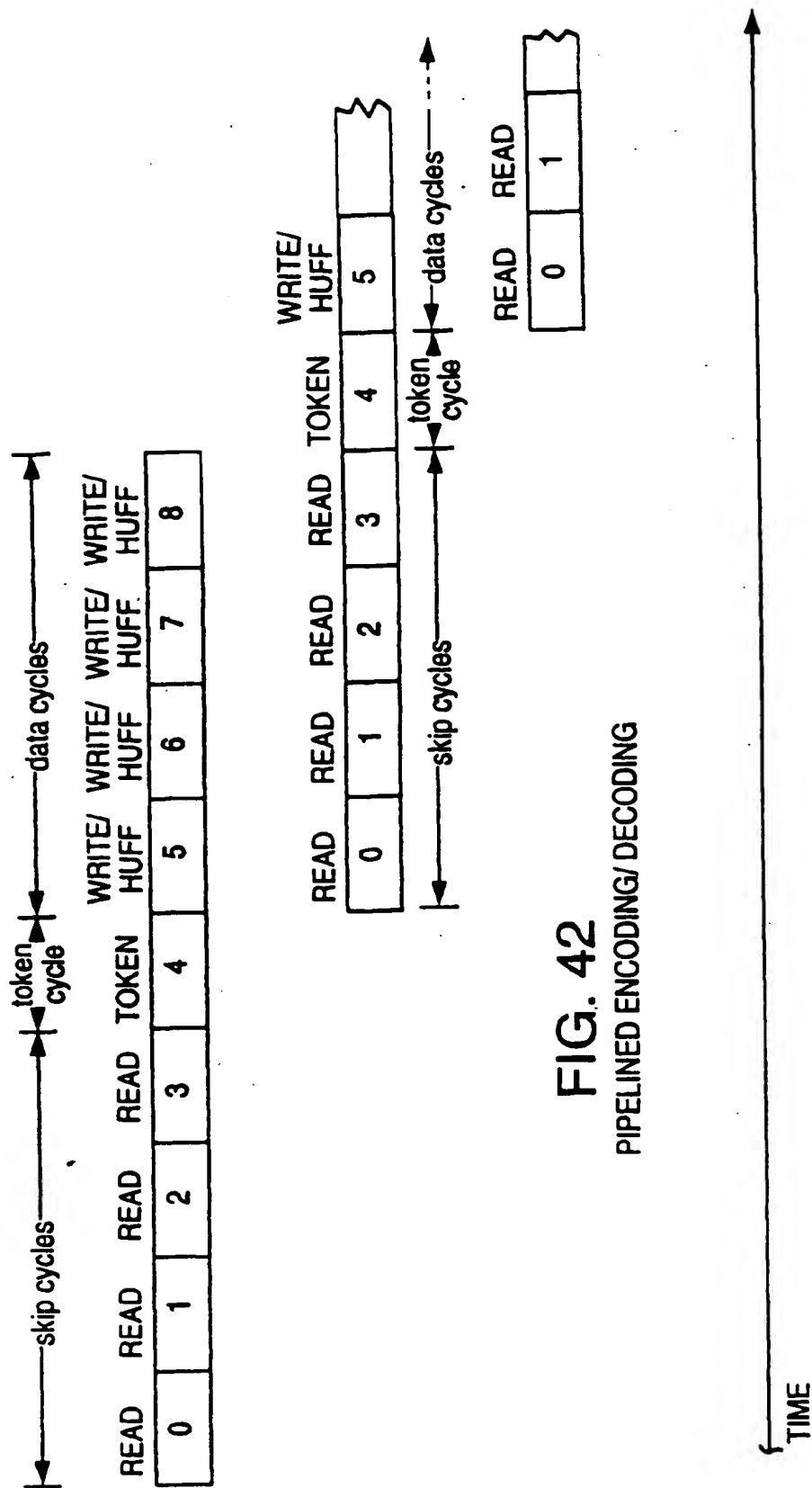


FIG. 40
OUTPUT OF BARREL SHIFTER
OF BUFFER BLOCK IN ENCODER MODE





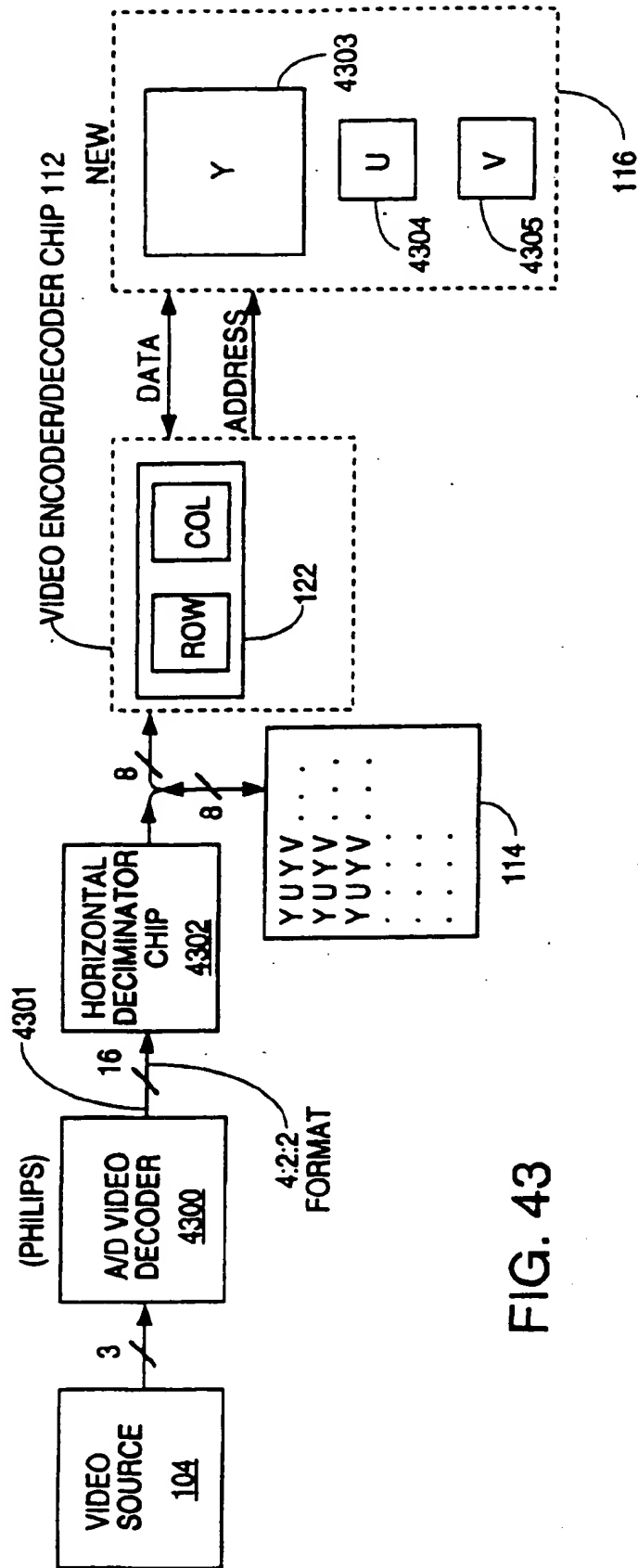
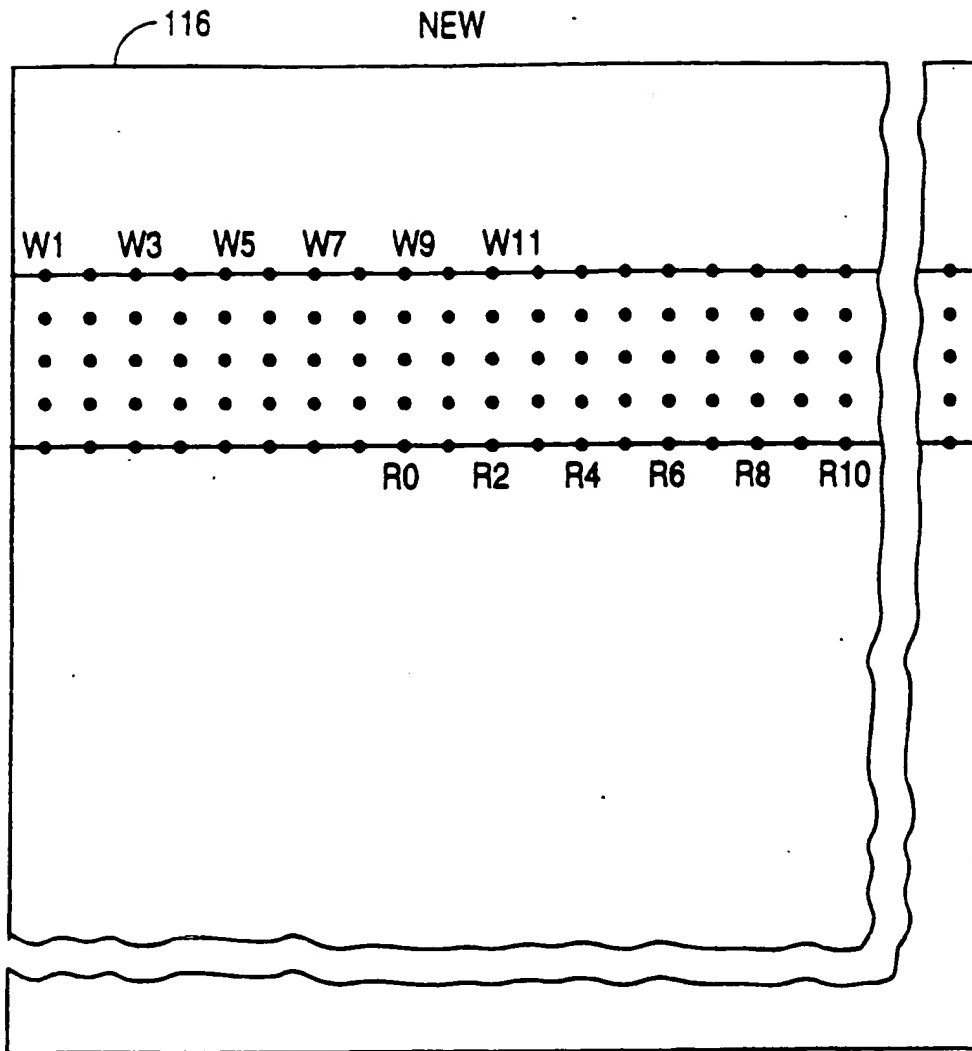
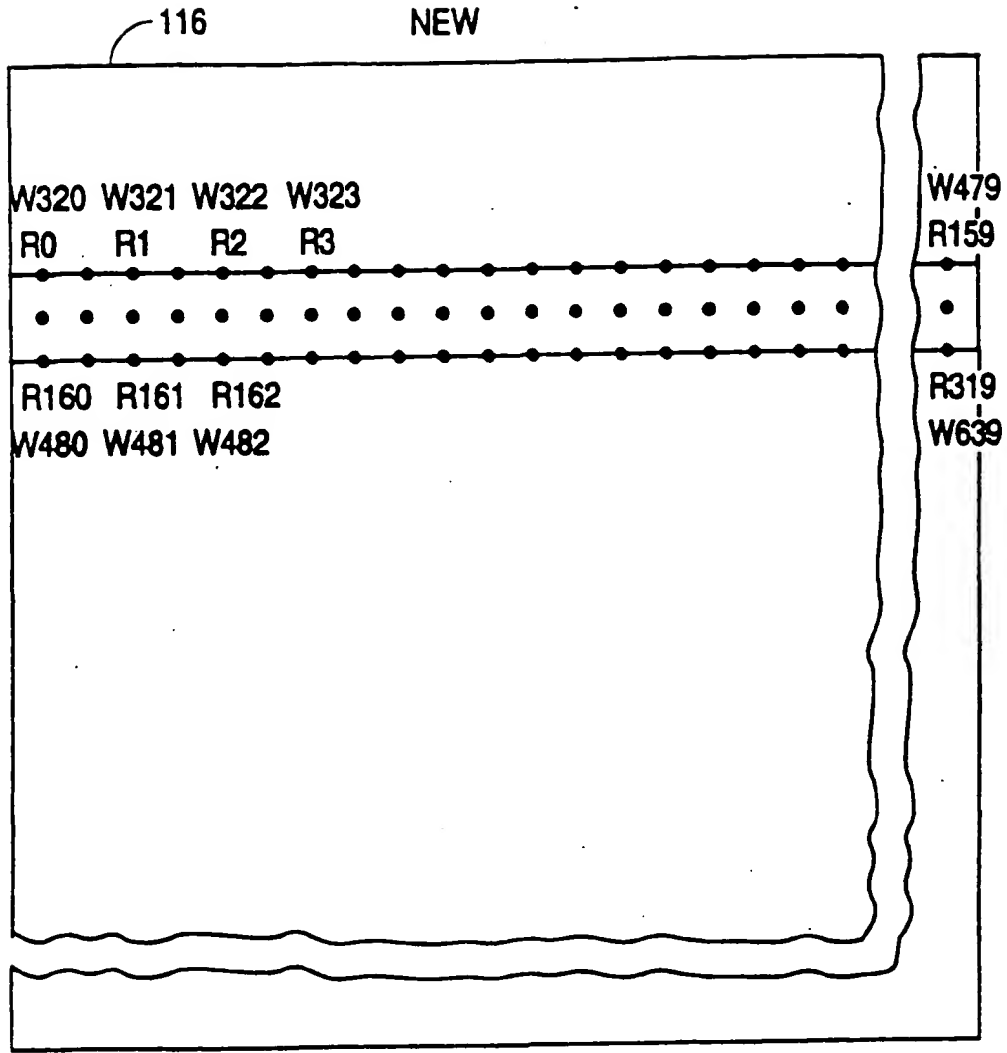


FIG. 43



OCTAVE 1

FIG. 44



OCTAVE 1

FIG. 45

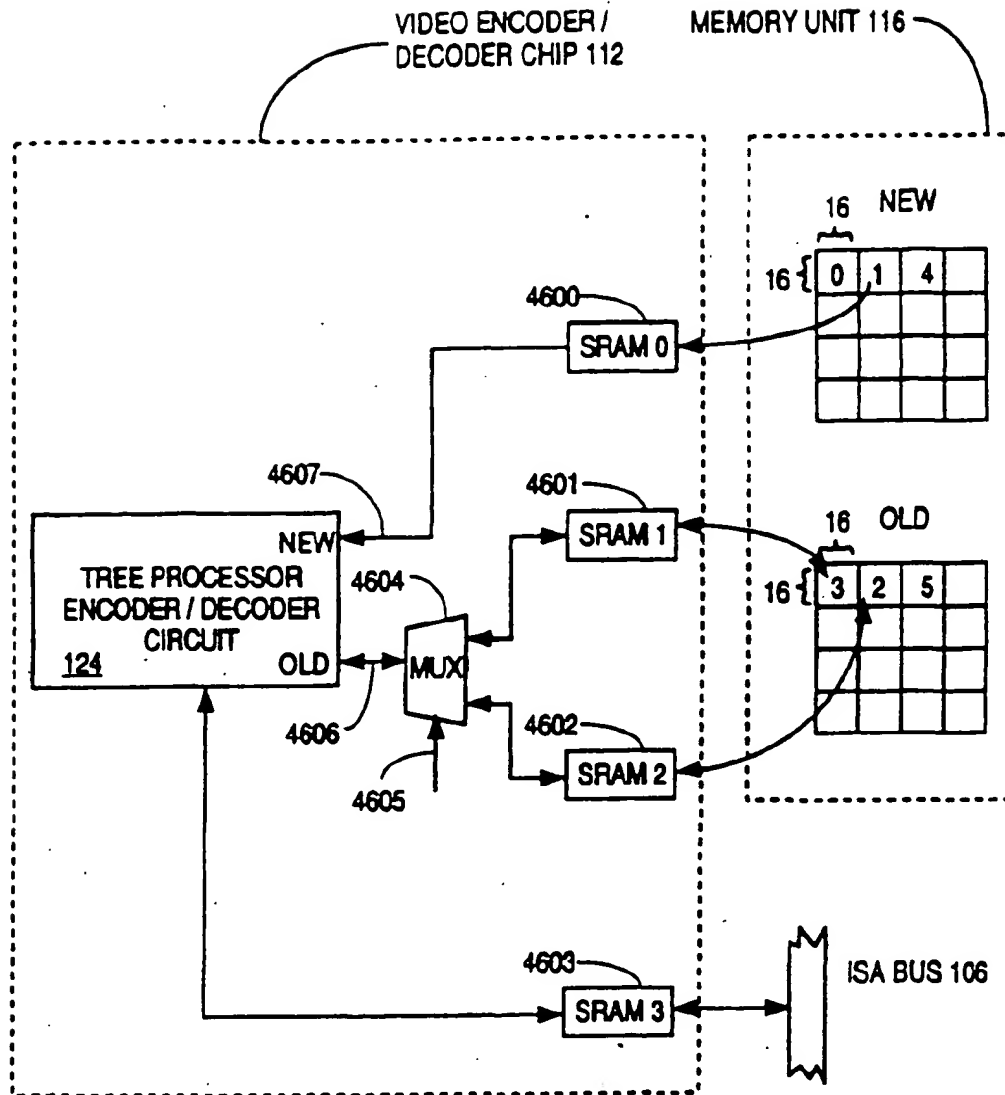


FIG. 46

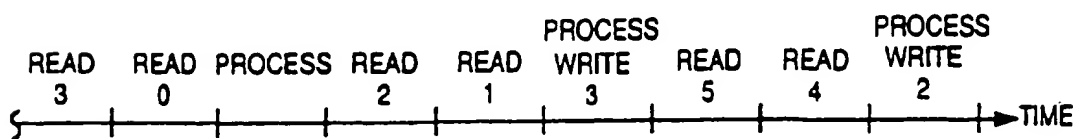


FIG. 47

**This Page is Inserted by IFW Indexing and Scanning
Operations and is not part of the Official Record**

BEST AVAILABLE IMAGES

Defective images within this document are accurate representations of the original documents submitted by the applicant.

Defects in the images include but are not limited to the items checked:

- ☐ BLACK BORDERS
- ☐ IMAGE CUT OFF AT TOP, BOTTOM OR SIDES
- ☐ FADED TEXT OR DRAWING
- ☐ BLURRED OR ILLEGIBLE TEXT OR DRAWING
- ☐ SKEWED/SLANTED IMAGES
- ☐ COLOR OR BLACK AND WHITE PHOTOGRAPHS
- ☐ GRAY SCALE DOCUMENTS
- ☐ LINES OR MARKS ON ORIGINAL DOCUMENT
- ☐ REFERENCE(S) OR EXHIBIT(S) SUBMITTED ARE POOR QUALITY
- ☐ OTHER: _____

IMAGES ARE BEST AVAILABLE COPY.

As rescanning these documents will not correct the image problems checked, please do not report these problems to the IFW Image Problem Mailbox.